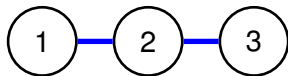
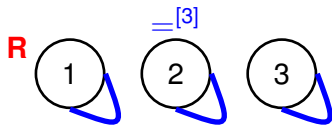


CS250: Discrete Math for Computer Science

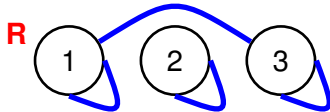
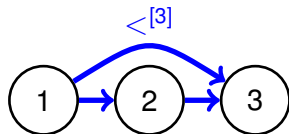
L26: Equivalence Relations

Equivalence Relations

Reflexive $\equiv \forall x E(x, x)$



succ/pred

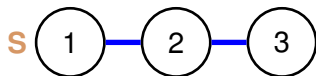
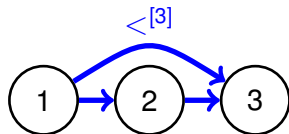
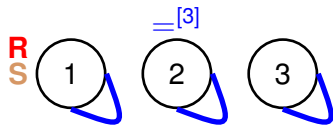


$\equiv (\text{mod } 2)$

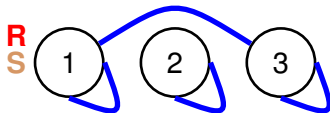
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succ/pred



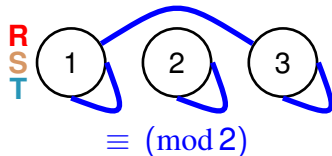
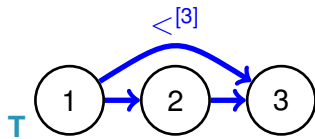
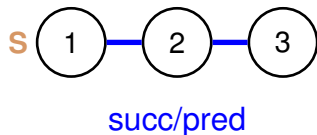
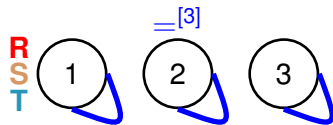
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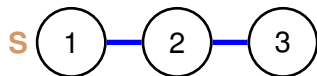
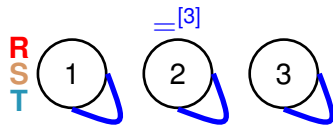


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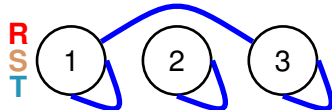
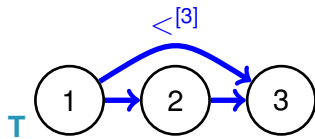
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succ/pred



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Def. An **equivalence relation** is a relation that is **reflexive**, **symmetric** and **transitive**.

Equivalence Relations on 3-Element Sets

Reflexive $\equiv \forall x E(x, x)$

Symmetric $\equiv \forall xy (E(x, y) \rightarrow E(y, x))$

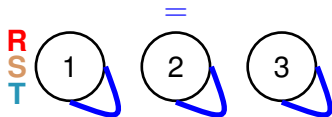
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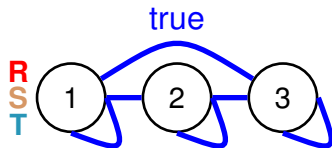
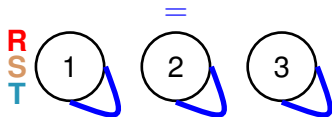


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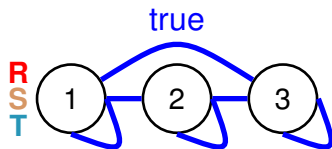
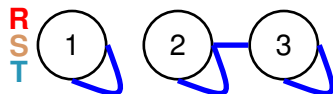
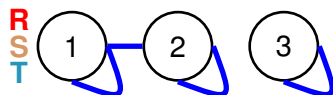
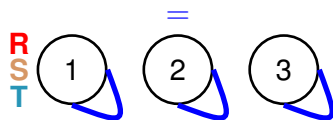


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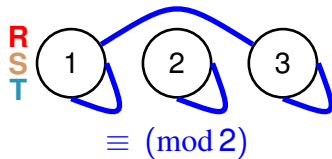
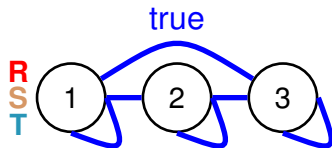
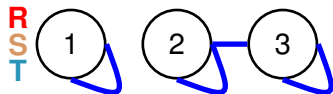
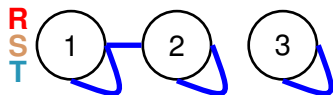
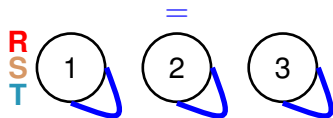


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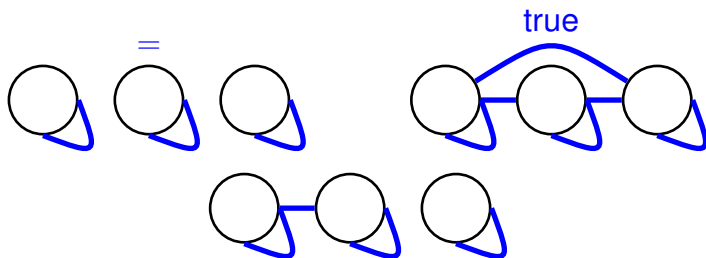


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Partitions

Def. A **partition** of a non-empty set V is collection of pairwise disjoint, non-empty subsets, (P_1, P_2, \dots) of V whose union is V :

$$\emptyset \neq P_i \subseteq V$$

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Prop. There are exactly five partitions of $V = \{1, 2, 3\}$:

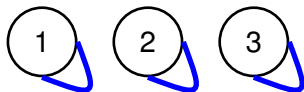
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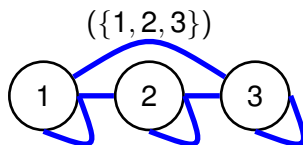
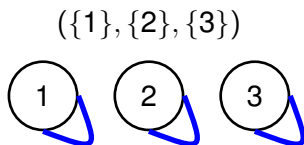


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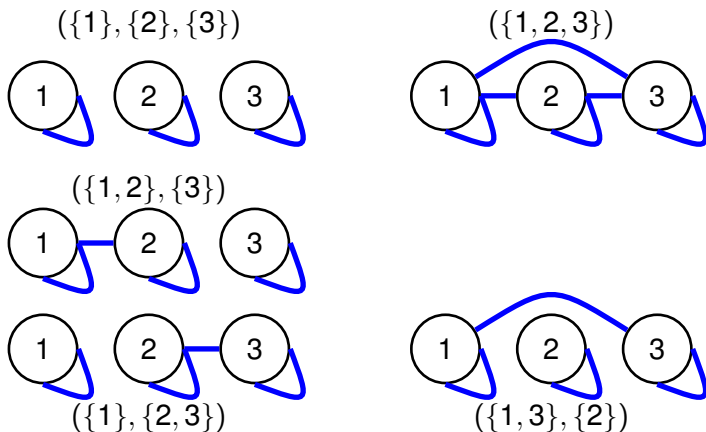


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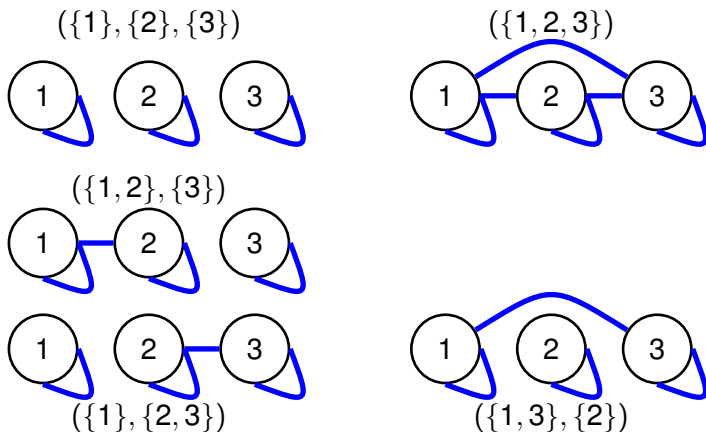


Partitions relate to Equivalence Relations, How?

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2. $a \cdot b \equiv a' \cdot b' \pmod{m}$
3. $a^n \equiv (a')^n \pmod{m}$

Recall L17: $a^{-1} \bmod m$ exists iff $\gcd(a, m) = 1$

$\cdot \mathbb{Z}/6\mathbb{Z}$	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	0	2	4
3	0	3	0	3	0	3
4	0	4	2	0	4	2
5	0	5	4	3	2	1

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$$|\mathbf{Z}_6^*| = \{1, 5\}$$

$\cdot \mathbb{Z}_6^*$	1	5
1	1	5
5	5	1

\mathbf{Z}_m^* is the multiplicative group mod m

$$|\mathbf{Z}_5^*| = \{1, 2, 3, 4\}$$

$\cdot \mathbf{Z}_5^*$	1	2	3	4
1	1	2	3	4
2	2	4	1	3
3	3	1	4	2
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$$\Sigma_{\text{group}} \stackrel{\text{def}}{=} (; 1, \cdot [\text{infix}]^2, {}^{-1} [\text{postfix}]^1)$$

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Prop. For all $m > 1$, Z_m^* is a group.

Euler's phi function, φ

Def. For $m > 1$,

$$\varphi(m) \stackrel{\text{def}}{=} \|\mathbf{Z}_m^*\| = |\{a \in \mathbf{Z}/m\mathbf{Z} \mid \gcd(a, m) = 1\}|$$

m	$\varphi(m)$	$ \mathbf{Z}_m^* $
2	1	{1}
3	2	{1, 2}
4	2	{1, 3}
5	4	{1, 2, 3, 4}
6	2	{1, 5}
7	6	{1, 2, 3, 4, 5, 6}
8	4	{1, 3, 5, 7}
9	6	{1, 2, 4, 5, 7, 8}
10	4	{1, 3, 7, 9}
11	10	{1, 2, 3, 4, 5, 6, 7, 8, 9, 10}
12	4	{1, 5, 7, 11}

Fermat's Little Theorem

Thm: For p prime, $a \in \mathbf{Z}_p^*$, $a^{p-1} \equiv 1 \pmod{p}$

Proof:

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$$1 \equiv a^{p-1} \pmod{p} \quad \square$$