

# Discrete-Event Systems and Generalized Semi-Markov Processes

Reading: Section 1.4 in Shedler or Section 4.1 in Haas

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CS 590M: Simulation  
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## Discrete-Event Systems and Generalized Semi-Markov Processes

Discrete-Event Stochastic Systems

The GSMP Model

Simulating GSMPs

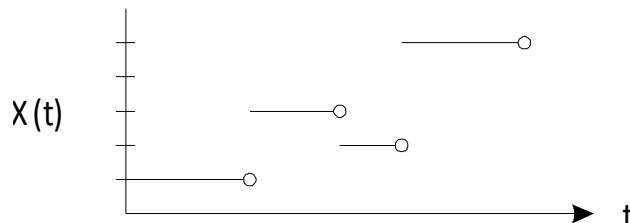
Generating Clock Readings: Inversion Method

Markovian and Semi-Markovian GSMPs

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## Discrete-Event Stochastic Systems

Stochastic state transitions occur at an increasing sequence of random times



How to model underlying process ( $X(t) : t \geq 0$ )?

- ▶ Generalized semi-Markov processes (GSMPs)
- ▶ Basic model of a discrete-event system

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## GSMP Overview

- ▶ **Events** associated with a state “compete” to trigger next state transition
- ▶ Each event has own distribution for determining the next state
- ▶ **New events**
  - ▶ Associated with new state but not old state, *or*
  - ▶ Associated with new state and just triggered state transition
  - ▶ **Clock** is set with time until event occurs (runs down to 0)
- ▶ **Old events**
  - ▶ Associated with old and new states, did not trigger transition
  - ▶ Clock continues to run down
- ▶ **Canceled events**
  - ▶ Associated with old state, but not new state
  - ▶ Clock reading is discarded
- ▶ Clocks can run down at state-dependent **speeds**

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## Clock-Reading Plot

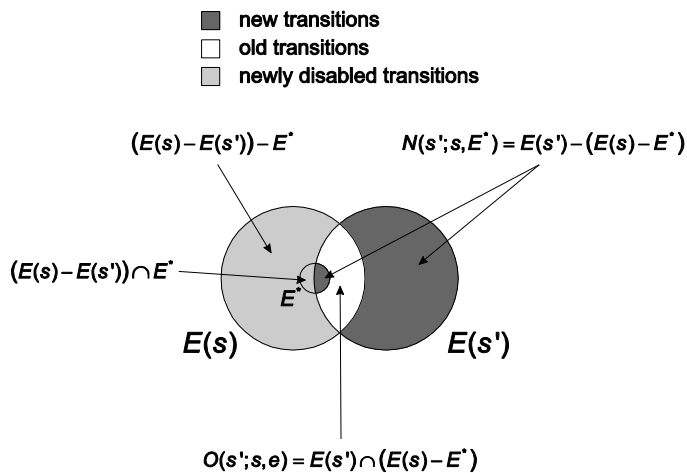
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## GSMP Building Blocks

- ▶  $S$ : a (finite or countably infinite) set of **states**
- ▶  $E = \{e_1, e_2, \dots, e_M\}$ : a finite set of **events**
- ▶  $E(s) \subseteq E$ : the set of **active events** in state  $s \in S$
- ▶  $p(s'; s, E^*)$ : probability that new state =  $s'$  when events in  $E^*$  simultaneously occur in  $s$ 
  - ▶ Write  $p(s'; s, e^*)$  if  $E^* = \{e^*\}$  (unique trigger event)
- ▶  $r(s, e)$ : the nonnegative finite **speed** at which clock for  $e$  runs down in state  $s$ 
  - ▶ Typically  $r(s, e) = 1$
  - ▶ Set  $r(s, e) = 0$  to model "preempt resume" service discipline
- ▶  $F(\cdot; s', e', s, E^*)$ : cdf of new clock-reading for  $e'$  after state transition  $s \xrightarrow{E^*} s'$
- ▶  $\mu$ : initial distribution for state and clock readings
  - ▶ Assume initial state  $s \stackrel{D}{\sim} \nu$  and clock readings  $\stackrel{D}{\sim} F_0(\cdot; e, s)$

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## New and Old Events



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## Example: GI/G/1 Queue

- Assume that interarrival-time dist'n  $F_a$  and service-time dist'n  $F_s$  are continuous (no simult. event occurrences)
- Assume that at time  $t = 0$  a job arrives to an empty system

$X(t) = \#$  of jobs in service or waiting in queue at time  $t$

Can define  $(X(t) : t \geq 0)$  as a GSMP:

- ▶  $S =$
- ▶  $E =$
- ▶  $E(s) =$
- ▶  $p:$
- ▶  $F(x; s', e', s, e^*) :$
- ▶  $r(s, e) =$
- ▶ Initial dist'n:

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## A More Complex Example: Patrolling Repairman

### See handout for details

- ▶ Provides an example of how to concisely express GSMP building blocks

### Specifying a GSMP can be complex and time-consuming, so why do it?

- ▶ Direct guidance for coding (helps catch “corner cases”)
- ▶ Communicates model at high level (vs poring through code)
- ▶ Theory for GSMPs can help in establishing important properties of the simulation
  - ▶ Stability (i.e., convergence to steady state), so that steady-state estimation problems are well defined
  - ▶ Validity of specific simulation output-analysis methods, so that estimates are correct

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## GSMPs and GSSMCs

### GSMP formally defined in terms of GSSMC $((S_n, C_n) : n \geq 0)$

- ▶  $S_n$  = state just after  $n$ th transition
- ▶  $C_n = (C_{n,1}, C_{n,2}, \dots, C_{n,M})$  = clock readings just after  $n$ th transition
- ▶ See Haas or Shedler books for definition of  $P((s, c), A)$  and  $\mu$

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## GSMP Definition

### Define

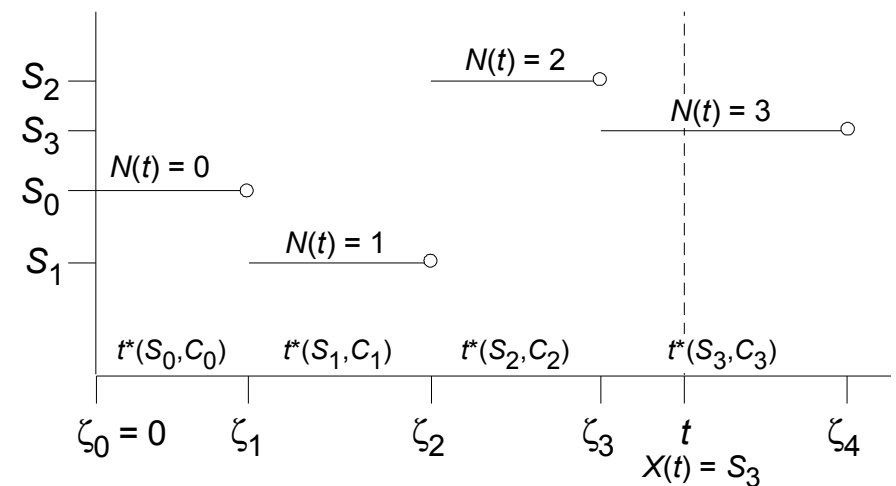
- ▶ Holding time:  $t^*(s, c) = \min_{\{i: e_i \in E(s)\}} c_i / r(s, e_i)$
- ▶  $n$ th state-transition time:  $\zeta_n = \sum_{k=0}^{n-1} t^*(s, c)$
- ▶ # of state transitions in  $[0, t]$ :  $N(t) = \max\{n \geq 0 : \zeta_n \leq t\}$

Let  $\Delta \notin S$  and set

$$X(t) = \begin{cases} S_{N(t)} & \text{if } N(t) < \infty; \\ \Delta & \text{if } N(t) = \infty \end{cases}$$

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## GSMP Definition in a Picture



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## Sample Path Generation

### GSMP Simulation Algorithm (Variable Time-Advance)

1. (Initialization) Select  $s \stackrel{D}{\sim} \nu$ . For each  $e_i \in E(s)$  generate a clock reading  $c_i \stackrel{D}{\sim} F_0(\cdot; e_i, s)$ . Set  $c_i = 0$  for  $e_i \notin E(s)$ .
2. Determine holding time  $t^*(s, c)$  and set of trigger events  $E^* = E^*(s, c) = \{e_i : c_i/r(s, e_i) = t^*(s, c)\}$ .
3. Generate next state  $s' \stackrel{D}{\sim} p(\cdot; s, E^*)$ .
4. For each  $e_i \in N(s'; s, E^*)$ , generate  $c'_i \stackrel{D}{\sim} F(\cdot; s', e_i, s, E^*)$ .
5. For each  $e_i \in O(s'; s, E^*)$ , set  $c'_i = c_i - t^*(s, c) r(s, e_i)$ .
6. For each  $e_i \in (E(s) - E^*) - E(s')$ , set  $c'_i = 0$  (i.e., cancel event  $e_i$ ).
7. Set  $s = s'$  and  $c = c'$ , and go to Step 2.  
(Here  $c = (c_1, c_2, \dots, c_M)$  and similarly for  $c'$ .)

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## Sample Path Generation, Continued

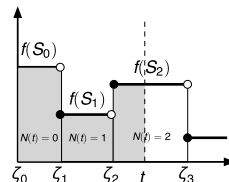
**Algorithm generates sequence of states ( $S_n : n \geq 0$ ), clock-reading vectors ( $C_n : n \geq 0$ ), and holding times ( $t^*(S_n, C_n) : n \geq 0$ )**

**Transition times ( $\zeta_n : n \geq 0$ ) and continuous-time process ( $X(n) : n \geq 0$ ) computed as described previously**

**Use usual techniques to estimate quantities like  $E[f(X(t))]$  or even**

$$\alpha = E \left[ \frac{1}{t} \int_0^t f(X(u)) du \right]$$

$$= E \left[ \frac{1}{t} \left( \sum_{n=0}^{N(t)-1} f(S_n) t^*(S_n, C_n) + f(S_{N(t)}) (t - \zeta_{N(t)}) \right) \right]$$



**Flow charts and diagrams can be helpful (see Law, p. 30–32 for an example)**

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## Generating Clock Readings: Example

**Exponential distribution with rate (intensity)  $\lambda$**

$$f(x; \lambda) = \begin{cases} \lambda e^{-\lambda x} & \text{if } x \geq 0; \\ 0 & \text{if } x < 0 \end{cases} \quad \text{and} \quad F(x; \lambda) = \begin{cases} 1 - e^{-\lambda x} & \text{if } x \geq 0; \\ 0 & \text{if } x < 0 \end{cases}$$

**Mean** =  $1/\lambda$

**Claim:**

If  $U \stackrel{D}{\sim} \text{Uniform}(0, 1)$  and  $V = \frac{-\ln U}{\lambda}$ , then  $V \stackrel{D}{\sim} \exp(\lambda)$

**Proof:**

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## The Inversion Method: Special Case

**Suppose that cdf  $F(x) = P(V \leq x)$  is increasing and continuous**

**Claim:**

If  $U \stackrel{D}{\sim} \text{Uniform}(0, 1)$  and  $V = F^{-1}(U)$ , then  $V \stackrel{D}{\sim} F$

**Proof:**

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## Example: Exponential Distribution

$$F(x) = 1 - e^{-\lambda x}$$

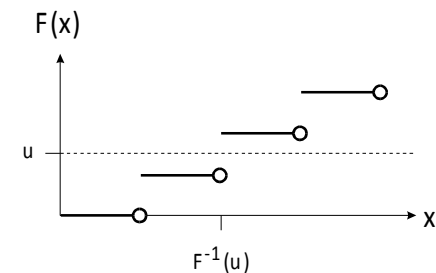
$$F^{-1}(u) =$$

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## The Inversion Method: General Case

**Generalized inverse**

$$F^{-1}(u) = \min\{x : F(x) \geq u\}$$



**Claim still holds:**  $F^{-1}(u) \leq x \Leftrightarrow u \leq F(x)$  by definition

**Exercise:** Show that inversion method = naive method for discrete RVs

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## Markovian GSMPs

### Properties of the Exponential Distribution

If  $X \stackrel{D}{\sim} \exp(\lambda)$  and  $Y \stackrel{D}{\sim} \exp(\mu)$  then

1.  $\min(X, Y) \stackrel{D}{\sim} \exp(\lambda + \mu)$  [indep. of whether  $\min = X$  or  $Y$ ]
2.  $P(X < Y) = \frac{\lambda}{\lambda + \mu}$
3.  $P(X > a + b \mid X > a) = e^{-\lambda b}$  [memoryless property]

**Properties 1 and 2 generalize to multiple exponential RVs**

**Simple** GSMP event  $e'$

$F(\cdot; s', e', s, E^*) \equiv F(\cdot; e')$  and  $F_0(\cdot; e'; s) \equiv F(\cdot; e')$

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## Markovian GSMPs, Continued

**Suppose that all events in a GSMP are simple with exponential clock-setting distn's**

**Key observation:** By memoryless property, whenever GSMP jumps into a state  $s$ , clock readings for events in  $E(s)$  are **mutually independent and exponentially distributed**

Simplified Simulation Algorithm (No clock readings needed)

1. (Initialization) Select  $s \stackrel{D}{\sim} \nu$
2. Generate holding time  $t^* \stackrel{D}{\sim} \exp(\lambda)$ , where  $\lambda = \lambda(s) = \sum_{e_i \in E(s)} \lambda_i$
3. Select  $e_j \in E(s)$  as trigger event with probability  $\lambda_j / \lambda$
4. Generate the next state  $s' \stackrel{D}{\sim} p(\cdot; s, e_j)$
5. Set  $s = s'$  and go to Step 2

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## Markovian GSMPs, Continued

### Structure of a Markovian GSMP

- ▶ Sequence  $(S_n : n \geq 0)$  is a **DTMC** with transition matrix  $R(s, s') = \sum_{e_i \in E(s)} p(s'; s, e_i)(\lambda_i / \lambda)$
- ▶ Given  $(S_n : n \geq 0)$ , holding times are mutually independent with holding time in  $S_n \stackrel{D}{\sim} \exp(\lambda(S_n))$

**Often, occurrence of  $e_i$  in  $s$  causes state to change to a unique state  $y_i = y_i(s)$  with probability 1**

Super-Simplified Simulation Algorithm

1. (Initialization) Select  $s \stackrel{D}{\sim} \nu$
2. Generate holding time  $t^* \stackrel{D}{\sim} \exp(\lambda)$ , where  $\lambda = \sum_{e_i \in E(s)} \lambda_i$
3. Set  $s' = y_i(s)$  with probability  $\lambda_i / \lambda$
4. Set  $s = s'$  and go to Step 2

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## Markovian GSMPs, Continued

A GSMP  $(X(t) : t \geq 0)$  with simple, exponential transitions is a **continuous-time Markov chain (CTMC)** [Ross, Ch. 6]

- ▶ Finite or countable state space
- ▶ Continuous-time Markov property

$$P(X(t+u) = s \mid X(s) : 0 \leq s \leq t) = P(X(t+u) = s \mid X(t))$$

All CTMCs have foregoing structure

- ▶ State sequence is a DTMC
- ▶ Holding times mutually independent and  $\exp(\lambda(s))$  in state  $s$

Q: What can go wrong if events are not simple?

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## Example of Markovian GSMP: Poisson Process

Definition of Poisson process  $(N(t) : t \geq 0)$  with rate  $\lambda$

- ▶  $S = \{0, 1, 2, \dots\}$
- ▶ Single  $\exp(\lambda)$  event
- ▶  $p(s+1; s, e) = 1$

Can show that

$$P(N(t+s) = m+n \mid N(t) = m) = \frac{e^{-\lambda s} (\lambda s)^n}{n!}$$

Examples: # arrivals to a queue, # of lightbulb replacements

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## Semi-Markovian GSMPs

GSMP  $(X(t) : t \geq 0)$  with **simple events** such that  $|E(s)| = 1$  for all  $s \in S$  is a **semi-Markov process**

Definition of semi-Markov process

- ▶ Discrete state space  $S$
- ▶ State sequence  $(X_n : n \geq 0)$  is a DTMC with transition matrix, say,  $R$
- ▶ Holding time in  $s \stackrel{D}{\sim} F(\cdot; s)$
- ▶ "Markov property holds only at state-transition times"

Example: Renewal counting process

- ▶  $S = \{0, 1, 2, \dots\}$
- ▶  $R(s, s+1) = 1$  for all  $s \in S$
- ▶  $F(\cdot; s) \equiv G(\cdot)$  for some  $G$

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