COMPSCI 514: ALGORITHMS FOR DATA SCIENCE

Cameron Musco University of Massachusetts Amherst. Spring 2020. Lecture 13

LOGISTICS

- · Midterm is on Thursday.
- No calculators, cheatsheets, or other aids.
- Very important to do some practice problems and to try them first with no resources, to simulate the exam.
- Make sure you can recognize when to apply the fundamentals: union bound, linearity of expectation and variance, Markov's inequality, Chebyshev's inequality, indicator random variables.
- Understand the goal of each algorithm/data structure. I.e., what problem it solves with what guarantees. No need to memorize proofs.

Random hashing /hash tables) formally malyse hash table using probability tools > markovs inequality, livenity of exp. (2-univernit + pirmice hushing) Optimizing hash tables 2 level histing to broom filters - more prudice is grow for huch finding Locality susitive theship: approximate groups -vant collisions, sinthish, Minthish - length - signatures + + high tildes, 5-cerve. . uny does ninthich . Elements (Minthigh) (E, J) Streading algos - Distinct Elements (Cantimon sketch) 2

SUMMARY

Last Few Classes:

- given Er & Et some other parameter to achike these bounds - prove some prob. bonz Pr(15-Is))=1/K. - what are the requirements for each inequality - markovs: EUXJ, X>O (-law or lunge numbers Whetay: Var(X), E[X] Berns: Sun of ind. r.VS., E, Var, [-M, M]) 3

Last Few Classes:

The Johnson-Lindenstrauss Lemma

- Reduce *n* data points in any dimension *d* to $O\left(\frac{\log n/\delta}{\epsilon^2}\right)$ dimensions and preserve (with probability $\geq 1 \delta$) all pairwise distances up to $1 \pm \epsilon$.
- Compression is linear via multiplication with a random, data oblivious, matrix (linear compression)



5

Last Few Classes:

The Johnson-Lindenstrauss Lemma

- Reduce *n* data points in any dimension *d* to $O\left(\frac{\log n/\delta}{\epsilon^2}\right)$ dimensions and preserve (with probability $\geq 1 \delta$) all pairwise distances up to $1 \pm \epsilon$.
- Compression is linear via multiplication with a random, data oblivious, matrix (linear compression)

High-Dimensional Geometry

- Why high-dimensional space is so different than low-dimensional space.
- How the JL Lemma can still work.

Next Few Classes: Low-rank approximation, the SVD, and principal component analysis (PCA).

- \cdot Reduce *d*-dimesional data points to a smaller dimension *m*.
- Like JL, compression is linear by applying a matrix.

Chose this matrix carefully, taking into account structure of the dataset.

· Can give better compression than random projection.

Next Few Classes: Low-rank approximation, the SVD, and principal component analysis (PCA).

- Reduce *d*-dimesional data points to a smaller dimension *m*.
- Like JL, compression is linear by applying a matrix.
- Chose this matrix carefully, taking into account structure of the dataset.
- $\cdot\,$ Can give better compression than random projection.

Will be using a fair amount of linear algebra: orthogonal basis, column/row span, eigenvectors, etc,

RANDOMIZED ALGORITHMS UNIT TAKEAWAYS

- Randomization is an important tool in working with large datasets.
- Lets us solve 'easy' problems that get really difficult on massive datasets. Fast/space efficient look up (hash tables and bloom filters), distinct items counting, frequent items counting, near neighbor search, etc.
- The analysis of randomized algorithms leads to complex output distributions, which we can't compute exactly.
- We use concentration inequalities to bound these distributions and behaviors like accuracy, space usage, and runtime.
- Concentration inequalities and probability tools used in randomized algorithms are also fundamental in statistics, machine learning theory, probabilistic modeling of complex systems, etc.

Assume that data points $\vec{x}_1, \ldots, \vec{x}_n$ lie in any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Assume that data points $\vec{x}_1, \dots, \vec{x}_n$ lie in any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Claim: Let $\vec{v}_1, \ldots, \vec{v}_k$ be an orthonormal basis for \mathcal{V} and $\mathbf{V} \in \mathbb{R}^{d \times k}$ be the matrix with these vectors as its columns. For all \vec{x}_i, \vec{x}_j :

$$\underbrace{\left\{\begin{array}{c} \left[\begin{array}{c} \mathbf{v}_{i} & \mathbf{v}_{i} \\ \mathbf{v}_{i} & \mathbf{v}_{i} \\ \mathbf{$$

Assume that data points $\vec{x}_1, \ldots, \vec{x}_n$ lie in any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Claim: Let $\vec{v}_1, \ldots, \vec{v}_k$ be an orthonormal basis for \mathcal{V} and $\mathbf{V} \in \mathbb{R}^{d \times k}$ be the matrix with these vectors as its columns. For all \vec{x}_i, \vec{x}_i :

$$\|\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \vec{x}_j\|_2 = \|\vec{x}_i - \vec{x}_j\|_2.$$

• $\mathbf{V}^T \in \mathbb{R}^{k \times d}$ is a linear embedding of $\vec{x}_1, \dots, \vec{x}_n$ into k dimensions with no distortion.

Assume that data points $\vec{x}_1, \ldots, \vec{x}_n$ lie in any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Claim: Let $\vec{v}_1, \ldots, \vec{v}_k$ be an orthonormal basis for \mathcal{V} and $\mathbf{V} \in \mathbb{R}^{d \times k}$ be the matrix with these vectors as its columns. For all \vec{x}_i, \vec{x}_i :

$$\|\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \vec{x}_j\|_2 = \|\vec{x}_i - \vec{x}_j\|_2.$$

- $\mathbf{V}^T \in \mathbb{R}^{k \times d}$ is a linear embedding of $\vec{x}_1, \dots, \vec{x}_n$ into k dimensions with no distortion.
- An actual projection, analogous to a JL random projection $\mathbf{\Pi}$.

JG-G **Claim:** Let $\vec{v}_1, \ldots, \vec{v}_k$ be an orthonormal basis for \mathcal{V} and $\mathbf{V} \in \mathbb{R}^{d \times k}$ be the matrix with these vectors as its columns. For all, $\vec{x}_i, \vec{x}_j \in \mathcal{V}$: $\|\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \vec{x}_j\|_2^2 = \|\vec{x}_i - \vec{x}_j\|_2^2.$ $\|\nabla^T V_{ci} - \nabla^T V_{ci}\|_2^2 = \|V_{ci} - V_{cj}\|_2^2$ VIT - =11V(ci-cj)||2 ITC: - TCill2 (TV) $=(c_1-c_1)VV(c_1-c_1)$ $\|c_i - c_j\|_{L^2}$ Juk Kxl $\chi_i^{\cdot} =$ $\int (c_i - c_j)$ Sor 5 llci-cilla (\sqrt{V})

Main Focus of Today: Assume that data points $\vec{x_1}, \ldots, \vec{x_n}$ lie close to any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Main Focus of Today: Assume that data points $\vec{x_1}, \ldots, \vec{x_n}$ lie close to any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Main Focus of Today: Assume that data points $\vec{x_1}, \ldots, \vec{x_n}$ lie close to any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Letting $\vec{v}_1, \ldots, \vec{v}_k$ be an orthonormal basis for \mathcal{V} and $\mathbf{V} \in \mathbb{R}^{d \times k}$ be the matrix with these vectors as its columns, $\mathbf{V}^T \vec{x}_i \in \mathbb{R}^k$ is still a good embedding for $x_i \in \mathbb{R}^d$.

Main Focus of Today: Assume that data points $\vec{x_1}, \ldots, \vec{x_n}$ lie close to any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Letting $\vec{v}_1, \ldots, \vec{v}_k$ be an orthonormal basis for \mathcal{V} and $\mathbf{V} \in \mathbb{R}^{d \times k}$ be the matrix with these vectors as its columns, $\mathbf{V}^T \vec{x}_i \in \mathbb{R}^k$ is still a good embedding for $x_i \in \mathbb{R}^d$. The key idea behind low-rank approximation and principal component analysis (PCA).

Main Focus of Today: Assume that data points $\vec{x_1}, \ldots, \vec{x_n}$ lie close to any *k*-dimensional subspace \mathcal{V} of \mathbb{R}^d .



Letting $\vec{v}_1, \ldots, \vec{v}_k$ be an orthonormal basis for \mathcal{V} and $\mathbf{V} \in \mathbb{R}^{d \times k}$ be the matrix with these vectors as its columns, $\mathbf{V}^T \vec{x}_i \in \mathbb{R}^k$ is still a good embedding for $x_i \in \mathbb{R}^d$. The key idea behind low-rank approximation and principal component analysis (PCA).

- · How do we find ${\cal V}$ and V?
- How good is the embedding?