

### Answers to Sample Midterm Questions

1. Claims processing. (This is an example of both a “fork-join” queuing system and a workflow system.)

(a) GSMP specification:

(i) **State space:**  $S =$  the set of all  $(n_1, n_2, m) \in \{0, 1, 2, \dots\}^3$  such that  $n_1 \leq m$  and  $n_2 \leq m$ .

(ii) **Event set mapping:** For  $s = (n_1, n_2, m) \in S$

$e_1 \in E(s)$  for all  $s$

$e_2 \in E(s)$  if and only if  $n_1 > 0$

$e_3 \in E(s)$  if and only if  $n_2 > 0$

$e_4 \in E(s)$  if and only if  $m - n_1 > 0$  and  $m - n_2 > 0$

[equivalently:  $e_4 \in E(s)$  iff  $\min(m - n_1, m - n_2) > 0$ ;  $e_4 \in E(s)$  iff  $m - \max(n_1, n_2) > 0$ ]

(iii) **Speeds:** All speeds  $r(s, e)$  are equal to 1.

(iv) **State-transition probabilities:**

$p(s'; s, e_1) = 1$  when  $s = (n_1, n_2, m)$  and  $s' = (n_1 + 1, n_2 + 1, m + 1)$

$p(s'; s, e_2) = 1$  when  $s = (n_1, n_2, m)$  with  $n_1 > 0$  and  $s' = (n_1 - 1, n_2, m)$

$p(s'; s, e_3) = 1$  when  $s = (n_1, n_2, m)$  with  $n_2 > 0$  and  $s' = (n_1, n_2 - 1, m)$

$p(s'; s, e_4) = 1$  when  $s = (n_1, n_2, m)$  with  $\min(m - n_1, m - n_2) > 0$  and  $s' = (n_1, n_2, m - 1)$

(b) The table is as follows (new clock readings are in **boldface**.)

$n$	$\zeta_n$	$X(\zeta_n)$	$C_{n,1}$	$C_{n,2}$	$C_{n,3}$	$C_{n,4}$	$e^*$
0	0.0	(1,1,1)	<b>4</b>	<b>2</b>	<b>3</b>	--	$e_2$
1	2.0	(0,1,1)	2	--	1	--	$e_3$
2	3.0	(0,0,1)	1	--	--	<b>4</b>	$e_1$
3	4.0	(1,1,2)	<b>5</b>	<b>6</b>	<b>4</b>	3	$e_4$
4	7.0	(1,1,1)	2	3	1	--	$e_3$

(c) The computations are as follows:

$$\hat{\mu} = \frac{1}{5}(1.0 + 3.0 + 7.0 + 5.0 + 4.0) = 4.0$$

$$s^2 = \frac{1}{4}((1.0 - 4.0)^2 + (3.0 - 4.0)^2 + (7.0 - 4.0)^2 + (5.0 - 4.0)^2 + (4.0 - 4.0)^2) = 5.0$$

$$z_p = 1.645 \quad [0.95 \text{ quantile for } N(0,1)]$$

$$\epsilon = 0.1$$

$$n = \frac{z_p^2 s^2}{\epsilon^2 \hat{\mu}^2} = \frac{(1.645)^2 (5)}{(0.1)^2 (4)^2} \approx 85,$$

so we need at least 85 runs.

2. Random number generation.

(a) The inverse of the cdf is given by  $F_X^{-1}(u) = (u/\beta)^{1/\alpha}$  for  $0 \leq u \leq 1$ . Using inversion, generate  $U$  and return  $X = (U/\beta)^{1/\alpha}$ .

(b) The pdf is  $f_X(x) = dF_X(x)/dx = \alpha\beta x^{\alpha-1}$  so the expected value of the distribution is

$\int_0^{\beta^{-1/\alpha}} x f_X(x) dx = \int_0^{\beta^{-1/\alpha}} \alpha\beta x^\alpha dx = \frac{\alpha}{\alpha+1} \beta^{-1/\alpha}$ . For  $\alpha = 2$  and  $\beta = 4$ , the expected value equals  $1/3$ . The given value is less than half of this amount, so the generator is untrustworthy.