

Parallel & Concurrent Programming: ZPL

Emery Berger
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Outline

- Previously:
 - **MPI** – point-to-point & collective
 - Complicated, far from problem abstraction
 - **OpenMP** - parallel directives
 - Language extensions to Fortran/C/C++
 - Questionable semantics, error-prone
- Today:
 - Something **way** better: ZPL



lecture material from ZPL project, UW

ZPL

- **Parallel array language**
 - Implicitly parallel
 - No parallel constructs *per se*
 - Very high level
 - Assignments work at array level, as in
 $A := B + C$
 - Machine independent
 - Compiles to ANSI C +
communication library calls (e.g., MPI)
 - Efficient



Comparison

- Matrix-multiplication:

- C

- triply-nested loop

```
for (i=0;i<n;i++){
  for (j=0;j<n;j++){
    for (k=0;k<n;k++){
      c[i][j]=c[i][j]
        +a[i][k]*b[k][j];
    }
  }
}
```

- ZPL

```
[1..n,1..n] for k := 1 to n do
  C += (>>[,k] A) * (>>[k,] B);
end;
```

- dot-product of rows & columns
- efficiently implemented on parallel machines



ZPL Outline

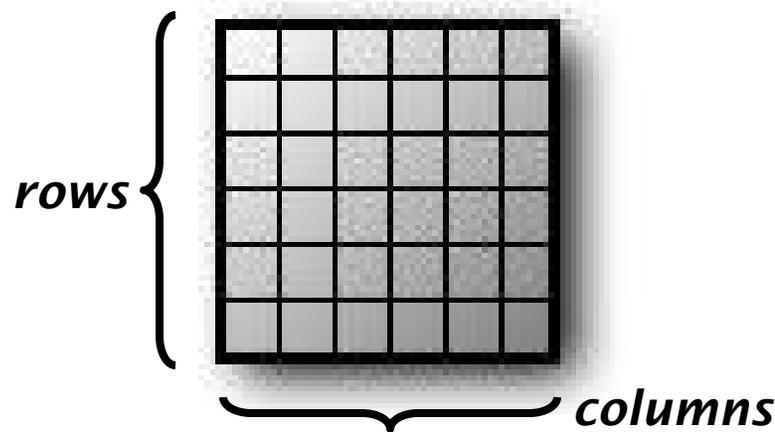
- Language overview
 - Regions
 - Directions
 - Parallel array operations
 - Handling boundary conditions
- ZPL programs & performance



Regions

- Key abstraction in ZPL: **regions**
 - **Index sets** (*rows,cols*) partition matrices
 - Operate on regions, not indexed items!

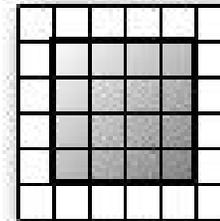
```
region  
R = [1..n,1..n];
```



Region Examples

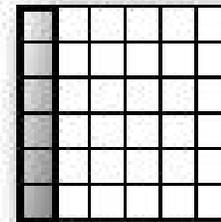
- Interior of matrix

```
region  
  InTR =  
    [2..n-1,2..n-1];
```



- Left-most column

```
region  
  Left = [1..n,1];
```



Directions

- **Directions:**
 - Offset vectors used to manipulate regions & array data

```
direction
north = [-1, 0];
east  = [ 0, 1];
south = [ 1, 0];
west  = [ 0,-1];
```

```
nw = [-1, -1];
ne = [-1,  1];
sw = [ 1, -1];
se = [ 1,  1];
```



Creating New Regions

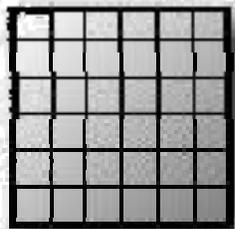
- **Prepositions** create new regions:
 - **in**
 - Applies direction to select part of region
 - **of**
 - Creates new region outside existing region
 - **at**
 - Shifts a region by a direction
 - **by**
 - Creates new region **strided** by direction



Applying Directions

- Use "in" to apply direction to region

```
region  
R = [1..n,1..d];
```



+

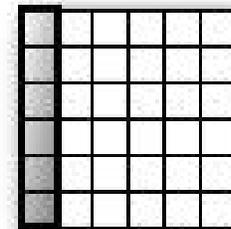
```
direction
```

```
west = [ 0, -1];
```



=

```
region  
Left = west in R;
```



Create Region Outside

- Use “of” to create region **outside** existing region
 - Extends region

```
region  
IntrR =  
[2..n-1,2..n-1];
```

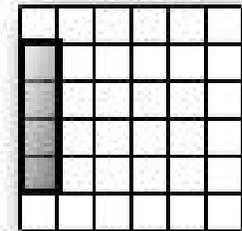


```
direction
```

```
west = [ 0, -1];
```



```
region  
SmallLeft =  
west of IntrR;
```



Shifting Regions

- Use “at” to create new region shifted by a direction

```
region  
Intr =  
[2..n-1,2..n-1];
```

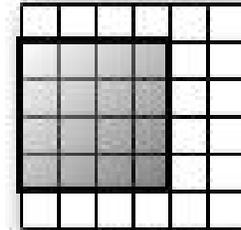


```
direction
```

```
west = [ 0, -1];
```



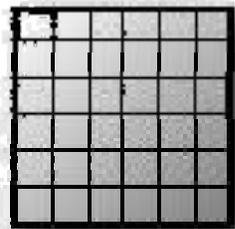
```
region  
IntrLeft =  
Intr at west;
```



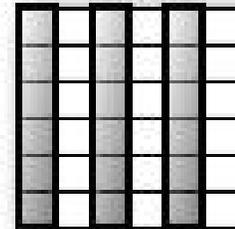
Striding Regions

- Use “by” to create new region strided by a direction

```
region  
R = [1..n.1..m];
```



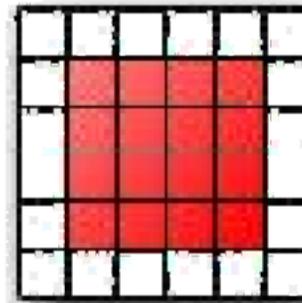
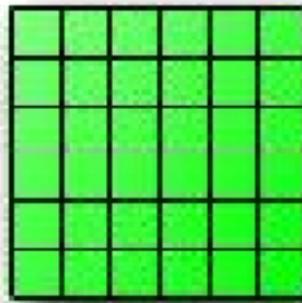
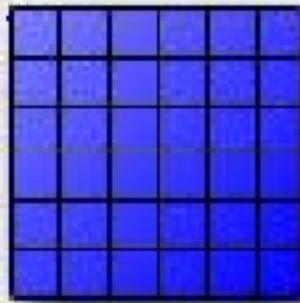
```
direction  
step = [1,2];  
region  
SR = R by step;
```



Parallel Arrays

- Parallel arrays declared over regions

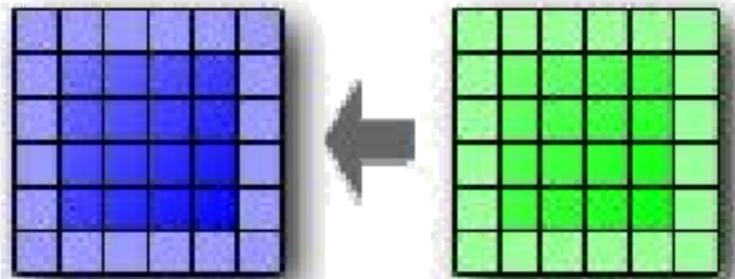
```
var  
  A, B : [R] double;  
  C : [IntR] double;
```



Computing Over Arrays

- Can use regions as **modifiers** that define computations over arrays:

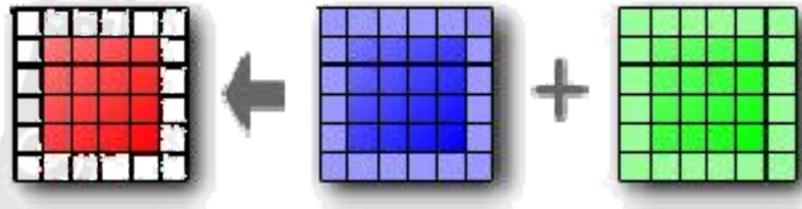
`[IntR] A := B;`



Arrays & Communication

- Most computations in ZPL do not involve communication

[IntR] C := A + B;



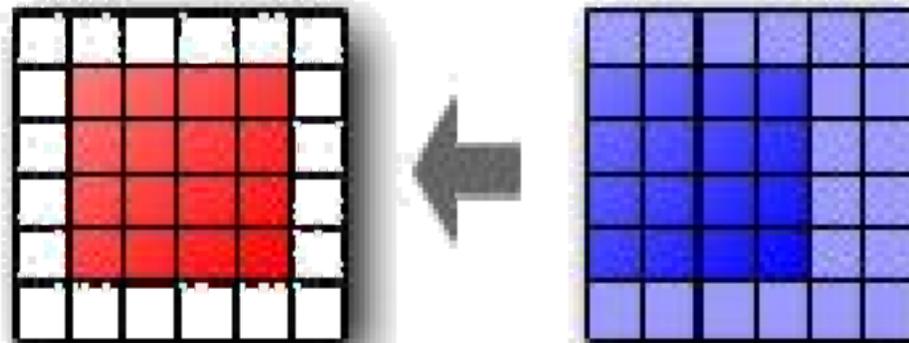
- Exceptions include:
 - Shifting
 - Reduction
 - Broadcast
 - All-to-all



Shifting Arrays

- @ operator shifts data in direction
 - This translation induces point-to-point communication

```
[IntR] C := A@west;
```

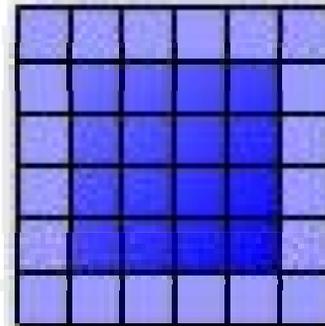


Reduction

- **Op<<** computes reductions
 - Reduction (tree-style) communication
 - +<< (sum), *<< (times), min<<, max<<...
 - For prefix (scan), use **op||**

```
[IntR] sum := +<< A;
```

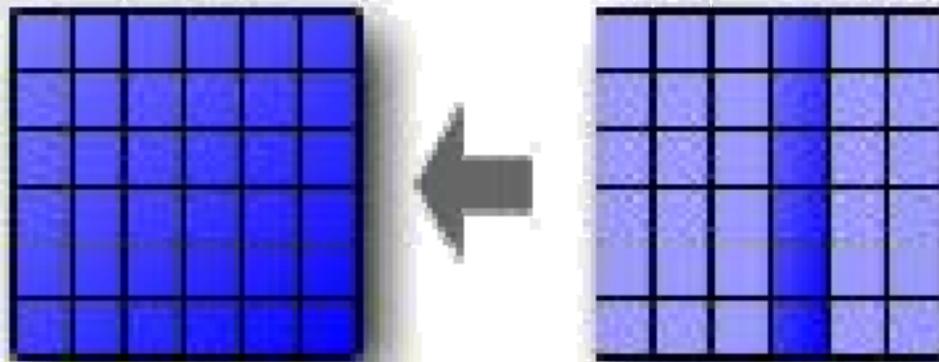
sum



Broadcast (Flooding)

- \gg (**flood**) replicates data across dimensions of array
 - Triggers broadcast operation

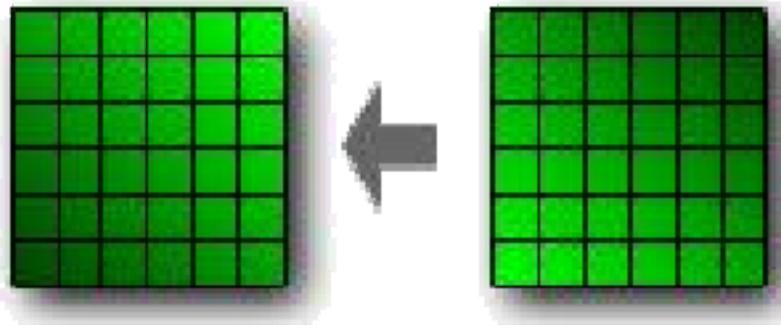
```
[E] A :=  $\gg$ [1..n,i] A;
```



Mapping

- **Remap (#)** moves data between arrays
 - Specified by “map” arrays
 - Built-in **Index1**, **Index2**
 - **Index1** = row indices, **Index2** = col indices

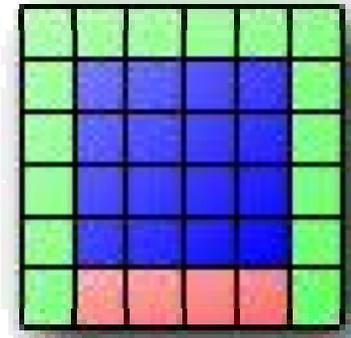
```
[R] B := B#[Index2, Index1];
```



Boundary Conditions

- Boundary conditions (“corner cases”)
 - Usually tedious, error-prone
 - **Very simple in ZPL**

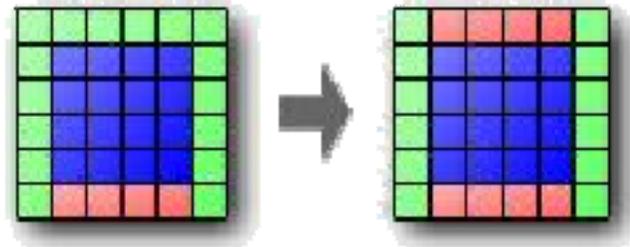
```
[north in R] A == 0.0;  
[east in R] A == 0.0;  
[west in R] A == 0.0;  
[south of IntR] A == 1.0;
```



Boundary Conditions

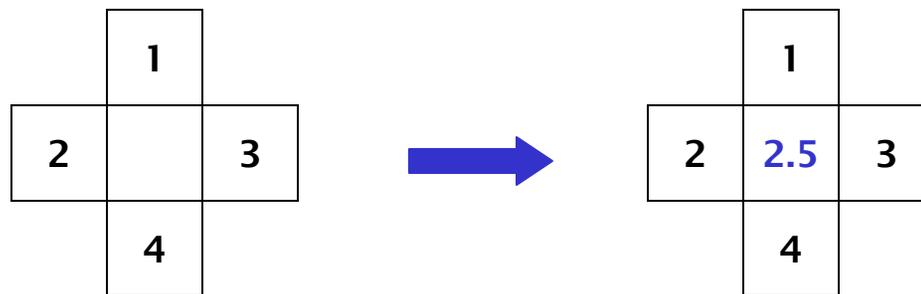
- **Periodic boundary conditions with wrap**

[north of IntR] wrap A;



ZPL Example

- **Jacobi iteration** – replace elements in array with average of four nearest neighbors, until largest change $< \delta$
 - Consider difficulty of parallelizing with MPI/OpenMP
 - boundary conditions, etc.



ZPL Example

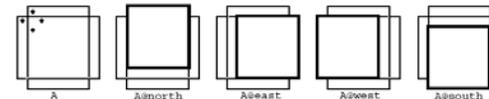
```
program Jacobi;
/*          Jacobi Iteration
   Written by L. Snyder, May 1994      */
config var  n      : integer  = 512; -- Declarations
            delta : float    = 0.000001;

region      R = [1..n, 1..n];
var         A, Temp: [R] float;
            err   : float;

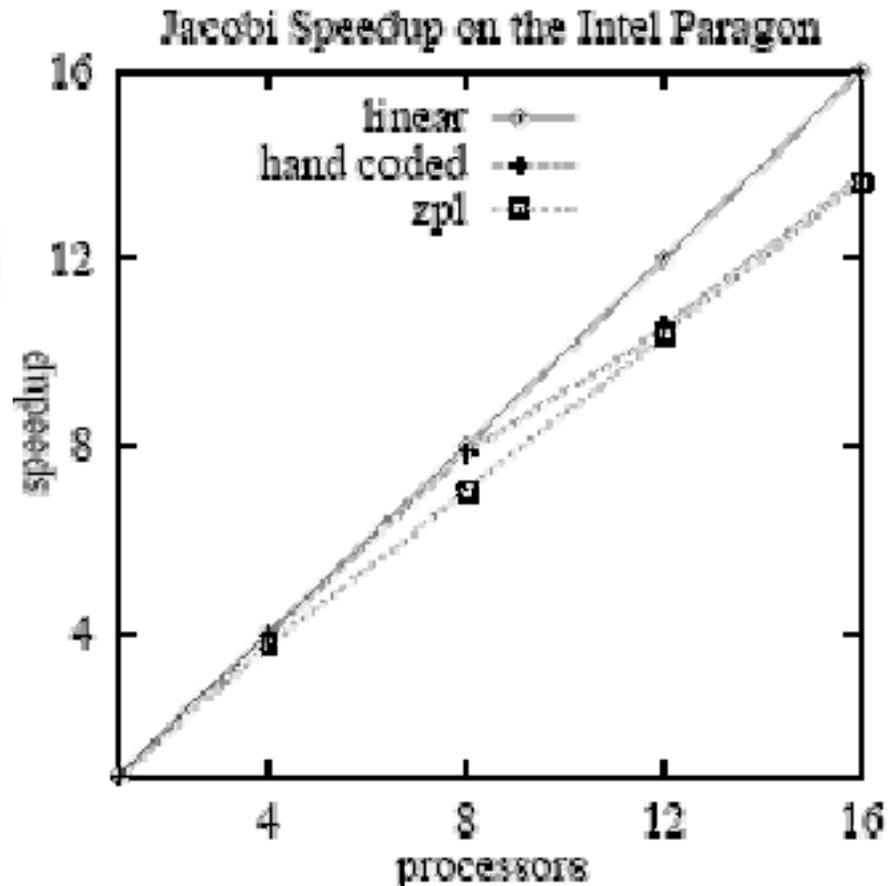
direction   north = [-1, 0];
            east  = [ 0, 1];
            west  = [ 0,-1];
            south = [ 1, 0];

procedure Jacobi();
begin
  [R]          A := 0.0;           -- Initialization
  [north of R] A := 0.0;
  [east  of R] A := 0.0;
  [west  of R] A := 0.0;
  [south of R] A := 1.0;

  [R]          repeat           -- Body
                Temp := (A@north + A@east
                        + A@west + A@south)/4.0;
                err := max<< abs(A - Temp);
                A := Temp;
  until err < delta;
end;
```



ZPL Performance



ZPL Example: Life

Conway's Game of Life simulates cells which can live, die, and reproduce according to the following rules:

Rule I. (Survival) A cell survives only if it has two or three live neighbors.



Rule II. (Birth) A cell is born in any empty square with exactly three live neighbors.



ZPL Example: Life

```

program Life;
config var
  n : integer = 100;
region
  High = {0..n+1, 0..n+1};
  R    = {1..n, 1..n};
direction
  nw = {-1, -1}; north = [-1, 0]; ne = [-1, 1];
  west = [ 0, -1]; east = [ 0, 1];
  sw = [ 1, -1]; south = [ 1, 0]; se = [ 1, 1];
var
  TW : [High] boolean; == The World
  NN : [R] integer; == Number of Neighbors
procedure Life();
begin
  == initialize the world
  [R] repeat
    NN := TW@nw + TW@north + TW@ne +
          TW@west + TW@east +
          TW@sw + TW@south + TW@se;
    TW := (TW & NN = 2) | (NN = 3);
  until !(<< NN);
end;

```

A configuration variable can be set on the command line.

Count the live neighbors.

Update the world.

Is this a bleak metaphor or what?

Conway's Game of Life simulates cells which can live, die, and reproduce according to the following rules:

Rule 1. (Survival) A cell survives only if it has two or three live neighbors.



Rule 2. (Birth) A cell is born in any empty square with exactly three live neighbors.



The End

- Next time:
 - Your turn!
 - Occam & Multilisp

