

# COMPSCI 514: ALGORITHMS FOR DATA SCIENCE

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Lecture 14

- We will be grading midterms soon, and plan to return before the add/drop deadline.
- No quiz this week.

## Last Few Classes:

### The Johnson-Lindenstrauss Lemma

- Reduce  $n$  data points in any dimension  $d$  to  $O\left(\frac{\log n/\delta}{\epsilon^2}\right)$  dimensions and preserve (with probability  $\geq 1 - \delta$ ) all pairwise distances up to  $1 \pm \epsilon$ .
- Compression is linear via multiplication with a random, data oblivious, matrix (linear compression)

### High-Dimensional Geometry

- Why high-dimensional space is so different than low-dimensional space.
- How the JL Lemma can still work.

**Next Few Classes: Low-rank approximation, the SVD, and principal component analysis (PCA).**

- Reduce  $d$ -dimensional data points to a smaller dimension  $m$ .
- Like JL, **compression is linear** – by applying a matrix.
- Chose this matrix carefully, taking into account **structure of the dataset**.
- Can give better compression than random projection.

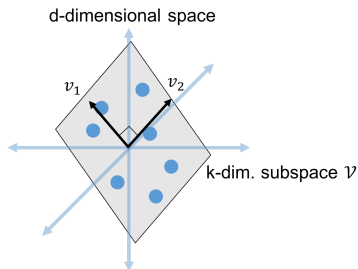
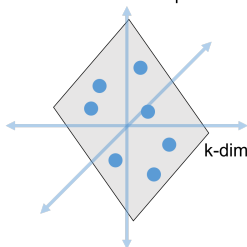
Will be using a fair amount of linear algebra: orthogonal basis, column/row span, eigenvectors, etc.

## RANDOMIZED ALGORITHMS UNIT TAKEAWAYS

- Randomization is an important tool in working with large datasets.
- Lets us solve 'easy' problems that get really difficult on massive datasets. Fast/space efficient look up (hash tables and bloom filters), distinct items counting, frequent items counting, near neighbor search (LSH), etc.
- The analysis of randomized algorithms leads to complex output distributions, which we can't compute exactly.
- We've covered many of the key ideas used through a small number of example applications/algorithms.
- We use concentration inequalities to bound these distributions and behaviors like accuracy, space usage, and runtime.
- Concentration inequalities and probability tools used in randomized algorithms are also fundamental in statistics, machine learning theory, probabilistic modeling of complex systems, etc.

## EMBEDDING WITH ASSUMPTIONS

Assume that data points  $\vec{x}_1, \dots, \vec{x}_n$  lie in any  $k$ -dimensional subspace  $\mathcal{V}$  of  $\mathbb{R}^d$  d-dimensional space



**Claim:** Let  $\vec{v}_1, \dots, \vec{v}_k$  be an orthonormal basis for  $\mathcal{V}$  and  $\mathbf{V} \in \mathbb{R}^{d \times k}$  be the matrix with these vectors as its columns. For all  $\vec{x}_i, \vec{x}_j$ :

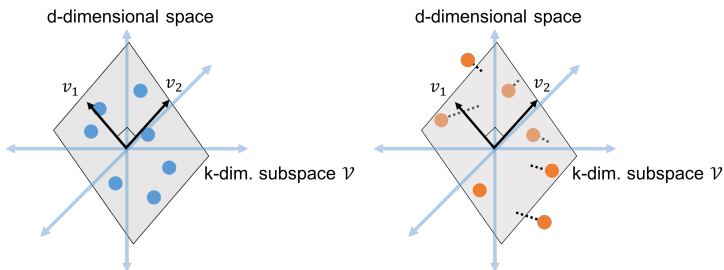
$$\|\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \vec{x}_j\|_2 = \|\vec{x}_i - \vec{x}_j\|_2.$$

- $\mathbf{V}^T \in \mathbb{R}^{k \times d}$  is a linear embedding of  $\vec{x}_1, \dots, \vec{x}_n$  into  $k$  dimensions with **no distortion**.

**Claim:** Let  $\vec{v}_1, \dots, \vec{v}_k$  be an orthonormal basis for  $\mathcal{V}$  and  $\mathbf{V} \in \mathbb{R}^{d \times k}$  be the matrix with these vectors as its columns. For all  $\vec{x}_i, \vec{x}_j \in \mathcal{V}$ :

$$\|\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \vec{x}_j\|_2 = \|\vec{x}_i - \vec{x}_j\|_2.$$

**Main Focus of Upcoming Classes:** Assume that data points  $\vec{x}_1, \dots, \vec{x}_n$  lie **close to** any  $k$ -dimensional subspace  $\mathcal{V}$  of  $\mathbb{R}^d$ .



Letting  $\vec{v}_1, \dots, \vec{v}_k$  be an orthonormal basis for  $\mathcal{V}$  and  $\mathbf{V} \in \mathbb{R}^{d \times k}$  be the matrix with these vectors as its columns,  $\mathbf{V}^T \vec{x}_i \in \mathbb{R}^k$  is **still a good embedding** for  $x_i \in \mathbb{R}^d$ . The key idea behind low-rank approximation and principal component analysis (PCA).

- How do we find  $\mathcal{V}$  and  $\mathbf{V}$ ?
- How good is the embedding?



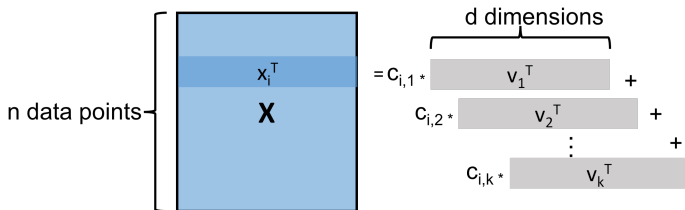
# LOW-RANK FACTORIZATION

**Claim:**  $\vec{x}_1, \dots, \vec{x}_n$  lie in a  $k$ -dimensional subspace  $\mathcal{V} \Leftrightarrow$  the data matrix  $\mathbf{X} \in \mathbb{R}^{n \times d}$  has rank  $\leq k$ .

- Letting  $\vec{v}_1, \dots, \vec{v}_k$  be an orthonormal basis for  $\mathcal{V}$ , can write any  $\vec{x}_i$  as:

$$\vec{x}_i = \mathbf{V}\vec{c}_i = c_{i,1} \cdot \vec{v}_1 + c_{i,2} \cdot \vec{v}_2 + \dots + c_{i,k} \cdot \vec{v}_k.$$

- So  $\vec{v}_1, \dots, \vec{v}_k$  span the rows of  $\mathbf{X}$  and thus  $\text{rank}(\mathbf{X}) \leq k$ .

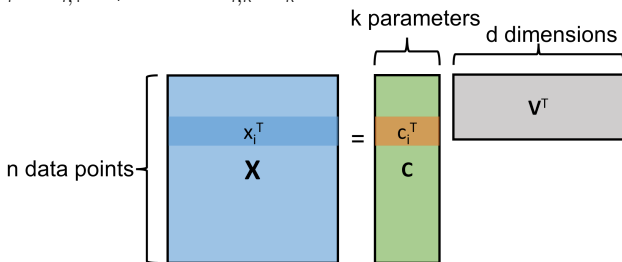


$\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^d$ : data points,  $\mathbf{X} \in \mathbb{R}^{n \times d}$ : data matrix,  $\vec{v}_1, \dots, \vec{v}_k \in \mathbb{R}^d$ : orthogonal basis for subspace  $\mathcal{V}$ .  $\mathbf{V} \in \mathbb{R}^{d \times k}$ : matrix with columns  $\vec{v}_1, \dots, \vec{v}_k$ .

**Claim:**  $\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^d$  lie in a  $k$ -dimensional subspace  $\mathcal{V} \Leftrightarrow$  the data matrix  $\mathbf{X} \in \mathbb{R}^{n \times d}$  has rank  $\leq k$ .

- Every data point  $\vec{x}_i$  (row of  $\mathbf{X}$ ) can be written as

$$\vec{x}_i = \mathbf{V}\vec{c}_i = c_{i,1} \cdot \vec{v}_1 + \dots + c_{i,k} \cdot \vec{v}_k.$$

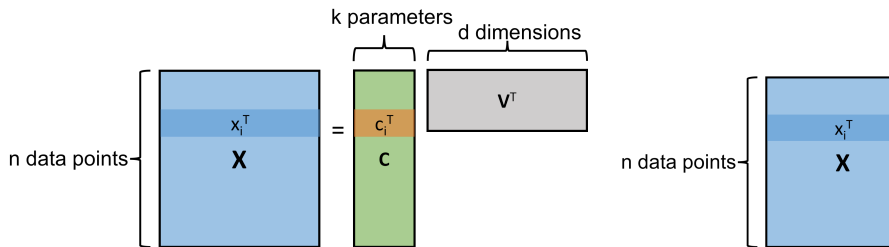


- $\mathbf{X}$  can be represented by  $(n + d) \cdot k$  parameters vs.  $n \cdot d$ .
- The rows of  $\mathbf{X}$  are spanned by  $k$  vectors: the columns of  $\mathbf{V} \implies$  the columns of  $\mathbf{X}$  are spanned by  $k$  vectors: the columns of  $\mathbf{C}$ .

$\vec{x}_1, \dots, \vec{x}_n$ : data points (in  $\mathbb{R}^d$ ),  $\mathcal{V}$ :  $k$ -dimensional subspace of  $\mathbb{R}^d$ ,  $\vec{v}_1, \dots, \vec{v}_k \in \mathbb{R}^d$ : orthogonal basis for  $\mathcal{V}$ .  $\mathbf{V} \in \mathbb{R}^{d \times k}$ : matrix with columns  $\vec{v}_1, \dots, \vec{v}_k$ .

# LOW-RANK FACTORIZATION

**Claim:** If  $\vec{x}_1, \dots, \vec{x}_n$  lie in a  $k$ -dimensional subspace with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be written as  $\mathbf{X} = \mathbf{C}\mathbf{V}^T$ .



**Exercise:** What is this coefficient matrix  $\mathbf{C}$ ? **Hint:** Use that  $\mathbf{V}^T\mathbf{V} = \mathbf{I}$ .

$$\bullet \mathbf{X} = \mathbf{C}\mathbf{V}^T \implies \mathbf{X}\mathbf{V} = \mathbf{C}\mathbf{V}^T\mathbf{V} \implies \mathbf{X}\mathbf{V} = \mathbf{C}$$

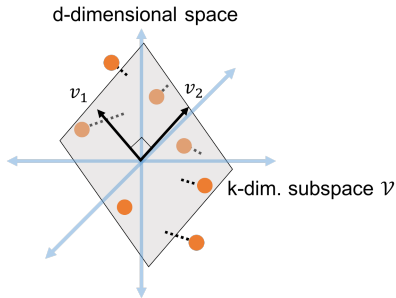
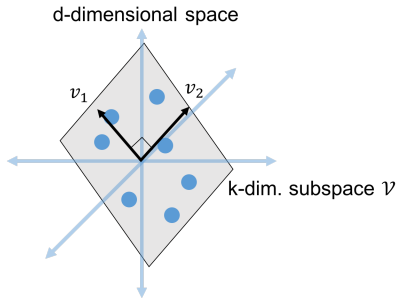
$\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^d$ : data points,  $\mathbf{X} \in \mathbb{R}^{n \times d}$ : data matrix,  $\vec{v}_1, \dots, \vec{v}_k \in \mathbb{R}^d$ : orthogonal basis for subspace  $\mathcal{V}$ .  $\mathbf{V} \in \mathbb{R}^{d \times k}$ : matrix with columns  $\vec{v}_1, \dots, \vec{v}_k$ .

# PROJECTION VIEW

**Claim:** If  $\vec{x}_1, \dots, \vec{x}_n$  lie in a  $k$ -dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be written as

$$\mathbf{X} = \mathbf{C}\mathbf{V}^T\mathbf{X}\mathbf{V}^T.$$

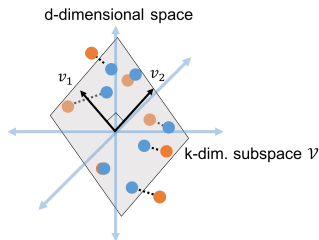
- $\mathbf{V}\mathbf{V}^T$  is a **projection matrix**, which projects the rows of  $\mathbf{X}$  (the data points  $\vec{x}_1, \dots, \vec{x}_n$ ) onto the subspace  $\mathcal{V}$ .



# LOW-RANK APPROXIMATION

**Claim:** If  $\vec{x}_1, \dots, \vec{x}_n$  lie **close to** a  $k$ -dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be **approximated as:**

$$\mathbf{X} \approx \mathbf{XV}^T$$



**Note:**  $\mathbf{XV}^T$  has rank  $k$ . It is a **low-rank approximation** of  $\mathbf{X}$ .

$$\mathbf{XV}^T = \underset{\mathbf{B} \text{ with rows in } \mathcal{V}}{\text{arg min}} \|\mathbf{X} - \mathbf{B}\|_F^2 = \sum_{i,j} (\mathbf{X}_{i,j} - \mathbf{B}_{i,j})^2.$$

$\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^d$ : data points,  $\mathbf{X} \in \mathbb{R}^{n \times d}$ : data matrix,  $\vec{v}_1, \dots, \vec{v}_k \in \mathbb{R}^d$ : orthogonal basis for subspace  $\mathcal{V}$ .  $\mathbf{V} \in \mathbb{R}^{d \times k}$ : matrix with columns  $\vec{v}_1, \dots, \vec{v}_k$ .

**So Far:** If  $\vec{x}_1, \dots, \vec{x}_n$  lie close to a  $k$ -dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be approximated as:

$$\mathbf{X} \approx \mathbf{XV}^T.$$

This is the closest approximation to  $\mathbf{X}$  with rows in  $\mathcal{V}$  (i.e., in the column span of  $\mathbf{V}$ ).

- Letting  $(\mathbf{XV}^T)_i, (\mathbf{XV}^T)_j$  be the  $i^{\text{th}}$  and  $j^{\text{th}}$  projected data points,
 
$$\|(\mathbf{XV}^T)_i - (\mathbf{XV}^T)_j\|_2 = \|[(\mathbf{XV})_i - (\mathbf{XV})_j]\mathbf{V}^T\|_2 = \|[(\mathbf{XV})_i - (\mathbf{XV})_j]\|_2.$$
- Can use  $\mathbf{XV} \in \mathbb{R}^{n \times k}$  as a compressed approximate data set.

Key question is how to find the subspace  $\mathcal{V}$  and correspondingly  $\mathbf{V}$ .

$\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^d$ : data points,  $\mathbf{X} \in \mathbb{R}^{n \times d}$ : data matrix,  $\vec{v}_1, \dots, \vec{v}_k \in \mathbb{R}^d$ : orthogonal basis for subspace  $\mathcal{V}$ .  $\mathbf{V} \in \mathbb{R}^{d \times k}$ : matrix with columns  $\vec{v}_1, \dots, \vec{v}_k$ .

**Quick Exercise:** Show that  $\mathbf{V}\mathbf{V}^T$  is **idempotent**. I.e.,  $(\mathbf{V}\mathbf{V}^T)(\mathbf{V}\mathbf{V}^T)\vec{y} = (\mathbf{V}\mathbf{V}^T)\vec{y}$  for any  $\vec{y} \in \mathbb{R}^d$ .

Why does this make sense intuitively?

**Less Quick Exercise: (Pythagorean Theorem)** Show that:

$$\|\vec{y}\|_2^2 = \|(\mathbf{V}\mathbf{V}^T)\vec{y}\|_2^2 + \|\vec{y} - (\mathbf{V}\mathbf{V}^T)\vec{y}\|_2^2.$$

## A STEP BACK: WHY LOW-RANK APPROXIMATION?

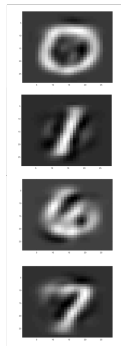
**Question:** Why might we expect  $\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^d$  to lie close to a  $k$ -dimensional subspace?

- The rows of  $\mathbf{X}$  can be approximately reconstructed from a basis of  $k$  vectors.

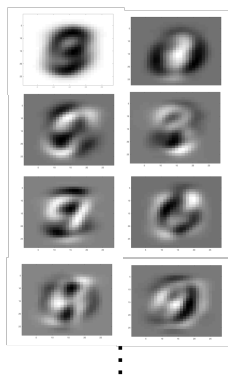
784 dimensional vectors



projections onto 15 dimensional space



orthonormal basis  $v_1, \dots, v_{15}$





# DUAL VIEW OF LOW-RANK APPROXIMATION

**Question:** Why might we expect  $\vec{x}_1, \dots, \vec{x}_n \in \mathbb{R}^d$  to lie close to a  $k$ -dimensional subspace?

- Equivalently, the columns of  $\mathbf{X}$  are approx. spanned by  $k$  vectors.

**Linearly Dependent Variables:**

	bedrooms	bathrooms	sq.ft.	floors	list price	sale price
home 1	2	2	1800	2	200,000	195,000
home 2	4	2.5	2700	1	300,000	310,000
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
home n	5	3.5	3600	3	450,000	450,000

	bedrooms
home 1	2
home 2	4
⋮	⋮
⋮	⋮
⋮	⋮
home n	5 <sup>16</sup>