CMPSCI 711: More Advanced Algorithms Section 1-5: Sparse Approximations and Algebraic Approximations

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Sparse Recovery

- ► Goal: Find g such that ||f g||₁ is minimized subject to the constraint that g has at most k non-zero entries.
- Define $\operatorname{Err}^{k}(f) = \min_{g: \|g\|_{0} \leq k} \|f g\|_{1}$
- Exercise: $\operatorname{Err}^{k}(f) = \sum_{i \notin S} |f_i|$ where S are indices of k largest f_i
- Using $O(\epsilon^{-1}k \log n)$ space, we can find \tilde{g} such that $\|\tilde{g}\|_0 \leq k$ and

$$\|\tilde{g} - f\|_1 \leq (1 + \epsilon) \operatorname{Err}^k(f)$$

Count-Min Revisited

▶ Consider Count-Min sketch with depth $d = O(\log n)$, width $w = \frac{4k}{\epsilon}$

▶ For
$$i \in [n]$$
, let $\tilde{f}_i = c_{j,h_j(i)}$ for some row $j \in [d]$.

Let S = {i₁,..., i_k} be the indices with maximum frequencies. Let A_i be the event that there doesn't exist k ∈ S \ i, with h_j(i) = h_j(k).
Then for i ∈ [n],

$$\mathbb{P}\left[|f_{i} - \tilde{f}_{i}| \ge \epsilon \frac{\mathsf{Err}^{k}(f)}{k}\right] = \mathbb{P}\left[\neg A_{i}\right] \times \mathbb{P}\left[|f_{i} - \tilde{f}_{i}| \ge \epsilon \frac{\mathsf{Err}^{k}(f)}{k}|\neg A_{i}\right] + \mathbb{P}\left[A_{i}\right] \times \mathbb{P}\left[|f_{i} - \tilde{f}_{i}| \ge \epsilon \frac{\mathsf{Err}^{k}(f)}{k}|A_{i}\right] \\ \le \mathbb{P}\left[\neg A_{i}\right] + \mathbb{P}\left[|f_{i} - \tilde{f}_{i}| \ge \epsilon \frac{\mathsf{Err}^{k}(f)}{k}|A_{i}\right] \\ \le k/w + 1/4 < 1/2$$

• With high probability, all f_i are approximated up to error $\epsilon \operatorname{Err}^k(f)/k$

Sparse Recovery Algorithm

▶ Consider a Count-Min sketch with depth *d* and width w = 4k/ϵ
 ▶ Let f' = (f̃₁, f̃₂,..., f̃_n) be frequency estimates where for all *i*

$$|f_i - \tilde{f}_i| \le \epsilon \frac{\mathsf{Err}^k(f)}{k}$$

• Let \tilde{g} be f' with all but the kth largest entries replaced by 0.

• *Lemma*: $\|\tilde{g} - f\|_1 \le (1 + 3\epsilon) \operatorname{Err}^k(f)$

Proof of Lemma

- ▶ Let $S, T \subset [n]$ be indices corresponding to largest values of f_i and \tilde{f}_i .
- For a vector x ∈ ℝⁿ and I ⊂ [n], write x_I as the vector formed by zeroing out all entries of x except for those indices in I.

► Then.

$$\begin{split} \|f - f'_{T}\|_{1} &\leq \|f - f_{T}\|_{1} + \|f_{T} - f'_{T}\|_{1} \\ &= \|f\|_{1} - \|f_{T}\|_{1} + \|f_{T} - f'_{T}\|_{1} \\ &= \|f\|_{1} - \|f'_{T}\|_{1} + (\|f'_{T}\|_{1} - \|f_{T}\|_{1}) + \|f_{T} - f'_{T}\|_{1} \\ &\leq \|f\|_{1} - \|f'_{T}\|_{1} + 2\|f_{T} - f'_{T}\|_{1} \\ &\leq \|f\|_{1} - \|f'_{S}\|_{1} + 2\|f_{T} - f'_{T}\|_{1} \\ &\leq \|f\|_{1} - \|f_{S}\|_{1} + (\|f_{S}\|_{1} - \|f'_{S}\|_{1}) + 2\|f_{T} - f'_{T}\|_{1} \\ &\leq \|f - f_{S}\|_{1} + \|f_{S} - f'_{S}\|_{1} + 2\|f_{T} - f'_{T}\|_{1} \\ &\leq \|f - f_{S}\|_{1} + \|f_{S} - f'_{S}\|_{1} + 2\|f_{T} - f'_{T}\|_{1} \\ &\leq \text{Err}^{k}(f) + k\epsilon \text{Err}^{k}(f)/k + 2k\epsilon \text{Err}^{k}(f)/k \\ &= (1 + 3\epsilon) \text{Err}^{k}(f) \end{split}$$

Similar Result for ℓ_2

► Goal: Find g such that ||f - g||₂ is minimized subject to the constraint that g has at most k non-zero entries.

• Define
$$\operatorname{Err}_2^k(f) = \min_{g: \|g\|_0 \le k} \|f - g\|_2^2$$

▶ Using $O(\epsilon^{-2}k \log n)$ space, we can find \tilde{g} such that $\|\tilde{g}\|_0 \leq k$ and

$$\|\tilde{g} - f\|_2^2 \leq (1 + \epsilon) \operatorname{Err}_2^k(f)$$

Outline

Wavelets

Haar Wavelets

• *Defn:* For n = 8, Haar Wavelet basis consists of rows of the matrix.

$$M = \begin{pmatrix} \frac{1}{\sqrt{8}} & \frac{1}{\sqrt{8}}$$

and for $n = 2^r$, the construction generalizes in the natural way.

▶ Note that the basis is orthonormal and $MM^T = M^TM = I$. Hence, any signal $f \in \mathbb{R}^n$ can be expressed in the Haar basis.

Sparse Representation in Haar Basis

• Let $\mathcal{H} = \{\phi_1, \dots, \phi_n\}$ be the Haar basis.

► Goal: Find g that minimizes ||f - g||₂ subject to the constraint that g can be expressed as the sum of at most k Haar basis vectors, i.e., g is k-sparse in the Haar basis.

• Write
$$f = \sum_i \lambda_i \phi_i$$
 where $\lambda_i = \phi_i \cdot f$.

• Suppose
$$g = \sum_{i \in I} \mu_i \phi_i$$
 for some $I \subset [n]$ of size at most k .

Then

$$\|f-g\|_2^2 = \sum_{i \in I} (\lambda_i - \mu_i)^2 + \sum_{i \notin I} \lambda_i^2$$

• Hence, we want to find k values of i that maximize $\mu_i = \phi_i \cdot f$.

Time Series Model

- ► Suppose coordinations of f are presented in order (f₁, f₂,..., f_n). This is called the *time-series model*.
- Can easily compute $\mu_i = \phi_i \cdot f$
- At any given time,
 - We've calculated μ_i exactly for some $i \in A$
 - We've calculated μ_i partially for some $i \in B$
 - We haven't started computing μ_i for $i \notin A \cup B$
- Lemma: The size of B is at most $\log_2 n$.
- Algorithm: Maintain only the k largest values of μ_i for $i \in A$.
- We find the optimal k term representation in $O(k + \log n)$ space.

General Update Model

- Can express goal in terms of standard basis...
- Because *M* is unitary,

$$\|f - g\|_2^2 = (f - g)^T (f - g) = (f - g)^T M^T M (f - g) = \|Mf - Mg\|_2^2$$

and g is k-sparse in Haar basis iff Mg is k-sparse in standard basis.

- ► Hence, finding best g is same as finding h = Mg with ||h||₀ ≤ k that minimizes ||Mf h||₂
- ▶ Using Count-Sketch algorithm, can find \tilde{h} with $\|\tilde{h}\|_0 \leq k$ such that

$$\begin{split} Mf - \tilde{h}\|_{2}^{2} &\leq (1+\epsilon) \min_{\substack{h:h \text{ is } k-\text{sparse in standard basis}}} \|Mf - h\|_{2}^{2} \\ &= (1+\epsilon) \min_{g:g \text{ is } k-\text{sparse in Haar basis}} \|Mf - Mg\|_{2}^{2} \\ &= (1+\epsilon) \min_{g:g \text{ is } k-\text{sparse in Haar basis}} \|f - g\|_{2}^{2} \end{split}$$