Instructions. Limited collaboration is allowed while solving problems, but you must write solutions yourself. List collaborators on your submission.

You can choose which problems to complete, but must submit at least one problem per assignment. See the course page for information about how challenge problems are graded and contribute to your homework grade. Since you don’t need to complete every problem, you are encouraged to focus your efforts on producing high-quality solutions to the problems you feel confident about. There is no benefit to guessing or writing vague answers.

If you are asked to design an algorithm, please (a) give a precise description of your algorithm using either pseudocode or language, (b) explain the intuition of the algorithm, (c) justify the correctness of the algorithm; give a proof if needed, (d) state the running time of your algorithm, (e) justify the running-time analysis.

Submissions. Please submit a PDF file. You may submit a scanned handwritten document, but a typed submission is preferred. Please assign pages to questions in Gradescope.

Problem 1. Stable Matching Running Time. In class, we saw that the Propose-and-reject algorithm terminates in at most $n^2$ iterations, when there are $n$ students and $n$ colleges.

1. It seems possible that the algorithm may complete in fewer rounds if the preference lists have a certain structure. Describe how to construct preference lists for any number $n$ of colleges and students, such that the propose-and-reject algorithm will complete in only $O(n)$ rounds when run on these preference lists.

2. Could it be the case that the running time is actually $O(n)$ for all preference lists? Show that this is not true by designing preference lists so that the number of rounds of the algorithm is $\Omega(n^2)$.

Problem 2. Stable Matchings: K&T Ch 1, Ex 5. Consider a version of the stable matching problem where there are $n$ students and $n$ colleges as before. Assume each student ranks the colleges (and vice versa), but now we allow ties in the ranking. In other words, we could have a school that is indifferent two students $s_1$ and $s_2$, but prefers either of them over some other student $s_3$ (and vice versa). We say a student $s$ prefers college $c_1$ to $c_2$ if $c_1$ is ranked higher on the $s$’s preference list and $c_1$ and $c_2$ are not tied.

1. **Strong Instability.** A strong instability in a matching is a student-college pair, each of which prefer each other to their current pairing. In other words, neither is indifferent about the switch. Does there always exist a matching with no strong instability? Either give an example instance for which all matchings have a strong instability (and prove it), or give and analyze an algorithm that is guaranteed to find a matching with no strong instabilities.

2. **Weak Instability.** A weak instability in a matching is a student-college pair where one party prefers the other, and the other may be indifferent. Formally, a student $s$ and a college $c$ with pairs $c'$ and $s'$ form a weak instability if either
   
   - $s$ prefers $c$ to $c'$ and $c$ either prefers $s$ to $s'$ or is indifferent between $s$ and $s'$.
   - $c$ prefers $s$ to $s'$ and $s$ either prefers $c$ to $c'$ or is indifferent between $c$ and $c'$.

   Does there always exist a perfect matching with no weak instability? Either give an example instance for which all matchings have a weak instability (and prove it), or give and analyze an algorithm that is guaranteed to find a matching with no weak instabilities.
Problem 3. Asymptotics. K&T Ch 2, Ex 6. Given an array $A$ of $n$ integers, you’d like to output a two-dimensional $n \times n$ array $B$ in which $B[i, j] = \max\{A[i], A[i+1], \ldots, A[j]\}$ for each $i < j$.
For $i \geq j$ the value of $B[i, j]$ can be left as is.

\begin{align*}
\text{for } i = 1, 2, \ldots, n \\
\quad \text{for } j = i + 1, \ldots, n \\
\quad \quad \text{Compute the maximum of the entries } A[i], A[i+1], \ldots, A[j]. \\
\quad \quad \text{Store the maximum value in } B[i, j].
\end{align*}

1. Find a function $f$ such that the running time of the algorithm is $O(f(n))$, and clearly explain why.
2. For the same function $f$ argue that the running time of the algorithm is also $\Omega(f(n))$. (This establishes an asymptotically tight bound $\Theta(f(n))$.)
3. Design and analyze a faster algorithm for this problem. You should give an algorithm with running $O(g(n))$, where $\lim_{n \to \infty} g(n)/f(n) = 0$.

Problem 4. Highest Safe Rung. You want to stress test glass jars. You have a ladder with $n$ rungs, and want to find the highest rung from which you can drop a jar and not have it break. We call this the highest safe rung.

Your goal is to find the highest safe rung with the fewest number of drops possible. However, you have a limited supply of jars, and need to find the highest safe rung before breaking all of them.

For example, with one jar you could drop it from the first rung, then the second, then the third, and so on, until it breaks, and you are guaranteed to find the highest safe rung with at most $n$ drops. With any other strategy, you would risk breaking the jar before finding the highest safe rung.

Now, suppose you have two (identical) jars, so you can break one and still find the highest safe rung. Describe a strategy for finding the highest safe rung that requires you to drop a jar at most $f(n)$ times, for some function $f(n)$ that grows slower than linearly. (In other words, it should be the case that $\lim_{n \to \infty} f(n)/n = 0$.)