## COMPSCI 514: ALGORITHMS FOR DATA SCIENCE

Cameron Musco University of Massachusetts Amherst. Fall 2020. Lecture 24

#### LOGISTICS

- · Problem Set 4 is due tomorrow at 8pm.
- · Optional Problem Set 5 will be released tomorrow, due 11/30.
- Exam will span December 3-4. Any two hour period.
- Exam review guide, practice problems, logistical details have been posted under the schedule tab on the course page.
- I am holding an optional SRTI (course reviews) for this class and would really appreciate your feedback (closes Dec 6).
- http://owl.umass.edu/partners/ courseEvalSurvey/uma/.

#### Last Class:

- · Analysis of gradient descent for optimizing convex functions.
- · Introduction to convex sets and projection functions.
- (The same) analysis of projected gradient descent for optimizing under convex functions under (convex) constraints.

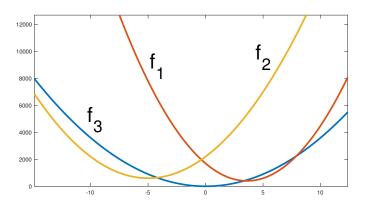
#### This Class:

- · Online learning, regret, and online gradient descent.
- · Application to stochastic gradient descent.

### **QUIZ REVIEW**

Consider the function  $f(\vec{\theta}) = \vec{x}^T \vec{\theta}$  for x = [1, -1, -2]. Give the minimum value of G such that  $f(\vec{\theta})$  is G-Lipschitz

# What does $f_1(\theta) + f_2(\theta) + f_3(\theta)$ look like?



A sum of convex functions is always convex (good exercise).

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#### ONLINE GRADIENT DESCENT

In reality many learning problems are online.

- Websites optimize ads or recommendations to show users, given continuous feedback from these users.
- Spam filters are incrementally updated and adapt as they see more examples of spam over time.
- Face recognition systems, other classification systems, learn from mistakes over time.

Want to minimize some global loss  $L(\vec{\theta}, \mathbf{X}) = \sum_{i=1}^{n} \ell(\vec{\theta}, \vec{x}_i)$ , when data points are presented in an online fashion  $\vec{x}_1, \vec{x}_2, \dots, \vec{x}_n$  (similar to streaming algorithms)

Stochastic gradient descent is a special case: when data points are considered a random order for computational reasons.

#### ONLINE OPTIMIZATION FORMAL SETUP

**Online Optimization:** In place of a single function *f*, we see a different objective function at each step:

$$f_1, f_2, \ldots, f_t : \mathbb{R}^d \to \mathbb{R}$$

- · At each step, first pick (play) a parameter vector  $\vec{\theta}^{(i)}$ .
- Then are told  $f_i$  and incur cost  $f_i(\vec{\theta}^{(i)})$ .
- Goal: Minimize total cost  $\sum_{i=1}^{t} f_i(\vec{\theta}^{(i)})$ .

Our analysis will make no assumptions on how  $f_1, \ldots, f_t$  are related to each other!

#### ONLINE OPTIMIZATION EXAMPLE

UI design via online optimization.



- · Parameter vector  $\vec{\theta}^{(i)}$ : some encoding of the layout at step i.
- Functions  $f_1, \ldots, f_t$ :  $f_i(\vec{\theta}^{(i)}) = 1$  if user does not click 'add to cart' and  $f_i(\vec{\theta}^{(i)}) = 0$  if they do click.
- Want to maximize number of purchases. I.e., minimize  $\sum_{i=1}^{t} f_i(\vec{\theta}^{(i)})$ .

## Home pricing tools.





 $\vec{x} = [\#baths, \#beds, \#floors...]$ 

- Parameter vector  $\vec{\theta}^{(i)}$ : coefficients of linear model at step *i*.
- Functions  $f_1, \ldots, f_t$ :  $f_i(\vec{\theta}^{(i)}) = (\langle \vec{x}_i, \vec{\theta}^{(i)} \rangle price_i)^2$  revealed when  $home_i$  is listed or sold.
- Want to minimize total squared error  $\sum_{i=1}^{t} f_i(\vec{\theta}^{(i)})$  (same as classic least squares regression).

In normal optimization, we seek  $\hat{\theta}$  satisfying:

$$f(\hat{\theta}) \leq \min_{\vec{\theta}} f(\vec{\theta}) + \epsilon.$$

In online optimization we will ask for the same.

$$\sum_{i=1}^{t} f_i(\vec{\theta}^{(i)}) \le \min_{\vec{\theta}} \sum_{i=1}^{t} f_i(\vec{\theta}) + \epsilon = \sum_{i=1}^{t} f_i(\vec{\theta}^{off}) + \epsilon$$

 $\epsilon$  is called the regret.

- This error metric is a bit 'unfair'. Why?
- Comparing online solution to best fixed solution in hindsight.  $\epsilon$  can be negative!

#### INTUITION CHECK

What if for  $i = 1, ..., t, f_i(\theta) = |\theta - 1000|$  or  $f_i(\theta) = |\theta + 1000|$  in an alternating pattern?

How small can the regret  $\epsilon$  be?  $\sum_{i=1}^{t} f_i(\vec{\theta}^{(i)}) \leq \sum_{i=1}^{t} f_i(\vec{\theta}^{off}) + \epsilon$ .

What if for i = 1, ..., t,  $f_i(\theta) = |\theta - 1000|$  or  $f_i(\theta) = |\theta + 1000|$  in no particular pattern? How can any online learning algorithm hope to achieve small regret?

#### ONLINE GRADIENT DESCENT

### Assume that:

- $f_1, \ldots, f_t$  are all convex.
- Each  $f_i$  is G-Lipschitz (i.e.,  $\|\vec{\nabla}f_i(\vec{\theta})\|_2 \leq G$  for all  $\vec{\theta}$ .)
- $\|\vec{\theta}^{(1)} \vec{\theta}^{off}\|_2 \le R$  where  $\theta^{(1)}$  is the first vector chosen.

## Online Gradient Descent

- Pick some initial  $\vec{\theta}^{(1)}$ .
- Set step size  $\eta = \frac{R}{G\sqrt{t}}$ .
- For  $i = 1, \ldots, t$ 
  - Play  $\vec{\theta}^{(i)}$  and incur cost  $f_i(\vec{\theta}^{(i)})$ .
  - $\cdot \vec{\theta}^{(i+1)} = \vec{\theta}^{(i)} \eta \cdot \vec{\nabla} f_i(\vec{\theta}^{(i)})$

Theorem – OGD on Convex Lipschitz Functions: For convex *G*-Lipschitz  $f_1, \ldots, f_t$ , OGD initialized with starting point  $\theta^{(1)}$  within radius R of  $\theta^{off}$ , using step size  $\eta = \frac{R}{G\sqrt{t}}$ , has regret bounded by:

$$\left[\sum_{i=1}^t f_i(\theta^{(i)}) - \sum_{i=1}^t f_i(\theta^{off})\right] \le RG\sqrt{t}$$

Upper bound on average regret goes to 0 and  $t \to \infty$ . No assumptions on  $f_1, \ldots, f_t$ !

Step 1.1: For all 
$$i$$
,  $\nabla f_i(\theta^{(i)})(\theta^{(i)} - \theta^{off}) \le \frac{\|\theta^{(i)} - \theta^{off}\|_2^2 - \|\theta^{(i+1)} - \theta^{off}\|_2^2}{2\eta} + \frac{\eta G^2}{2}$ .

Convexity  $\implies$  Step 1: For all i,

$$f_i(\theta^{(i)}) - f_i(\theta^{off}) \le \frac{\|\theta^{(i)} - \theta^{off}\|_2^2 - \|\theta^{(i+1)} - \theta^{off}\|_2^2}{2\eta} + \frac{\eta G^2}{2}.$$

#### ONLINE GRADIENT DESCENT ANALYSIS

Theorem – OGD on Convex Lipschitz Functions: For convex *G*-Lipschitz  $f_1, \ldots, f_t$ , OGD initialized with starting point  $\theta^{(1)}$  within radius R of  $\theta^{off}$ , using step size  $\eta = \frac{R}{G\sqrt{t}}$ , has regret bounded by:

$$\left[\sum_{i=1}^t f_i(\theta^{(i)}) - \sum_{i=1}^t f_i(\theta^{off})\right] \le RG\sqrt{t}$$

Step 1: For all 
$$i$$
,  $f_i(\theta^{(i)}) - f_i(\theta^{off}) \le \frac{\|\theta^{(i)} - \theta^{off}\|_2^2 - \|\theta^{(i+1)} - \theta^{off}\|_2^2}{2\eta} + \frac{\eta G^2}{2} \Longrightarrow$ 

$$\left[\sum_{i=1}^t f_i(\theta^{(i)}) - \sum_{i=1}^t f_i(\theta^{off})\right] \le \sum_{i=1}^t \frac{\|\theta^{(i)} - \theta^{off}\|_2^2 - \|\theta^{(i+1)} - \theta^{off}\|_2^2}{2\eta} + \frac{t \cdot \eta G^2}{2}.$$