

# COMPSCI 514: Algorithms for Data Science

---

Cameron Musco

University of Massachusetts Amherst. Spring 2026.

Lecture 19

# Logistics

- Midterm 2 is this Thursday in class, 1-2:15pm.
- Closed book – no cheatsheets, calculators, or other aids allowed.
- Study guide is posted on the course webpage (under the Midterm 2 row in the schedule tab).
- Past finals, which cover mostly the same material as this midterm, are posted in Canvas.
- Ignore content on optimization, which we have not covered yet. *convexity, gradient, Lipschitz*
- Solutions to study guide questions also posted in Canvas.
- I will hold regular office hours today after class, focused on midterm review.

## Suggested Studying Approach:

- Review the study guide to get a sense of what you need to know, and then mostly focus on doing practice questions from the past finals and the study guide.
- Review slides as needed.
- Do some practice exams, under time constraints, with no material in front of you.
- Note that some of the past exams were 2 hours long, but designed to be completed in 90 minutes. This exam will be a bit shorter.

# Midterm Format

## Rough Outline: (subject to changes)

- Question 1: 4-5 True/False questions. No justification needed.
- Question 2: 4-5 numerical answers, like quiz questions. No justification needed.
- Question 3: 4-5 part question on analyzing an algorithm. Similar in style to but easier than a homework question.
- Question 4: More challenging 4-5 part question on analyzing an algorithm – more similar to a homework question.
- Potentially some extra credit subquestions on Q3/Q4.

Content or Format Questions?

# Questions

# Questions

# Questions

# Spectral Graph Theory

# Quiz Review

2 Multiple Choice 1 point

Let graph  $G$  consist of a clique on 10 nodes along with a cycle on 10 different nodes. There are no other edges. Let  $L \in \mathbb{R}^{20 \times 20}$  be the Laplacian matrix of  $G$ . What is the minimum eigenvalue of  $L$ ?

- 0  
 10  
 20  
 -10  
 -100

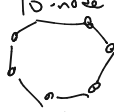
$$v^T L v \geq 0 \quad \forall v$$

$$= \lambda \quad \text{wh}$$

$$v \text{ is an eigenvector}$$



$\lambda_n(L) = 0$



$$v^T L v = \sum_{(i,j) \in E} [v(i) - v(j)]^2$$

$Lv = (D - A)v$

$= Dv - Av$

$$= \begin{bmatrix} d_1 \\ d_2 \\ \vdots \\ d_n \end{bmatrix} - \begin{bmatrix} a_1 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

$\lambda_{n-1}(L)$

$v_{n-1}^T \mathbf{1} = 0$

$\sum_{i=1}^{20} v_{n-1}(i) = 0$

$$v_{n-1}^T L v_{n-1} = \sum_{(i,j) \in E} (v(i) - v(j))^2 = 0$$

$$v_{n-1} = \left. \begin{bmatrix} \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \end{bmatrix} \right\} 10 \text{ clique nodes}$$

$$\left. \begin{bmatrix} \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \\ \vdots \end{bmatrix} \right\} 10 \text{ cycle nodes}$$

$$v^T L v = v^T (\lambda v)$$

$$= \lambda \|v\|_2^2$$

$$= \lambda \quad \text{if } \|v\|_2^2 = 1$$

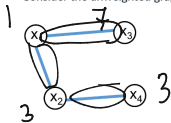
$\lambda_{n-1}(\text{clique}) > \lambda_{n-1}(\text{cycle})$

# Quiz Review

3 Formula 1 point



Consider the unweighted graph  $G$  shown below and let  $L$  be its graph Laplacian. Let  $\vec{x} = [1, 3, 7, 3]$ . What is  $\vec{x}^T L \vec{x}$ ?



$$(7-1)^2 + (3-1)^2 + (5-3)^2 = \underline{\underline{40}}$$

Hint: You don't need to explicitly write down  $L$  or do any linear algebra.

Answer

$$\begin{aligned} v_{n-1} &= \text{argmin}_T v^T L v \\ &\text{s.t. } v_{n-1} \neq 0 \\ &\|v\|_2 = 1 \\ c &= \text{argmin}_T c^T L c \\ &\text{s.t. } c^T \mathbf{1} = 0 \\ c &\in \{-1, 1\} \end{aligned}$$

# Quiz Review

4

Multiple Choice 1 point

$$\|z\|_2 = 1 \quad z^T \mathbf{1} = 0$$



Consider solving the optimization problem:  $\min_{z \in \{-1, 1\}^n}$ : not all entries of  $z$  are equal  $z^T L z$ .

What is this optimization problem commonly known as?

- Computing the lowest eigenvalue of the graph Laplacian.
- Computing the second lowest eigenvalue of the graph Laplacian.
- Computing the minimum cut.
- Computing the smallest node degree.
- Computing the maximum eigenvalue of the graph Laplacian.

$$z \in \{-1, 1\}^n$$

$$z^T L z = 4 \cdot \text{cut}$$

$$z^T \mathbf{1} = 0$$

↳ balance

# Quiz Review

5

Multiple Answer 1 point



Which of the following vectors can be used to recover the two communities in the stochastic block model? Check all that apply.

- Highest eigenvector of the adjacency matrix  $A$ .
- Second highest eigenvector of the adjacency matrix  $A$ .
- Second highest eigenvector of the Laplacian matrix  $L$ .
- Lowest eigenvector of the Laplacian matrix  $L$ .
- Second lowest eigenvector of the Laplacian matrix  $L$ .

# More Spectral Graph Theory Review

$$L = D - A = (n-1)I - A = (n-1)I - (J - I) = nI - J$$

3. Let  $G$  be a fully connected graph (a complete graph), with self-loops. What are the eigenvalues of its corresponding adjacency matrix  $A$  and Laplacian  $L$ ? What if there are no self-loops?

$$A = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{pmatrix} = \begin{pmatrix} n \\ n \\ n \\ n \end{pmatrix}$$

$$\lambda_1 = n, \quad \vec{v}_1 = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

$$\lambda_2 = 0$$

$$\vdots = 0$$

$$\lambda_n = 0$$

$v_2, \dots, v_n$  } any basis for the null space

$$L = D - A = nI - A$$

$$\text{tr}(A) = \sum_{i=1}^n A_{ii} = \sum_{i=1}^n \lambda_i$$

$$= n = n + ? + ? + ?$$

$$A = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$$\lambda_1 = n-1$$

$$\lambda_2 = -1$$

$$\lambda_n = -1$$

$$\begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} = (J - I)v = Av = Jv - v = \lambda v - v = (\lambda - 1)v$$

null( $A$ ) has dimension  $n - \text{rank}(A) = n - 1$   
 $\hookrightarrow \{x : Ax = 0\}$

# More Spectral Graph Theory Review

1. Consider a graph  $G$  with Laplacian matrix  $\mathbf{L}$ . Consider the problem:  $x_* = \arg \min_{x: \|x\|=1} x^T \mathbf{L} x$ .  
What is  $x_*$ ? What value of  $x_*^T \mathbf{L} x_*$  does it achieve?

# General Linear Algebra

# Views of Matrix Multiplication

$$\left[ A \begin{matrix} v^{(1)} \\ \vdots \\ v^{(n)} \end{matrix} \right] \quad (\text{column combination})$$

$$A_1 v^{(1)} + A_2 v^{(2)} + \dots + A_n v^{(n)}$$

(row inner products)

$$\begin{bmatrix} a_1^T \\ \vdots \\ a_n^T \end{bmatrix} \begin{bmatrix} v^{(1)} \\ \vdots \\ v^{(n)} \end{bmatrix} = \begin{bmatrix} a_1^T v \\ a_2^T v \\ \vdots \\ a_n^T v \end{bmatrix}$$

$$\left[ A \right] \begin{bmatrix} d_1 \\ \vdots \\ d_n \end{bmatrix} = \begin{bmatrix} | & | \\ d_1 A_1 & d_2 A_2 \\ | & | \end{bmatrix}$$

$$\left[ A \right] \begin{bmatrix} | \\ \vdots \\ | \end{bmatrix} = \begin{bmatrix} a_1^T | \\ \vdots \\ a_n^T | \end{bmatrix}$$

# Views of Matrix Multiplication

## Row Span/Column Span

$$n \times d \quad n \times p \quad p \times d$$

Let  $X = AB$ . Fill in the blank:

$X$ 's columns are spanned by ...  $A$ 's columns

$$X_i = [A] [b_i]$$

$$\text{col}(X) \subseteq \text{col}(A) \quad \text{rank}(X) \leq \text{rank}(A)$$

$X$ 's rows are spanned by ... rows  $B$

$$[X_i^T] = [a_i^T] [B]$$

$$X = \begin{bmatrix} \vdots \\ c \\ \vdots \end{bmatrix}^k [v^T]^d_k$$

$$\text{row}(X) \subseteq \text{row}(B) \quad \text{rank}(X) \leq \text{rank}(B)$$

## Row Span/Column Span

Let  $X \in \mathbb{R}^{n \times n}$  be a symmetric matrix with  $\text{rank}(X) = k$ .

What is the dimension of  $X$ 's row span?

$k$

$$X = \underbrace{C V^T}_{\text{basis for row span}}$$

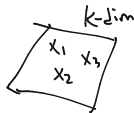
What is the dimension of  $X$ 's column span?

$k$

What is the dimension of  $X$ 's null space?

$n - k$

$$\begin{bmatrix} X \\ \vdots \\ X \end{bmatrix} \begin{bmatrix} z \\ z \end{bmatrix} = 0 = \begin{bmatrix} x_1^T z \\ \vdots \\ x_n^T z \end{bmatrix}$$



How many of  $X$ 's eigenvalues are equal to 0?

$n - k$

# SVD And Eigendecomposition

Let  $X = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^T$  be the SVD of  $X$ . Prove that the  $i^{\text{th}}$  column of  $\mathbf{V}$ ,  $\mathbf{v}_i$ , is an eigenvector of  $\mathbf{X}^T\mathbf{X}$  with eigenvalue  $\sigma_i^2$ .

$$X = \mathbf{V}\mathbf{\Lambda}\mathbf{V}^T$$

$$X \in \mathbb{R}^{n \times d}$$

$$\left[ \begin{array}{c} \mathbf{U} \\ \mathbf{\Sigma} \\ \mathbf{V}^T \end{array} \right] \left\{ \begin{array}{l} \text{orthogonal cols} \\ \text{diag} \end{array} \right.$$

$\mathbf{V}$  has ortho columns

$$\begin{array}{c} \mathbf{X}^T\mathbf{X} \\ \mathbf{X}\mathbf{X}^T \end{array}$$

$$\mathbf{X}^T\mathbf{X} = (\mathbf{U}\mathbf{\Sigma}\mathbf{V}^T)^T (\mathbf{U}\mathbf{\Sigma}\mathbf{V}^T)$$

$$\mathbf{V}\mathbf{\Sigma}\mathbf{V}^T\mathbf{U}\mathbf{\Sigma}\mathbf{V}^T$$

$$\mathbf{X}^T\mathbf{X} = \left[ \mathbf{V}\mathbf{\Sigma}^2\mathbf{V}^T \right] \left[ \mathbf{v}_i \right]$$

$$\left( \begin{array}{c} \mathbf{v}_i \\ \sigma_i^2 \end{array} \right) \mathbf{X}\mathbf{X}^T$$

$$\left[ \begin{array}{c} \mathbf{v}_i^T \\ \vdots \\ \mathbf{v}_i \end{array} \right] \left[ \mathbf{v}_i \right]$$

$$\mathbf{V} \left[ \begin{array}{c} \sigma \\ \vdots \\ \sigma_i^2 \\ \vdots \\ \sigma \end{array} \right] = \sigma_i^2 \cdot \mathbf{v}_i$$

## SVD And Eigendecomposition

Let  $\mathbf{X} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^T$  be the SVD of  $\mathbf{X}$ . Prove that the  $i^{\text{th}}$  column of  $\mathbf{U}$ ,  $\mathbf{u}_i$ , is an eigenvector of  $\mathbf{X}\mathbf{X}^T$  with eigenvalue  $\sigma_i^2$ .

## Low Rank Approximation Example

5.  $\mathbf{X} \in \mathbb{R}^{500 \times 50}$  contains 500 well-clustered data points as its rows. In particular, there are ten cluster centers  $\vec{y}_1, \dots, \vec{y}_{10} \in \mathbb{R}^{50}$ , such that each row  $\vec{x}_i$  lies within Euclidean distance at most 1 of a center. Give an *upper bound* on  $\min_{\mathbf{B}: \text{rank}(\mathbf{B})=10} \|\mathbf{X} - \mathbf{B}\|_F^2$ .

6. Consider two matrices  $\mathbf{A} = \begin{bmatrix} 1.01 & 0 \\ 0 & 1 \end{bmatrix}$  or  $\mathbf{B} = \begin{bmatrix} 1.1 & 0 \\ 0 & 1 \end{bmatrix}$ .

- (a) What are their eigenvalues and eigenvectors?
- (b) On which matrix will power method converge more quickly?

# Power Method

3. Let  $\mathbf{X} \in \mathbb{R}^{n \times d}$  have SVD  $\mathbf{X} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^T$  with singular values  $\sigma_1(\mathbf{X}), \dots, \sigma_d(\mathbf{X})$ .
- (a) What are the eigenvalues of the matrix  $(\mathbf{X}^T\mathbf{X})^2 + (\mathbf{X}^T\mathbf{X})^3$ ? What are its eigenvectors? How about the matrix  $(\mathbf{X}\mathbf{X}^T)^2 + (\mathbf{X}\mathbf{X}^T)^3$ ?
  - (b) What is the runtime required to compute  $[(\mathbf{X}^T\mathbf{X})^2 + (\mathbf{X}^T\mathbf{X})^3] \vec{v}$  for any  $\vec{v} \in \mathbb{R}^d$ .
  - (c) Name one method discussed in class which relies on efficiently applying a polynomial in  $\mathbf{X}^T\mathbf{X}$  to a vector (or more generally, applying a polynomial in a matrix  $\mathbf{A} \in \mathbb{R}^{d \times d}$  to a vector).

Johnson-Lindenstrauss

## Low Distortion Embedding Definition and JL Statement

$$x_1 \dots x_n \in \mathbb{R}^2 \xrightarrow{\text{EM}} \tilde{x}_1 \dots \tilde{x}_n \in \mathbb{R}^m \quad m \ll n$$

$$(1-\varepsilon) \|x_i - x_j\|_2 \leq \|\tilde{x}_i - \tilde{x}_j\|_2 \leq (1+\varepsilon) \|x_i - x_j\|_2$$

---

$$\|X - XVV^T\|_F$$

---

JL Lemma: always LDE into

$$m = O\left(\frac{\log(n)}{\varepsilon^2}\right) \text{ dimensions}$$

(random linear map works)

# Random Norm

Let  $\pi \in \{-1, 1\}^n$  have uniform random  $\pm 1$  entries. Let  $y \in \mathbb{R}^n$

be an arbitrary vector. What is

$$\text{Var}(X) = \mathbb{E}X^2 - (\mathbb{E}X)^2$$

$$\mathbb{E}[(\pi^T y)^2] = \text{Var}(\pi^T y)$$

$$\mathbb{E}\|\pi y\|_2^2 = \|y\|_2^2$$

$$\begin{bmatrix} 1 & -1 & -1 & \dots & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ \vdots \\ y_n \end{bmatrix}$$

$$\mathbb{E}[(\pi^T y)] = \mathbb{E} \sum_{i=1}^n \pi(i) y(i) = \sum_{i=1}^n \mathbb{E} \pi(i) y(i) = 0$$

$$\begin{aligned} \text{Var}(\pi^T y) &= \text{Var}\left(\sum_{i=1}^n \pi(i) y(i)\right) = \sum_{i=1}^n \text{Var}(\pi(i) y(i)) \\ &= \sum_{i=1}^n y(i)^2 \cdot \text{Var}(\pi(i)) \\ &= \sum_{i=1}^n y(i)^2 = \|y\|_2^2 \end{aligned}$$

$$\begin{aligned} \tilde{x}_i &= \pi x_i \\ \forall y \quad \|\pi y\|_2 &\approx \|y\|_2 \\ \text{set } y_{ij} &= x_i - x_j \end{aligned}$$

distributional JL