

## COMPSCI 501 Spring 2019 - Homework 5

Released 4/1/2019

Due in Gradescope on Fri Apr 12, 11:59pm

### Instructions.

You may work in groups, but you must write solutions yourself. List your collaborators.

All your answers should contain a rigorous argument of correctness.

Please submit a PDF file. You may submit a scanned handwritten document, but a typed submission is preferred. Please assign pages to questions in Gradescope.

**Q1** (5p). Let *SPARSE-SAT* be the problem: given a formula in CNF form and an integer  $k$ , find a satisfying assignment in which at most  $k$  variables are TRUE, if one exists. Prove that *SPARSE-SAT* is NP-complete.

**Q2** (10p). Consider the problem of determining, given an undirected graph  $G = (V, E)$ , whether  $G$  contains a clique of size  $\lfloor V/2 \rfloor$ . Is this problem in NP? Is it NP-complete? Prove your statements.

**Q3** (15p). Consider the following problem, called *THIRD-FRACTION*: given a set  $\{a_1, \dots, a_n\}$  of positive integers, decide whether we can find indices  $i_1, \dots, i_k$  such that  $\sum_j a_{i_j} = \frac{1}{3} \sum_j a_j$ . Show that *THIRD-FRACTION* is NP-complete.

**Q4** (5p). For  $k \geq 1$ , let  $k$ -*COLOR* be

$\{G \mid G \text{ is an undirected graph whose vertices can be properly colored with at most } k \text{ colors}\}.$

For which values of  $k$  is the problem  $k$ -*COLOR* in P? For which values is it NP-complete?

Prove your statements.

**Q5** (10p). Are all NP-hard problems necessarily PSPACE-complete? What about the converse?

**Q6** (15p). Suppose  $T$  is a set of “tile types”,  $T = \{t_1, t_2, \dots, t_k\}$ , and  $H$  and  $V$  be two relations  $H, V \subseteq (T \times T)$ , with  $H(t_i, t_j)$  meaning  $t_j$  can be placed immediately to the right of  $t_i$ , and  $V(t_i, t_j)$  meaning  $t_j$  can be placed immediately below  $t_i$ . Given an integer  $n$ , let  $R$  be a length- $n$  horizontal row of tiles (possibly with repeats) from  $T$  satisfying all  $H$ -constraints. Consider the language  $EXTEND-TILING = \{\langle n, T, H, V, R \rangle \mid \text{there exists a tiling of the } n \times n \text{ grid satisfying all } H \text{ and } V \text{ constraints, with } R \text{ as its topmost row}\}$ . Show that *EXTEND-TILING* is NP-complete.