## Privacy and Reliability in an Untrusted Cloud

A private and secure cloud


Distributing computation onto untrusted machines.

## Today's focus on privacy: sTile

## sTile

A technique for privately solving computationally-intensive problems (3-SAT) on untrusted computers.

## Our approach: intelligent distribution

Obstacle: Private computation is hard and inefficient [Childs 2005; Gentry 2009].


Solution:
(1) Divide computation into elemental subcomputations.
(2) Distribute subcomputations onto network.

## Computing with tiles

Input:


Program:


Computation: Copies of the program tiles self-attach to the input.

## Addition with tiles

adding program

## Addition with tiles

## adding program

Encode input to add $10\left(=1010_{2}\right)$ and $11\left(=1011_{2}\right)$


## Addition with tiles

## adding program

Add the two least significant bits


## Addition with tiles

## adding program

Add the rest of the bits, one at a time: $10+11=21\left(=10101_{2}\right)$


## Addition with tiles

## adding program

Suppose computers deployed tiles


## Addition with tiles

## adding program



Even if some were compromised, they couldn't learn private data


## 3-SAT with tiles [Winfree 1998]

## Addition [TCS'07]



3-SAT [Nat.Comp.'12]



## sTile intuition: computers simulate tiles



## sTile intuition: computers simulate tiles


discovery algorithm

## sTile intuition: computers simulate tiles



secure multi-party computation [Yao 1986]

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## sTile intuition: computers simulate tiles



## Evaluation plan

- Formally prove privacy
- Empirically demonstrate robustness to network delay
- Empirically demonstrate scalability


## Probability of reconstructing a $20-$, 38 -, and 56 -bit input


sTile provides highly-probable privacy

Threat model:
A Byzantine fraction of the cloud attempts to reconstruct private data.
sTile guarantee:
$P_{\text {compromise }}(c, n, s)=1-\left(1-c^{n}\right)^{s}$
$c$ - compromised fraction $n$ - bits in input $s$ - number of seeds

## TeraGrid example

Controlling $\frac{1}{8}$ of TeraGrid's 100,000 machines yields a probability of $10^{-10}$ of data compromise of a 17 -variable formula.

## Experimental Setup

- Mahjong: sTile implementation framework
- Java, 3K LoC, builds on Prism-MW [Malek et al. 2005]
- Input: NP-c problem instance $P$
- Output: Distributed software system to solve $P$
- Download: http://www.cs.umass.edu/~brun/Mahjong


## Experimental Setup

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- Networks
- 11-node private cluster (P4 1.5GHz, 512MiB, WinXP/2000)
- 186-node USC HPCC cluster [High Performance Computing and Communications] (P4 Xeon 3GHz, Linux)
- 100-node PlanetLab [Peterson et al. 2003] (global, varying speeds and resources)

Network Delay


VS.


Communication is $\sim 100-1000$ times faster in a CPU than on a network.

## Network Delay



Communication is $\sim 100-1000$ times faster in a CPU than on a network. But latency is not throughput!

## Robustness to Network Delay

| Problem | \# of Nodes | Network Delay | Execution Time |
| :---: | :---: | :---: | :---: |
| Mahjong |  |  |  |
| $\mathfrak{A}$ | 11 | Private Cluster | 20.1 sec . |
|  |  | HPCC | 19.3 sec . |
|  |  | PlanetLab | 18.5 sec . |
| $\mathfrak{B}$ | 11 | Private Cluster | 41.6 min . |
|  |  | HPCC | 41.2 min . |
|  |  | PlanetLab | 43.9 min . |
| Simjong |  |  |  |
| $\mathfrak{D}$ | 1,000,000 | 0ms | 65 min . |
|  |  | 10 ms | 57 min . |
|  |  | 100 ms | 64 min . |
|  |  | 500 ms | 60 min . |
|  |  | Gaussian | 68 min . |
|  |  | Distance-based | 59 min . |

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Network latency does not affect system throughput

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## Scalability: Speed $\propto$ Network Size

| Network \& Problem | \# of Nodes | Execution Time | Speed-up Ratio |
| :---: | ---: | :---: | :---: |
| Private Cluster | 5 | 43 sec. |  |
| $\mathfrak{A}$ | 10 | 23 sec. | 1.9 |
| HPCC | 93 | 220 min. |  |
| $\mathfrak{C}$ | 186 | 116 min. | 1.9 |
| PlanetLab | 50 | 9.2 min. |  |
| $\mathfrak{B}$ | 100 | 4.8 min. | 1.9 |
| Simjong | 125,000 | 8.7 hours |  |
|  | 250,000 | 4.5 hours | 1.9 |
| $\mathfrak{D}$ | 500,000 | 2.1 hours | 2.1 |
|  | $1,000,000$ | 64 min. | 2.0 |

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System speed scales almost linearly with network size

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## Related Work

- Private computation in quantum computing through entanglement [Childs 2005]
- Homomorphic encryption for private computation [Gentry 2009]
- Plethora of non-private distributed computation work [BOINC 2009; Korpela et al. 1996; Larson et al. 2002; Rosetta@home; Dean and Ghemawat 2004; Chakravarti and Baumgartner 2004]
- ... and fault-tolerant computation work
[Sarmenta 2002; Bondavalli et al. 1993, 2002; Felber and Schiper 2001; Koren and Krishna 2007; Hwang and Kesselman 2003]
- ... and private storage and access
[Ateniese et al. 2006; Wang et al. 2011; Yang et al. 2011; Yu et al. 2010]


## Contributions

## sTile

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For more, see "Entrusting Private Computation and Data to Untrusted Networks" by Y. Brun and N. Medvidovic. In IEEE Transactions on Dependable and Secure Computing (TDSC), 10(4):225-238, 2013. http://dx.doi.org/10.1109/TDSC.2013.13

# How do I compute a function using Byzantine machines? 

## How do I send you a message over a noisy channel?

## Environment model

A pool of network nodes

- some nodes are Byzantine
- Byzantine node identity and rate are unknown
- nodes may join, leave, fail, and become reliable


Applicable to problems with many independent subtasks that can be executed out of order.

## Example

- MapReduce / Hadoop [Dean and Ghemawat 2004]
- Globus Grid Toolkit [Foster et al. 2001]
- BOINC [Korpela et al. 1996]


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Crowdsourcing applications too

- reCAPTCHA [von Ahn et al. 2008]
- ESP Game [von Ahn and Dabbish 2004]
- Foldlt [Baker 2009]
- software verification [Schiller and Ernst 2010]
- AutoMan [Barowy et al. 2012]


## Voting redundancy

Assume (for now) we know average node reliability

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$1-0.3^{3}-3\left(0.3^{2}\right) 0.7 \approx 0.84$
- 19 nodes have to vote to get 0.97 reliability:
$1-\sum_{i=10}^{19}\binom{19}{i} 0.3^{i} 0.7^{19-i} \approx 0.97$


## Smart redundancy



## Smart redundancy example execution

| answers | reliability |  |  |
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| 1 | 0 |  | 0.70 |

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smart redundancy
(1) assumes best case and asks the minimum number of nodes
(2) asks more after learning how reality differs from best case.

## How many jobs to distribute?

## room 1

Flip a 70\% / 30\% coin 4 times get 4 heads and 0 tails.
room 2
Flip a 70\% / 30\% coin 1004 times get 504 heads and 500 tails.

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$$
\frac{\binom{1004}{504}\left(0.7^{504}\right) 0.3^{500}}{\binom{1004}{504}\left(0.7^{504}\right) 0.3^{500}+\binom{1004}{500}\left(0.3^{504}\right) 0.7^{500}}
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## Bayes theorem implies that given an

 a-b split of answers, only the difference affects the reliability.
## Inject redundancy only when it is needed

node reliability:
cost factor:
system reliability:


## Smart always outperforms voting redundancy

Theoretical results


## Simulation analysis confirms theoretical predictions

Simulated 1,000,000 task executions on 10,000 nodes using the XDEVS simulator [Edwards 2010] cost factor


## Empirical analysis confirms theoretical predictions

Deployed a SAT solver using BOINC [Anderson 2004] on PlanetLab [Peterson et al. 2003]


## Response time cost

response time


## Related work

## other redundancy techniques

- self-configuring optimistic programming [Bondavalli et al. 2002]
- credibility-based fault tolerance [Sarmenta 2002]
- checkpointing [Priya et al. 2007]
- crowdsourcing [Barowy et al. 2012]
- Byzantine faults in service-based computing (ZZ [Wood et al. 2011])


## complementary

- primary backup [Budhiraja et al. 1993]
- active replication [Schneider 1990]
- developer-defined fault detection [Hwang and Kesselman 2003]


## Contributions and Future Projects



## smart redundancy: using resources optimally to boost reliability

What's next?

- Channels with more bandwidth than 1 bit
- Using history to improve resource use (non-Byzantine)
- Crowdsourcing


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