COMPSCI 514: Algorithms for Data Science

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University of Massachusetts Amherst. Fall 2022. Lecture 25 (Final Lecture!)

- Problem Set 5 is due Dec 12 at 11:59pm.
- Exam is next Wednesday Dec 14, from 10:30am-12:30pm in class.
- I am holding office hours Friday 12/9 2:30-4:30pm and Monday 12/12 10am-12m. Both will be held in LGRC A215.
- It would be really helpful if you could fill out SRTIs for this class (they close Dec 23).
- http://owl.umass.edu/partners/courseEvalSurvey/uma/.

Summary



Last Class:

• Analysis of gradient descent for convex and Lipschitz functions.

This Class:

• Extend gradient descent to constrained optimization via projected gradient descent.

• Course wrap up and review.

GD Analysis Proof

Theorem – GD on Convex Lipschitz Functions: For convex G-Lipschitz function f, GD run with $t \ge \frac{R^2 G^2}{\epsilon^2}$ iterations, $\eta = \frac{R}{G\sqrt{t}}$, and starting point within radius R of $\vec{\theta}_*$, outputs $\hat{\theta}$ satisfying: $f(\hat{\theta}) \le f(\vec{\theta}_*) + \epsilon$. Step 1: For all $i, f(\vec{\theta}_i) - f(\vec{\theta}_*) \le \frac{\|\vec{\theta}_i - \vec{\theta}_*\|_2^2 - \|\vec{\theta}_{i+1} - \vec{\theta}_*\|_2^2}{2\eta} + \frac{\eta G^2}{2}$

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$$f(\hat{\theta}) \leq f(\vec{\theta}_*) + \epsilon.$$

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Constrained Convex Optimization

Often want to perform convex optimization with convex constraints.

$$\underline{\vec{\theta}^*} = \arg\min_{\vec{\theta} \in S} \underline{f(\vec{\theta})},$$

where \mathcal{S} is a convex set.

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Definition – Convex Set: A set $S \subseteq \mathbb{R}^d$ is convex if and only if, for any $\vec{\theta_1}, \vec{\theta_2} \in S$ and $\lambda \in [0, 1]$:

$$(1-\lambda)\vec{\theta_1} + \lambda \cdot \vec{\theta_2} \in \mathcal{S}$$

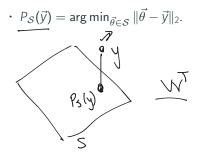
E.g. $S = \{\vec{\theta} \in \mathbb{R}^d : \|\vec{\theta}\|_2 \le 1\}.$

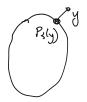
10, 12 < 1, 10, 11 < 1

$$||(1-\lambda)0, + \lambda 0z||_{2}^{\leq})$$

 $(triansle inequality_{5})$

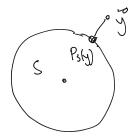
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- For $S = \{\vec{\theta} \in \mathbb{R}^d : \|\vec{\theta}\|_2 \le 1\}$ what is $P_S(\vec{y})$?







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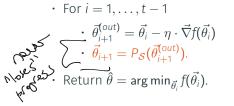
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- For $S = \{\vec{\theta} \in \mathbb{R}^d : \|\vec{\theta}\|_2 \le 1\}$ what is $P_S(\vec{y})$?
- For S being a k dimensional subspace of \mathbb{R}^d , what is $P_{\mathcal{S}}(\vec{y})$?

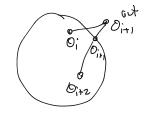
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Projected Gradient Descent

• Choose some initialization $\vec{\theta_1}$ and set $\eta = \frac{R}{G\sqrt{t}}$.





Convex Projections

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Theorem – Projection to a convex set: For any convex set $S \subseteq$ \mathbb{R}^d , $\vec{y} \in \mathbb{R}^d$, and $\vec{\theta} \in S$, $\|P_{\mathcal{S}}(\vec{y}) - \vec{\theta}\|_2 \le \|\vec{y} - \vec{\theta}\|_2.$ (1 (0 x (men proto Pel

Projected Gradient Descent Analysis

Theorem – Projected GD: For convex *G*-Lipschitz function *f*, and convex set *S*, Projected GD run with $t \ge \frac{R^2 G^2}{\epsilon^2}$ iterations, $\eta = \frac{R}{G\sqrt{t}}$, and starting point within radius *R* of $\vec{\theta}_*$, outputs $\hat{\theta}$ satisfying: $\mathcal{O}_{1,1}, \ldots, \mathcal{O}_{r} \stackrel{\text{(e)}}{\longrightarrow} f(\hat{\theta}) \le f(\vec{\theta}_*) + \epsilon = \min_{\vec{\theta} \in S} f(\vec{\theta}) + \epsilon$

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Recall:
$$\vec{\theta}_{i+1}^{(out)} = \vec{\theta}_i - \eta \cdot \vec{\nabla} f(\vec{\theta}_i)$$
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Course Review

Randomization as a computational resource for massive datasets.

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 Focus on problems that are easy on small datasets but hard at massive scale – set size estimation, load balancing, distinct elements counting (MinHash), checking set membership (Bloom Filters), frequent items counting (Count-min sketch), near neighbor search (locality sensitive hashing).

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- Just the tip of the iceberg on randomized streaming/sketching/hashing algorithms. Check out 690RA if you want to learn more.
- In the process covered probability/statistics tools that are very useful beyond algorithm design: concentration inequalities, higher moment bounds, law of large numbers, central limit theorem, linearity of expectation and variance, union bound, median as a robust estimator.

Methods for working with (compressing) high-dimensional data

• Started with randomized dimensionality reduction and the JL lemma: compression from *any* d-dimensions to $O(\log n/\epsilon^2)$ dimensions while preserving pairwise distances.

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- · Spectral graph theory nonlinear dimension reduction and spectral clustering for community detection.
- In the process covered linear algebraic tools that are very broadly useful in ML and data science: eigendecomposition, singular value decomposition, projection, norm transformations.

Foundations of continuous optimization and gradient descent.

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- Simple extension to projected gradient descent for optimization over a convex constraint set.
- Lots that we didn't cover: online and stochastic gradient descent, accelerated methods, adaptive methods, second order methods (quasi-Newton methods), practical considerations. Gave mathematical tools to understand these methods.

Thanks for a great semester!

Final Exam Questions/Review

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