

Power of software

What's going on

- User report grades posted
- 1.0 release due May 1, 11:55 PM
- Presentations April 30/May 1
 - managers will submit preference by Sunday
 - GeoMusic, Luper, and Pangea presented first last time
 - AutoAdvisor, Achieve!, and TeamAlphaStrikeForce's second
- Final team assessment due May 1 (late OK)

Today's plan

- Exam review
- Evaluations
- Power of computing
- Zero knowledge proofs

What'll be on the exam?

- testing
- debugging
- working in groups
- reasoning about programs
- ~~hacker cred~~

Testing

- Know about different kinds of tests
 - unit, integration, regression, etc.
- Know about different kinds of coverage
 - statement, path, etc.
- Know what's hard about testing
 - GUI, usability, covering all behavior, etc.

Debugging

- Know four kinds of defense against bugs
 - make impossible
 - don't introduce
 - make errors visible
 - last resort: debugging
- Rep invariants
- Assertions

Working in groups

- What's hard?
 - corner cases
 - complete specification covers **A LOT** of behavior
 - unless a spec is concise, it's hard to understand
 - precision is hard: language is ambiguous
 - communication is important

Reasoning about programs

- Ways to verify your code
 - testing, reasoning, proving
- Forward reasoning
- Backward reasoning
- Loop invariants
- Induction
- Practice some examples!

Loop example

Find the weakest precondition

```
for (int x = 1; x <> y;) {
  if (y > x) {
    y = y / 2;
    x=2*x;
  }
}
// postcondition: x=8, y=8, and x and y are ints
```

you can also find the loop invariant and decrement function

What's the test going to be like?

- 2 thinking questions
- the rest will be multiple choice
- will be designed to take less than an hour

On Tuesday, come with review questions and I will answer them.

Power of Computing

Can you write any program I describe to you?

Can you write:

A program HALTS? whose input is the body of a method, and that outputs 0 if the method enters an infinite loop, and 1 if it does not.

What's HALTS?(method)?

```
method() {
  print "hello world";
}
```

What's HALTS?(method)?

```
method() {
  for (int x=0; x<5; x++)
    print "hello world";
}
```

What's HALTS?(method)?

```
method() {
  for (int x=0; x<-1; x++)
    print "hello world";
}
```

What's HALTS?(method)?

```
method() {
  while (true);
}
```

What's HALTS?(method)?

```
method() {
  int x = 785th digit of  $\pi$ ;
  if (x == 7)
    while(true);
}
```

What's HALTS?(method)?

```
method() {
  int x = 785th digit of  $\pi$ ;
  int y =  $x^x^{x^{x^{x+1}}}$ ;
  int z = yth digit of  $\pi$ ;
  if (z == 0)
    while(true);
}
```

What's HALTS?(method)?

```
method() {
  int x = 785th digit of  $\pi$ ;
  int y =  $x^x \cdot x^{x+1}$ ;
  int[] z[] = the  $y^{\text{th}}$  through  $(x+y)^{\text{th}}$ 
  digits of  $\pi$ ;
  if (z ever repeats in  $\pi$  again)
    while(true);
}
```

How about the general case?

- Let's count programs. How many programs are there?
- And how many problems are there?
 - let's limit ourselves to simple problems:
 - given a set of numbers, e.g., {2, 4, 6}
 - on input i , return 1 if i is in the set, and 0 otherwise

First 64 programs

- How many of our problems can I solve with 64 programs?
 - 64
 - 32
 - 8
 - 6
 - 2

First 64 programs

- With 64 programs, how large can my sets get (if I am being compact)
 - 64
 - 32
 - 8
 - 6
 - 2
- Example: with 4 programs, I could cover:
 - {}, {1}, {2}, {1,2}

Scalability Problem

- To cover subsets of a set of n numbers, I need 2^n **programs**.
- But I only have as many **programs** as there are natural numbers.
- That's exponentially smaller than the number of **problems** there are.

Can't do it for all subsets!

Can HALTS? exist?

- Imagine that you wrote HALTS?
- I will write a new program NALTS?:


```
NALTS?(Method p) {
  if (HALTS?(p)==0) return 1;
  else while (true);
}
```

Key, run the program on (almost) itself
 What is the value of
 NALTS? (NALTS?)

What is the value of NALS?(NALS?)

- Two cases:
 1. If NALS?(NALS?) goes into an infinite loop, then $\text{HALTS?}(\text{NALS?})=1$, which means that NALS? terminates.
So case 1 is impossible.
 2. If NALS?(NALS?) does not go into an infinite loop, then $\text{HALTS?}(\text{NALS?})=0$, which means that NALS? does not terminate.
So case 2 is impossible.

Conclusion

- The program HALTS cannot exist!
- Many programs cannot exist!

- Learn more in CS 601

Zero-Knowledge Proofs

How can I prove to you I know X without telling you anything about X?