

CS 103: Lecture 8 Matching Markets

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Announcements

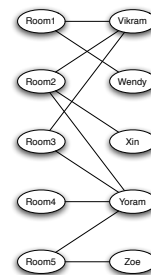
- ▶ Happy Mountain Day / Fall Break
- ▶ HW 3 assigned next week

Plan for today

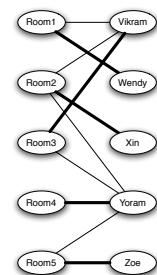
- ▶ Finish second-price auctions
- ▶ Matching markets
 - ▶ Perfect matchings
 - ▶ Constricted sets
 - ▶ Matching Theorem

Matching Market

Example: room preferences

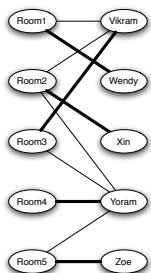


Bipartite graph



Perfect matching

Matching Market



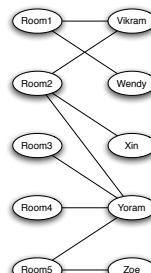
Bipartite graph

- ▶ two types of nodes
- ▶ edges from one type to other

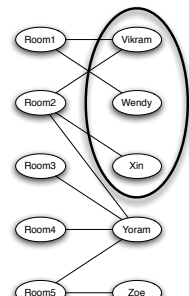
Perfect matching: each node L assigned to a single node on the R to which it is connected by an edge

- ▶ Baristas / shifts
- ▶ Classes / rooms
- ▶ Planes / gates

Perfect Matchings

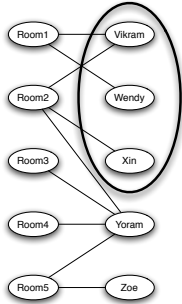


Perfect matching?



No: three people / two rooms

Constricted Set



S = set of nodes on the right
 $N(S)$ = set of left nodes with an edge to a node in S

S is **constricted** if $N(S)$ is smaller than S

Significance: if bipartite graph has a constricted set, then it has no perfect matching. Can there be other reasons?

Exercise

Draw examples on board

Find a perfect matching or constricted set

Matching Theorem (König 1931 / Hall 1935)

Matching Theorem: if a bipartite graph has no perfect matching, then it must contain a constricted set.

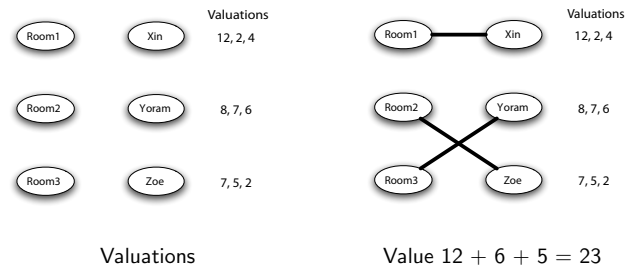
Discussion: very clean conceptual picture for resource allocation!

- ▶ perfect matching: problem solved!
- ▶ constricted set: simple explanation why it is impossible

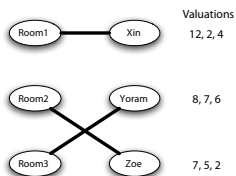
What is missing?

- ▶ 0/1 preferences: like a room or not
- ▶ want finer-grained notion of preference

Valuations



Optimal Assignment



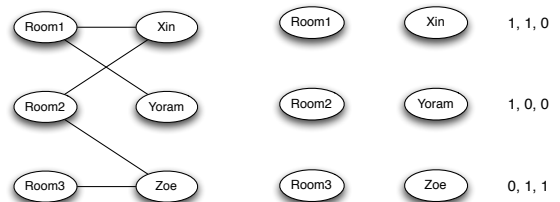
Total valuation of an assignment:
sum of each individual's valuation

Optimal assignment:
assignment that maximizes total valuation

Is this assignment optimal? Yes.

Do people get most preferred rooms? No.

Optimal Assignment / Perfect Matching



Optimal assignment problem generalizes perfect matching