

ON RECURRENCE AND TRANSIENCE IN HEAVY-TAILED GSMP'S

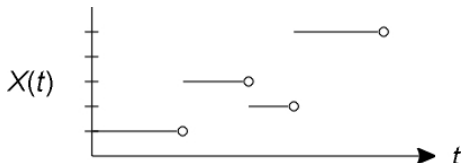
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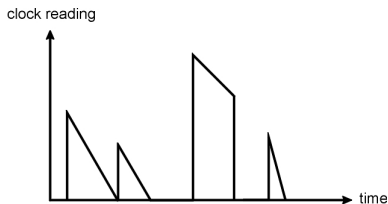
Discrete-Event Stochastic Systems



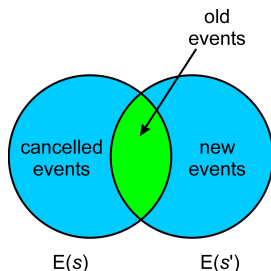
- ▶ System changes **state** when **events** occur
 - ▶ Stochastic state changes
 - ▶ At strictly increasing sequence of random times
- ▶ Underlying stochastic process $\{X(t): t \geq 0\}$
 - ▶ $X(t)$ = state of system at time t (a random variable)
 - ▶ Piecewise-constant sample paths
 - ▶ Typically **not** a continuous-time Markov chain

Generalized Semi-Markov Processes [Matthes62, Whitt80]

- ▶ Classical model for discrete-event stochastic systems
 - ▶ Subsumes: queueing networks, SMPs, CTMCs, SPNs
 - ▶ Central to simulation theory
- ▶ Building blocks
 - ▶ S = set of **states** (finite or countably infinite)
 - ▶ E = set of **events** (finite)
 - ▶ $E(s)$ = **active events** in state s
 - ▶ $p(s'; s, e^*)$ = **state-transition probability**
- ▶ One **clock** per event: records remaining time until occurrence



Clocks (Event Scheduling)

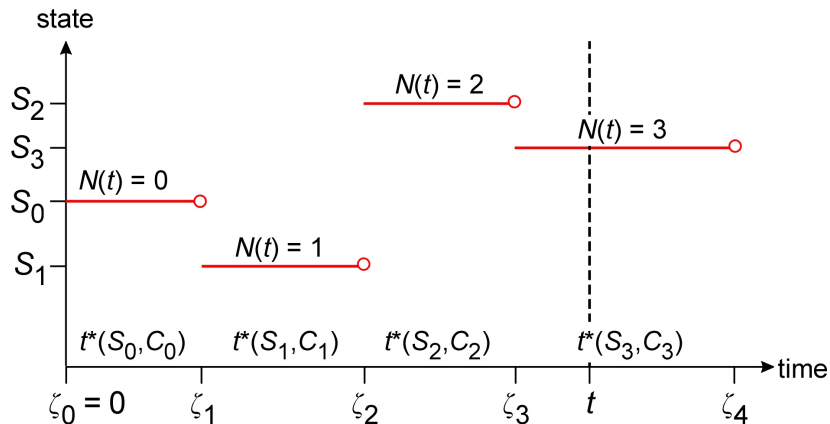


- ▶ Active events **compete** to trigger state transition
 - ▶ The clock that runs down to 0 first is the “winner”
 - ▶ Can have simultaneous event occurrence: $p(s'; s, E^*)$
- ▶ At a state transition $s \xrightarrow{e^*} s'$: three kinds of events
 - ▶ **New events**: Clock for e' set according to $F(x; e')$
 - ▶ **Old events**: Clocks continue to run down
 - ▶ **Cancelled events**: Clock readings are discarded

The GSMP Process

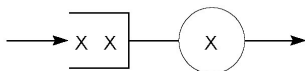
- ▶ **The continuous-time process:** $\{X(t) : t \geq 0\}$
 - ▶ $X(t)$ = the state at time t
 - ▶ A very complicated process
- ▶ Defined in terms of **Markov chain** $\{(S_n, C_n) : n \geq 0\}$
 - ▶ System observed after the n th state transition
 - ▶ S_n = the state
 - ▶ $C_n = (C_{n,1}, \dots, C_{n,M})$ = the clock-reading vector
 - ▶ Chain defined via GSMP building blocks
 - ▶ **initial distribution** μ on state and clocks

Definition of the GSMP



$$X(t) = S_{N(t)}$$

Example: GI/G/1 Queue



- ▶ $X(t)$ = “number of jobs in system at time t ”
- ▶ $S = \{0, 1, 2, \dots\}$
- ▶ $E = \{e_1, e_2\}$
 - ▶ e_1 = “arrival”
 - ▶ e_2 = “service completion”
- ▶ $E(0) = \{e_1\}$ and $E(s) = \{e_1, e_2\}$ for $s > 0$
- ▶ $p(s+1; s, e_1) = 1$ and $p(s-1; s, e_2) = 1$
- ▶ $F(\cdot; e_1) = F_{\text{interarrival}}$ and $F(\cdot; e_2) = F_{\text{service}}$

Harris Recurrence: A Basic Form of Stability

- ▶ Definition for general chain $\{Z_n : n \geq 0\}$ with state space Γ

$$P_z \{Z_n \in A \text{ i.o.}\} = 1, \quad z \in \Gamma \quad \text{whenever} \quad \phi(A) > 0$$

- ▶ ϕ is a **recurrence measure** (often “Lebesgue-like”)
- ▶ Every “dense enough” set is hit infinitely often w.p. 1
- ▶ No “wandering off to ∞ ”
- ▶ Chain admits invariant measure π_0 : $\int P(z, A) \pi_0(dz) = \pi_0(A)$
- ▶ **Positive** Harris recurrence:
 - ▶ Chain admits invariant **probability** measure π
 - ▶ $P_\pi \{Z_1 \in A\} = \pi(A)$
 - ▶ Implies stationarity when initial dist'n is π
- ▶ When is $\{(S_n, C_n) : n \geq 0\}$ (positive) Harris recurrent?

Some Stability Conditions

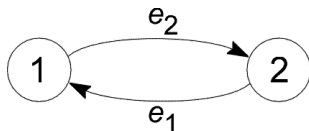
- ▶ **Density component** g of a cdf F : $F(t) \geq \int_0^t g(u) du$
- ▶ $s \rightarrow s'$ iff $p(s'; s, e) > 0$ for some e
- ▶ $s \rightsquigarrow s'$: either $s \rightarrow s'$ or $s \rightarrow s^{(1)} \rightarrow \dots \rightarrow s^{(n)} \rightarrow s'$
- ▶ **Assumption PD(q)**:
 - ▶ State space S is **finite**
 - ▶ GSMP is **irreducible**: $s \rightsquigarrow s'$ for all $s, s' \in S$
 - ▶ There exists $\bar{x} \in (0, \infty)$ s.t. all clock-setting dist'n functions
 - ▶ **Have finite q th moment**
 - ▶ Have density component positive on $[0, \bar{x}]$

Positive Harris Recurrence in Light-Tailed GSMPs [Haas99]

- ▶ $\bar{\phi}(\{s\} \times A) = \text{Lebesgue measure of } A \cap [0, \bar{x}]^M$
- ▶ **Theorem:** If Assumption PD(1) holds, then the (S_n, C_n) chain is positive Harris recurrent with recurrence measure $\bar{\phi}$
- ▶ Implies $P_\mu \{ S_n = s \text{ i.o.} \} = 1$ for all $s \in S$ with **finite** expected (continuous-time) hitting times
- ▶ Proof:
 - ▶ Show that chain is “ $\bar{\phi}$ -irreducible”
 - ▶ Establish Lyapunov drift condition and apply MC machinery (Meyn and Tweedie, 1993)
- ▶ Alternate approach to recurrence: geometric-trials arguments
 - ▶ Can drop positive-density assumption
 - ▶ Use detailed analysis of specific GSMP structure

Positive Recurrence in Heavy-Tailed GSMPs: It Depends

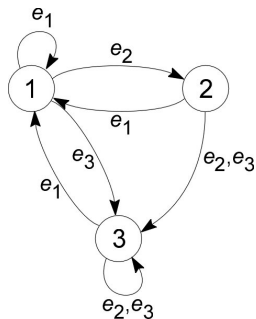
- ▶ Example 1: **Uninfluential events**
 - ▶ For all $s \in S$: $e \in E(s)$ and $p(s; s, e) = 1$
 - ▶ No effect on state or other clocks
 - ▶ If all heavy-tailed events are uninfluential and PD(1) holds otherwise, then positive recurrence
- ▶ Example 2: e_1 is heavy-tailed, so no positive-recurrence



Recurrence in Heavy-Tailed GSMPs

- ▶ S_n = state just after n th state transition
- ▶ Conjecture: $P\{S_n = s \text{ i.o.}\} = 1$ for each s under PD(0)
 - ▶ State space S is **finite**
 - ▶ GSMP is **irreducible**
 - ▶ $\exists \bar{x} > 0$ s.t. each $F(\cdot; e)$ has **positive density** on $(0, \bar{x})$
- ▶ Certainly true for CTMCs
- ▶ **CONJECTURE IS FALSE for GSMPs!**
 - ▶ In the presence of **heavy-tailed** clock-setting dist'ns

A Counterexample



- ▶ $S = \{1, 2, 3\}$ and $E = \{e_1, e_2, e_3\}$
- ▶ Event sets: $E(s) = \{e_1, e_2, e_3\}$ for all s (renewal processes)
- ▶ $p(s'; s, e^*) = 0$ or 1
- ▶ Clock-setting distributions:
 - ▶ $F(t; e_1) = 1 - (1 + t)^{-\alpha}$
 - ▶ $F(t; e_2) = 1 - (1 + t)^{-\beta}$
 - ▶ $F(\cdot; e_3)$ is Uniform $[0, a]$
- with $\beta > 1/2$ and $\alpha + \beta < 1$
- ▶ GSMP hits state $s = 2$ only if:
 - ▶ e_1 occurs and then e_2 occurs with no intervening occurrence of e_3
- ▶ Claim: $P\{S_n = 2 \text{ i.o.}\} = 0$
- ▶ Intuition: heavy clocks are rarely small simultaneously

Proof

- ▶ Observe: $P \{ S_n = 2 \text{ i.o.} \} = 0$ iff $P \{ B_n \text{ i.o.} \} = 0$
 - ▶ $B_n = \{ C_2(T_n) \leq C_3(T_n) \}$
 - ▶ $T_n = n$ th occurrence time of e_1
 - ▶ $C_i(t) =$ clock reading for e_i at time t
- ▶ Borel-Cantelli: suffices to show that $\sum_{n=1}^{\infty} P \{ B_n \} < \infty$
- ▶ Bound $P \{ B_n \}$ by an integral:

$$P \{ B_n \} \leq \int_0^{\infty} P \{ C_2(t) \leq a \} f_1^{*n}(t) dt$$

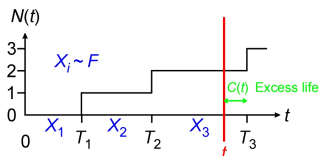
- ▶ $f_1^{*n} =$ density of $T_n = n$ -fold convolution of $f(\cdot; e_1)$
- ▶ Sum over n :

$$\sum_{n=1}^{\infty} P \{ B_n \} \leq \int_0^{\infty} h(t) u_1(t) dt$$

- ▶ $h(t) = P \{ C_2(t) \leq a \}$
 - ▶ $u_1 =$ renewal density function for $F(\cdot; e_1)$
- ▶ So suffices to show that $\int_0^{\infty} h(t) u_1(t) dt < \infty$

Proof–Continued

- ▶ To show: $\int_0^\infty h(t)u_1(t) dt < \infty$ where $h(t) = P\{C_2(t) \leq a\}$
- ▶ Renewal argument: $h = U_2 * Q$
 - ▶ $U_2 =$ renewal function for $F(\cdot; e_2)$
 - ▶ $Q(t) = F(t + a; e_2) - F(t; e_2)$
- ▶ Heavy-tail key renewal theorem [Erickson70]:
 - ▶ Light-tail KRT: $h(t) \rightarrow m_2(a)/m_2(\infty)$
 - ▶ where $m_2(t) = \int_0^t \bar{F}(u; e_2) du = (1+t)^{1-\beta}/(1-\beta)$
 - ▶ $\bar{F} = 1 - F$
 - ▶ Regular variation: $G \in RV_\lambda$ iff $\lim_{t \rightarrow \infty} G(tx)/G(t) = x^\lambda$
 - ▶ Erickson: $h(t) = O(1/m_2(t))$
 - ▶ Thus $h(t) = O(t^{\beta-1})$



$$U(t) = E[N(t)] = \sum_{n=1}^{\infty} F^n(t)$$
$$u(t) = dU(t)/dt = \sum_{n=1}^{\infty} f^n(t)$$

Proof–Continued

- ▶ To show: $\int_0^\infty h(t)u_1(t) dt < \infty$
where $h(t) = O(t^{\beta-1})$
- ▶ Tauberian theorems: $\bar{F}(\cdot; e_1) \in RV_{-\alpha} \Rightarrow U_1 \in RV_\alpha$
 - ▶ $U_1 =$ renewal function for $F(\cdot; e_1)$
- ▶ Landau's theorem: u_1 ultimately monotone $\Rightarrow u_1 \in RV_{\alpha-1}$
 - ▶ Temporarily assume ultimate monotonicity
 - ▶ Result in Feller $\Rightarrow u_1(t) = O(t^{\alpha-1+\epsilon})$
- ▶ Combine: $\int_0^\infty h(t)u_1(t) dt = \int_0^\infty O(t^{\alpha+\beta+\epsilon-2}) dt < \infty$
since $\alpha + \beta + \epsilon < 1$

Ultimate-Monotonicity Proof (Sketch)

- ▶ $\mathcal{L}(f_1)(s) = \alpha s^\alpha e^s \Gamma(-\alpha, s)$, where

$$\Gamma(a, s) = \int_s^\infty e^{-t} t^{a-1} dt$$

- ▶ $\mathcal{L}(u_1)(s) = \frac{\mathcal{L}(f_1)(s)}{1 - \mathcal{L}(f_1)(s)}$

- ▶ $\mathcal{L}(u_1')(s) = \frac{s\mathcal{L}(u_1)(s)}{1 - \mathcal{L}(u_1)(s)} - u_1(0+) = g(s)$, where

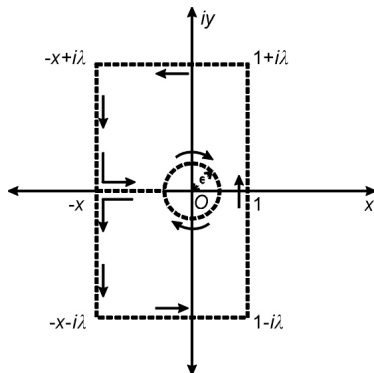
$$g(s) = \frac{\alpha s^{\alpha+1} e^s \Gamma(-\alpha, s)}{1 - \alpha s^\alpha e^s \Gamma(-\alpha, s)} - \alpha$$

- ▶ $g(s)$ is analytic except at origin

Ultimate-Monotonicity Proof — Continued

- ▶ Inversion formula: $u_1'(t) = \frac{1}{2\pi i} \int_{1-i\infty}^{1+i\infty} e^{st} g(s) ds$
- ▶ Apply Cauchy's integral theorem:

$$u_1'(t) \approx -\frac{ct^{\alpha-2} \sin(\pi(1-\alpha))}{\pi} < 0$$



Summary and Conjecture

# HT Clocks	RW Equiv.	Recurrent?	"Positive Recurrent"?
0	–	Yes	Yes
1	2	Yes	Depends
≥ 2	≥ 4	Depends	Depends

A Special Case

- ▶ Hazard rate: $h(x; e) = f(x; e) / \bar{F}(x; e)$
- ▶ Theorem: $P \{ S_n = s \text{ i.o.} \} = 1$ for each s if
 - ▶ State space S is **finite**
 - ▶ GSMP is **irreducible**
 - ▶ $\exists \bar{x} > 0$ s.t. each $F(\cdot; e)$ has **positive density** on $(0, \bar{x})$
 - ▶ At most **one active event** with heavy-tailed clock-setting dist'n
 - ▶ $\alpha(e) \leq h(x; e) \leq \beta(e)$ for each **light-tailed event** e
- ▶ Proof uses regenerative structure [Glynn89] + geometric trials

Questions?

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