# COMPSCI 514: ALGORITHMS FOR DATA SCIENCE

Cameron Musco University of Massachusetts Amherst. Spring 2020. Lecture 19

### LOGISTICS

- · Problem Set 3 due this upcoming Monday at 8pm.
- · Final to be held on Zoom: May 6th from 1:00pm-3:00pm.

# Last Class: Spectral Clustering

- Splitting a graph into communities is important in network analysis and non-linear data analysis.
- · Want to find a small cut that is also balanced.
- Argued that the second smallest eigenvector of the graph Laplacian matrix can be used to find such a cut.
- Intuitive argument but not a formal proof that the identified cut is 'good'.

  Tr A = C

VLV 5.7 V1=0 cut is belance

# Last Class: Spectral Clustering

- Splitting a graph into communities is important in network analysis and non-linear data analysis.
- · Want to find a small cut that is also balanced.
- Argued that the second smallest eigenvector of the graph Laplacian matrix can be used to find such a cut.
- Intuitive argument but not a formal proof that the identified cut is 'good'.

# This Class: The Stochastic Block Model

- A simple clustered graph model where we can prove the effectiveness of spectral clustering.
- · One of the most important random graph models.

### **GENERATIVE MODELS**

**So Far:** Have argued that spectral clustering partitions a graph effectively, along a small cut that separates the graph into large pieces. But it is difficult to give any formal guarantee on the 'quality' of the partitioning in general graphs.

### **GENERATIVE MODELS**

**So Far:** Have argued that spectral clustering partitions a graph effectively, along a small cut that separates the graph into large pieces. But it is difficult to give any formal guarantee on the 'quality' of the partitioning in general graphs.

**Common Approach:** Give a natural generative model for random inputs and analyze how the algorithm performs on inputs drawn from this model.

 Very common in algorithm design for data analysis/machine learning (can be used to justify least squares regression, k-means clustering, PCA, etc.)

### STOCHASTIC BLOCK MODEL

Stochastic Block Model (Planted Partition Model): Let  $G_n(p,q)$  be a distribution over graphs on n nodes, split randomly into two groups B and C, each with n/2 nodes.

### STOCHASTIC BLOCK MODEL

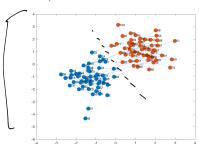
Stochastic Block Model (Planted Partition Model): Let  $G_n(p,q)$  be a distribution over graphs on n nodes, split randomly into two groups B and C, each with n/2 nodes.

- Any two nodes in the same group are connected with probability *p* (including self-loops).
- Any two nodes in different groups are connected with prob. q < p.
- Connections are independent.

### STOCHASTIC BLOCK MODEL

Stochastic Block Model (Planted Partition Model): Let  $G_n(p,q)$  be a distribution over graphs on n nodes, split randomly into two groups B and C, each with n/2 nodes. Let  $G_n(p,q)$  be a distribution over graphs on  $G_n(p,q)$  be a distribution over  $G_n(p,q)$  be a distribut

- Any two nodes in the same group are connected with probability *p* (including self-loops).
- Any two nodes in different groups are connected with prob. q < p.
- · Connections are independent.



Vn-1

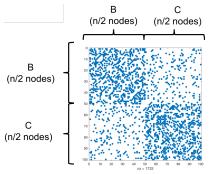
Let G be a stochastic block model graph drawn from  $G_n(p,q)$ .

Let G be a stochastic block model graph drawn from  $G_n(p,q)$ .

• Let  $A \in \mathbb{R}^{n \times n}$  be the adjacency matrix of G, ordered in terms of group ID.

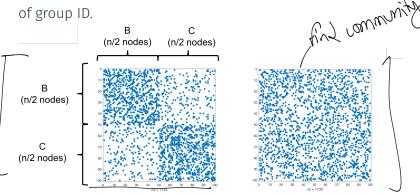
Let G be a stochastic block model graph drawn from  $G_n(p,q)$ .

• Let  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be the adjacency matrix of G, ordered in terms of group ID.



Let G be a stochastic block model graph drawn from  $G_n(p,q)$ .

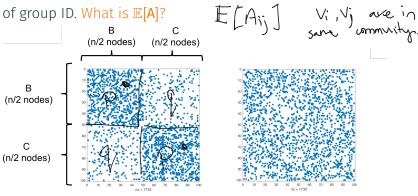
• Let  $A \in \mathbb{R}^{n \times n}$  be the adjacency matrix of G, ordered in terms of group ID.



 $G_n(p,q)$ : stochastic block model distribution. B, C: groups with n/2 nodes each. Connections are independent with probability p between nodes in the same group, and probability q between nodes not in the same group.

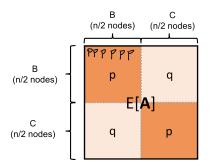
Let G be a stochastic block model graph drawn from  $G_n(p,q)$ .

· Let  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be the adjacency matrix of G, ordered in terms

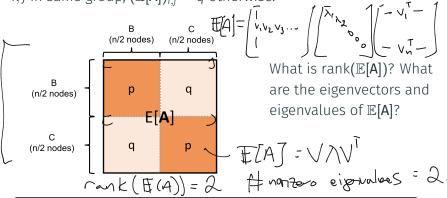


 $G_n(p,q)$ : stochastic block model distribution. B,C: groups with n/2 nodes each. Connections are independent with probability p between nodes in the same group, and probability q between nodes not in the same group.

Letting G be a stochastic block model graph drawn from  $G_n(p,q)$  and  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be its adjacency matrix.  $(\mathbb{E}[\mathbf{A}])_{i,j} = p$  for i,j in same group,  $(\mathbb{E}[\mathbf{A}])_{i,j} = q$  otherwise.

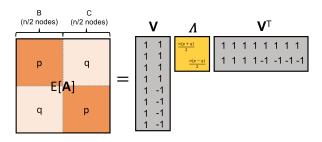


Letting G be a stochastic block model graph drawn from  $G_n(p,q)$  and  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be its adjacency matrix.  $(\mathbb{E}[\mathbf{A}])_{i,j} = p$  for i,j in same group,  $(\mathbb{E}[\mathbf{A}])_{i,j} = q$  otherwise.

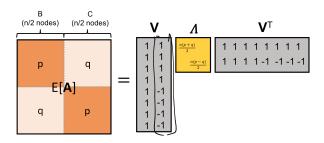


Letting G be a stochastic block model graph drawn from  $G_n(p,q)$  and  $A \in \mathbb{R}^{n \times n}$  be its adjacency matrix, what are the eigenvectors and eigenvalues of  $\mathbb{E}[A]$ ?

Letting G be a stochastic block model graph drawn from  $G_n(p,q)$  and  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be its adjacency matrix, what are the eigenvectors and eigenvalues of  $\mathbb{E}[\mathbf{A}]$ ?



If we compute  $\vec{v}_2$  then we recover the communities B and C!

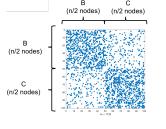


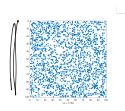
If we compute  $\vec{v}_2$  then we recover the communities B and C!

- Can show that for  $G \sim G_n(p,q)$ , **A** is close to  $\mathbb{E}[A]$  with high probability (matrix concentration inequality).
- Thus, the true second eigenvector of A is close to  $[1,1,1,\ldots,-1,-1]$  and gives a good estimate of the communities.

### SPECTRUM OF PERMUTED MATRIX

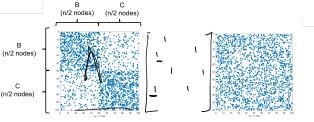
Goal is to recover communities – so adjacency matrix won't be ordered in terms of community ID (or our job is already done!)





### SPECTRUM OF PERMUTED MATRIX

Goal is to recover communities – so adjacency matrix won't be ordered in terms of community ID (or our job is already done!)



- Actual adjacency matrix is PAP<sup>T</sup> where P is a random permutation matrix and A is the ordered adjacency matrix.
- **Exercise:** The first two eigenvectors of PAP<sup>T</sup> are  $P\vec{v}_1$  and  $P\vec{v}_2$ .

$$\mathbf{P}\vec{v}_2 = [1, -1, 1, -1, \dots, 1, 1, -1]$$
 gives community ids.

Letting G be a stochastic block model graph drawn from  $G_n(p,q)$ ,  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be its adjacency matrix and  $\mathbf{L}$  be its Laplacian, what are the eigenvectors and eigenvalues of  $\mathbb{R}$ 

Laplacian, what are the eigenvectors and eigenvalues of 
$$\mathbb{E}[L]$$
?

$$\mathbb{E}[L] = \mathbb{E}[D] - \mathbb{E}[A] = \mathbb{E}[L] = \mathbb{E}[$$

$$P_{\frac{1}{2}}^{\frac{1}{2}} n \cdot P_{\frac{1}{2}}^{\frac{1}{2}} n \cdot Qn = 0$$
Letting G be a stochastic block model graph drawn from

 $G_n(p,q)$ ,  $A \in \mathbb{R}^{n \times n}$  be its adjacency matrix and L be its

Laplacian, what are the eigenvectors and eigenvalues of 
$$\mathbb{E}[L]$$
?

$$\mathbb{E}[A]: \quad \forall_1, \forall_2, \forall_3 \dots \forall_n \quad \forall_i = 0 \quad \forall i > 2.$$

$$\forall_1 \in \mathbb{F}[4] \quad \forall_1 \in \mathbb{F}[4] \quad \forall$$

$$E[L] = (P+q)^{\gamma} \vee_{l} - E[A]_{l} = (P+q)^{\gamma} \vee_{l} - (P+q)^{\gamma} \vee_{l} = (Q+q)^{\gamma} \vee_{l} = (Q+q)^{\gamma}$$

E[L] 
$$V_1 = (P+q)^n V_1 - E[AV_1 = (P+q)^n V_1 - (P+q)^n V_1 = (Q+q)^n V_1 - (P-q)^n V_2 = (Q+q)^n V_2$$

Second smallest algorithm

E[L]  $V_1 = (P+q)^n V_2 - (P-q)^n V_2 = (Q+q)^n V_2$ 

[Second smallest algorithm

E[L]  $V_1 = (P+q)^n V_1 - (P+q)^n V_2 = (Q+q)^n V_2$ 

**Upshot:** The second small eigenvector of  $\mathbb{E}[L]$  is  $\chi_{B,C}$  – the indicator vector for the cut between the communities.

**Upshot:** The second small eigenvector of  $\mathbb{E}[L]$  is  $\chi_{B,C}$  – the indicator vector for the cut between the communities.

• If the random graph *G* (equivilantly **A** and **L**) were exactly equal to its expectation, partitioning using this eigenvector would exactly recover the two communities *B* and *C*.

**Upshot:** The second small eigenvector of  $\mathbb{E}[L]$  is  $\chi_{B,C}$  – the  $\mathfrak B$  indicator vector for the cut between the communities.

• If the random graph *G* (equivilantly **A** and **L**) were exactly equal to its expectation, partitioning using this eigenvector would exactly recover the two communities *B* and *C*.

How do we show that a matrix (e.g., A) is close to its expectation? Matrix concentration inequalities.  $\sqrt{1}$ 

- Analogous to scalar concentration inequalities like Markovs,
   Chebyshevs, Bernsteins.
- Random matrix theory is a very recent and cutting edge subfield of mathematics that is being actively applied in computer science, statistics, and ML.

### MATRIX CONCENTRATION

**Matrix Concentration Inequality:** If  $p \ge O\left(\frac{\log^4 n}{n}\right)$ , then with high probability

$$\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \le O(\sqrt{pn}).$$

where  $\|\cdot\|_2$  is the matrix spectral norm (operator norm).

For any 
$$\mathbf{X} \in \mathbb{R}^{n \times d}$$
,  $\|\mathbf{X}\|_2 = \max_{z \in \mathbb{R}^d: \|z\|_2 = 1} \|\mathbf{X}z\|_2$ .

### MATRIX CONCENTRATION

**Matrix Concentration Inequality:** If  $p \ge O\left(\frac{\log^4 n}{n}\right)$ , then with high probability

$$\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \le O(\sqrt{pn}).$$

where  $\|\cdot\|_2$  is the matrix spectral norm (operator norm).

For any  $\mathbf{X} \in \mathbb{R}^{n \times d}$ ,  $\|\mathbf{X}\|_2 = \max_{z \in \mathbb{R}^d: \|z\|_2 = 1} \|\mathbf{X}z\|_2$ .

**Exercise:** Show that  $\|\mathbf{X}\|_2$  is equal to the largest singular value of  $\mathbf{X}$ . For symmetric  $\mathbf{X}$  (like  $\mathbf{A} - \mathbb{E}[\mathbf{A}]$ ) show that it is equal to the magnitude of the largest magnitude eigenvalue.

**Matrix Concentration Inequality:** If  $p \ge O\left(\frac{\log^4 n}{n}\right)$ , then with high probability

$$\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \leq O(\sqrt{pn}).$$

where  $\|\cdot\|_2$  is the matrix spectral norm (operator norm).

For any  $\mathbf{X} \in \mathbb{R}^{n \times d}$ ,  $\|\mathbf{X}\|_2 = \max_{z \in \mathbb{R}^d: \|z\|_2 = 1} \|\mathbf{X}z\|_2$ .

**Exercise:** Show that  $\|\mathbf{X}\|_2$  is equal to the largest singular value of  $\mathbf{X}$ . For symmetric  $\mathbf{X}$  (like  $\mathbf{A} - \mathbb{E}[\mathbf{A}]$ ) show that it is equal to the magnitude of the largest magnitude eigenvalue.

For the stochastic block model application, we want to show that the second eigenvectors of A and  $\mathbb{E}[A]$  are close. How does this relate to their difference in spectral norm?

Davis-Kahan Eigenvector Perturbation Theorem: Suppose  $A, \overline{A} \in \mathbb{R}^{d \times d}$  are symmetric with  $\|A - \overline{A}\|_2 \leq \epsilon$  and eigenvectors  $v_1, v_2, \ldots, v_d$  and  $\overline{v}_1, \overline{v}_2, \ldots, \overline{v}_d$ . Letting  $\theta(v_i, \overline{v}_i)$  denote the angle between  $v_i$  and  $\overline{v}_i$ , for all i:

$$sin[\theta(v_i, \bar{v}_i)] \le \frac{\epsilon}{\min_{j \ne i} |\lambda_i - \lambda_j|}$$

where  $\lambda_1, \ldots, \lambda_d$  are the eigenvalues of  $\overline{\mathbf{A}}$ .

The errors get large if there are eigenvalues with similar magnitudes.

# **EIGENVECTOR PERTURBATION**

Claim 1 (Matrix Concentration): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \le O(\sqrt{pn}).$$

Claim 2 (Davis-Kahan): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\sin \theta(v_2, \bar{v}_2) \le \frac{O(\sqrt{pn})}{\min_{j \ne i} |\lambda_i - \lambda_j|}$$

Claim 1 (Matrix Concentration): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \leq O(\sqrt{pn}).$$

Claim 2 (Davis-Kahan): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\sin \theta(v_2, \bar{v}_2) \le \frac{O(\sqrt{pn})}{\min_{j \ne i} |\lambda_i - \lambda_j|}$$

**Recall:**  $\mathbb{E}[A]$ , has eigenvalues  $\lambda_1 = \frac{(p+q)n}{2}$ ,  $\lambda_2 = \frac{(p-q)n}{2}$ ,  $\lambda_i = 0$  for  $i \ge 3$ .

Claim 1 (Matrix Concentration): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \leq O(\sqrt{pn}).$$

Claim 2 (Davis-Kahan): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\sin \theta(v_2, \bar{v}_2) \le \frac{O(\sqrt{pn})}{\min_{j \ne i} |\lambda_i - \lambda_j|}$$

**Recall:**  $\mathbb{E}[A]$ , has eigenvalues  $\lambda_1 = \frac{(p+q)n}{2}$ ,  $\lambda_2 = \frac{(p-q)n}{2}$ ,  $\lambda_i = 0$  for  $i \ge 3$ .

$$\min_{j\neq i} |\lambda_i - \lambda_j| = \min \left( qn, \frac{(p-q)n}{2} \right).$$

Claim 1 (Matrix Concentration): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,  $\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \le O(\sqrt{pn})$ .

Claim 2 (Davis-Kahan): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\sin \theta(v_2, \bar{v}_2) \le \frac{O(\sqrt{pn})}{\min_{j \ne i} |\lambda_i - \lambda_j|}$$

**Recall:**  $\mathbb{E}[A]$ , has eigenvalues  $\lambda_1 = \frac{(p+q)n}{2}$ ,  $\lambda_2 = \frac{(p-q)n}{2}$ ,  $\lambda_i = 0$  for  $i \ge 3$ .

$$\min_{j\neq i} |\lambda_i - \lambda_j| = \min \left( qn, \frac{(p-q)n}{2} \right).$$

Typically,  $\frac{(p-q)n}{2}$  will be the minimum of these two gaps.

Claim 1 (Matrix Concentration): For 
$$p \ge O\left(\frac{\log^4 n}{n}\right)$$
,  $\|\mathbf{A} - \mathbb{E}[\mathbf{A}]\|_2 \le O(\sqrt{pn})$ .

Claim 2 (Davis-Kahan): For  $p \ge O\left(\frac{\log^4 n}{n}\right)$ ,

$$\sin\theta(v_2, \overline{v}_2) \leq \frac{O(\sqrt{pn})}{\min_{j\neq i} |\lambda_i - \lambda_j|} \leq \frac{O(\sqrt{pn})}{(p-q)n/2} = O\left(\frac{\sqrt{p}}{(p-q)\sqrt{n}}\right)$$

**Recall:**  $\mathbb{E}[A]$ , has eigenvalues  $\lambda_1 = \frac{(p+q)n}{2}$ ,  $\lambda_2 = \frac{(p-q)n}{2}$ ,  $\lambda_i = 0$  for  $i \ge 3$ .

$$\min_{j\neq i} |\lambda_i - \lambda_j| = \min \left( qn, \frac{(p-q)n}{2} \right).$$

Typically,  $\frac{(p-q)n}{2}$  will be the minimum of these two gaps.

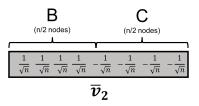
So Far:  $\sin \theta(v_2, \bar{v}_2) \leq O\left(\frac{\sqrt{p}}{(p-q)\sqrt{n}}\right)$ .

So Far:  $\sin \theta(v_2, \bar{v}_2) \leq O\left(\frac{\sqrt{p}}{(p-q)\sqrt{n}}\right)$ . What does this give us?

• Can show that this implies  $||v_2 - \bar{v}_2||_2^2 \le O\left(\frac{p}{(p-q)^2n}\right)$  (exercise).

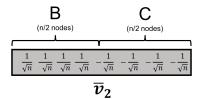
So Far:  $\sin \theta(v_2, \bar{v}_2) \leq O\left(\frac{\sqrt{p}}{(p-q)\sqrt{n}}\right)$ . What does this give us?

- · Can show that this implies  $\|v_2 \bar{v}_2\|_2^2 \le O\left(\frac{p}{(p-q)^2n}\right)$  (exercise).
- $\bar{v}_2$  is  $\frac{1}{\sqrt{n}}\chi_{B,C}$ : the community indicator vector.



So Far:  $\sin \theta(v_2, \bar{v}_2) \leq O\left(\frac{\sqrt{p}}{(p-q)\sqrt{p}}\right)$ . What does this give us?

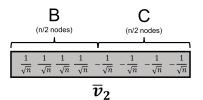
- Can show that this implies  $||v_2 \bar{v}_2||_2^2 \le O\left(\frac{p}{(p-q)^2n}\right)$  (exercise).
- $\bar{V}_2$  is  $\frac{1}{\sqrt{n}}\chi_{B,C}$ : the community indicator vector.



• Every i where  $v_2(i)$ ,  $\bar{v}_2(i)$  differ in sign contributes  $\geq \frac{1}{n}$  to  $||v_2 - \bar{v}_2||_2^2$ .

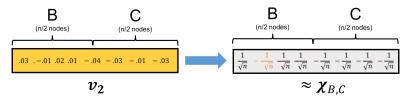
So Far:  $\sin \theta(v_2, \bar{v}_2) \leq O\left(\frac{\sqrt{p}}{(p-q)\sqrt{n}}\right)$ . What does this give us?

- Can show that this implies  $||v_2 \bar{v}_2||_2^2 \le O\left(\frac{p}{(p-q)^2n}\right)$  (exercise).
- $\bar{v}_2$  is  $\frac{1}{\sqrt{n}}\chi_{B,C}$ : the community indicator vector.

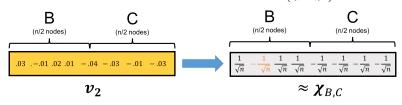


- Every *i* where  $v_2(i)$ ,  $\bar{v}_2(i)$  differ in sign contributes  $\geq \frac{1}{n}$  to  $||v_2 \bar{v}_2||_2^2$ .
- · So they differ in sign in at most  $O\left(\frac{p}{(p-q)^2}\right)$  positions.

**Upshot:** If *G* is a stochastic block model graph with adjacency matrix **A**, if we compute its second large eigenvector  $v_2$  and assign nodes to communities according to the sign pattern of this vector, we will correctly assign all but  $O\left(\frac{p}{(p-q)^2}\right)$  nodes.

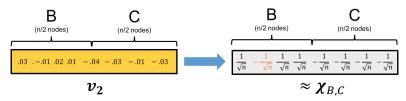


**Upshot:** If *G* is a stochastic block model graph with adjacency matrix **A**, if we compute its second large eigenvector  $v_2$  and assign nodes to communities according to the sign pattern of this vector, we will correctly assign all but  $O\left(\frac{p}{(p-q)^2}\right)$  nodes.



• Why does the error increase as q gets close to p?

**Upshot:** If *G* is a stochastic block model graph with adjacency matrix **A**, if we compute its second large eigenvector  $v_2$  and assign nodes to communities according to the sign pattern of this vector, we will correctly assign all but  $O\left(\frac{p}{(p-q)^2}\right)$  nodes.



- Why does the error increase as q gets close to p?
- Even when  $p-q=O(1/\sqrt{n})$ , assign all but an O(n) fraction of nodes correctly. E.g., assign 99% of nodes correctly.