# Constituency Parsing: CKY 

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## Context-Free Grammar

- CFG describes a generative process for an (infinite) set of strings
- I. Nonterminal symbols
- "S": START symbol / "Sentence" symbol
- 2.Terminal symbols: word vocabulary
- 3. Rules (a.k.a. Productions). Practically, two types:
"Grammar": one NT expands to >=I NT always one NT on left side of rulep

| $S \rightarrow$ | $N P V P$ | I + want a morning flight |
| :---: | :---: | :---: |
| $N P \rightarrow$ | Pronoun | I |
| Nominal $\xrightarrow{\mid}$ | Proper-Noun | Los Angeles |
|  | Det Nominal | a + flight |
|  | Nominal Noun | morning + flight |
|  | Noun | flights |
| $V P \rightarrow$ | Verb | do |
|  | Verb N | want + a flight |
|  | Verb NP P | leave + Boston + in the morning |
|  | Verb PP | leaving + on Thursday |
| $P P \rightarrow$ | Preposition NP | from + Los Angeles |

Lexicon: NT expands to a terminal


## Constituent Parse Trees



## Bracket notation

(12.2)

<=> Set of non-terminal spans (start,end positions) $\{(N P, 0, I),(V P, I, 5),(N P, 2,5), \ldots\}$

## Parsing with a CFG

- Task: given text and a CFG, answer:
- Does there exist at least one parse?
- Enumerate parses (backpointers)
- Problem: extremely high number of possible trees for a sentence, and even a large number of legal trees (licensed by the grammar) for a sentence
- Many parsing algorithms have been invented to tackle this
- Cocke-Kasami-Younger algorithm (CKY)
- Bottom-up dynamic programming:

Find possible nonterminals for short spans of sentence, then possible combinations for higher spans

- Maintains local ambiguity, representing many subtrees for each span. ("Packed forest" representation)
- Provably finds all possible parse trees (legal derivations), and correctly says when none exist.
- Requires converting to Chomsky Normal Form (binarization)


## Chomsky Normal Form

## CKY

| Grammar |
| :--- |
| Adj -> yummy |
| NP -> foods |
| NP -> store |
| NP -> NP NP |
| NP -> Adj NP |



For cell [i,j] (loop through them bottom-up) For possible splitpoint $k=(i+1) . .(j-I)$ : For every $B$ in $[i, k]$ and $C$ in $[k, j]$, If exists rule $A->B C$, add $A$ to cell [i,j] (Recognizer) ... or ... add (A,B,C, k) to cell [i,j] (Parser)

Recognizer: per span, record list of possible nonterminals

Parser: per span, record possible ways the nonterminal was constructed.

## CKY

| Grammar |
| :--- |
| Adj $->$ yummy |
| NP -> foods |
| NP -> store |
| NP $\rightarrow$ NP NP |
| NP $\rightarrow$ Adj NP |



For cell [ $\mathrm{i}, \mathrm{j}]$ (loop through them bottom-up) For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{l}) . .(\mathrm{j}-\mathrm{l})$ : For every B in $[i, k]$ and $C$ in $[k, j]$, If exists rule A -> B C, add $A$ to cell $[i, j] \quad$ Recognizer) ... or ... add (A,B,C, k) to cell [i,j] (Parser)

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## CKY

| Grammar |
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For cell [i,j] (loop through them bottom-up) For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{l}) . .(\mathrm{j}-\mathrm{l})$ : For every B in $[\mathrm{i}, \mathrm{k}]$ and C in $[\mathrm{k}, \mathrm{j}]$, If exists rule A -> B C, add $A$ to cell $[i, j] \quad$ Recognizer) ... or ... add (A,B,C, k) to cell [i,j] (Parser)

Recognizer: per span, record list of possible nonterminals

Parser: per span, record possible ways the nonterminal was constructed.

$$
\begin{aligned}
\mathrm{S} & \rightarrow \mathrm{NP} \mathrm{VP} \\
\mathrm{NP} & \rightarrow \mathrm{NP} P \mathrm{PP} \\
\mathrm{VP} & \rightarrow \mathrm{~V} \mathrm{NP} \\
\mathrm{VP} & \rightarrow \mathrm{VP} \mathrm{PP} \\
\mathrm{PP} & \rightarrow \mathrm{P} \mathrm{NP} \\
\mathrm{NP} & \rightarrow \text { she } \\
\mathrm{NP} & \rightarrow \text { fish } \\
\mathrm{NP} & \rightarrow \text { fork } \\
\mathrm{NP} & \rightarrow \text { chopsticks } \\
\mathrm{V} & \rightarrow \text { eats } \\
\mathrm{V} & \rightarrow \text { fish } \\
\mathrm{P} & \rightarrow \text { with }
\end{aligned}
$$

| she | eats | fish | with | chopsticks |
| :---: | :---: | :---: | :---: | :---: |
|  | eats |  |  | chopsticks |
|  |  |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |
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|  |  |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |
|  |  |  |  |  |
| she | eats | fish | with | chopsticks |

$\qquad$
Fill in the CYK dynamic programming table to parse the sentence below. In the bottom right corner, draw the two parse trees. Show the possible nonterminals in each cell. Optional: draw the backpointers too.


$$
\begin{array}{rl}
\mathrm{S} \rightarrow \mathrm{NP} \mathrm{VP} & \mathrm{NP} \rightarrow \text { she } \\
\mathrm{NP} \rightarrow \mathrm{NP} \mathrm{PP} & \mathrm{NP} \rightarrow \text { fish } \\
\mathrm{VP} \rightarrow \mathrm{~V} \mathrm{NP} & \mathrm{NP} \rightarrow \text { fork } \\
\mathrm{VP} \rightarrow \mathrm{VP} \mathrm{PP} & \mathrm{NP} \rightarrow \text { chopsticks } \\
\mathrm{PP} \rightarrow \mathrm{P} \mathrm{NP} & \mathrm{~V} \rightarrow \text { eats } \\
& \mathrm{V} \rightarrow \text { fish } \\
& \mathrm{P} \rightarrow \text { with }
\end{array}
$$

## 5

## For cell [i,j]

For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{I}) . .(\mathrm{j}-\mathrm{I})$ :
For every $B$ in $[i, k]$ and $C$ in $[k, j]$,
If exists rule $A->B C$, add A to cell $[\mathrm{i}, \mathrm{j}]$

## Computational Complexity?

How do we fill in $C(1,2)$ ?

123


,

## For cell [i,j]

For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{I}) . .(\mathrm{j}-\mathrm{I})$ :
For every $B$ in $[i, k]$ and $C$ in $[k, j]$, If exists rule $A->B C$, add A to cell $[\mathrm{i}, \mathrm{j}]$

## Computational Complexity ?

How do we fill in $C(1,2)$ ?
Put together $C(1,1)$ and $C(2,2)$.

[Example from Noah Smith]

## For cell [i,j]

For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{I}) . .(\mathrm{j}-\mathrm{I})$ :
For every $B$ in $[i, k]$ and $C$ in $[k, j]$,
If exists rule $A$-> B C, add A to cell $[\mathrm{i}, \mathrm{j}]$

## Computational Complexity?

How do we fill in $C(1,3)$ ?

1
23

[Example from Noah Smith]

## For cell [i,j]

For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{I}) . .(\mathrm{j}-\mathrm{I})$ :
For every $B$ in $[i, k]$ and $C$ in $[k, j]$,
If exists rule $A$-> B C, add A to cell $[\mathrm{i}, \mathrm{j}]$

## Computational Complexity?

How do we fill in $C(1,3)$ ?
One way ...

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23

[Example from Noah Smith]

## For cell [i,j]

For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{I}) . .(\mathrm{j}-\mathrm{I})$ :
For every $B$ in $[i, k]$ and $C$ in $[k, j]$,
If exists rule $A->B C$, add A to cell $[\mathrm{i}, \mathrm{j}]$

## Computational Complexity?

How do we fill in $C(1,3)$ ?
One way ...
Another way.

[Example from Noah Smith]

## For cell [i,j]

For possible splitpoint $\mathrm{k}=(\mathrm{i}+\mathrm{I}) . .(\mathrm{j}-\mathrm{I})$ :
For every $B$ in $[i, k]$ and $C$ in [ $k, j]$,
If exists rule A -> B C, add A to cell $[\mathrm{i}, \mathrm{j}]$

How do we fill in $C(1, n)$ ?

1
2

## Computational Complexity ?


[Example from Noah Smith]

## For cell [i,j]

For possible splitpoint $k=(i+I) . .(j-I)$ :
For every $B$ in $[i, k]$ and $C$ in $[k, j]$, If exists rule $A->B C$, add $A$ to cell [i,j]

How do we fill in $C(1, n)$ ?

## Computational Complexity?


[Example from Noah Smith]

## Probabilistic CFGs

| $S \rightarrow$ NP VP | $[.80]$ |
| :--- | :--- |
| $S \rightarrow$ Aux NP VP | $[.15]$ |
| $S \rightarrow V P$ | $[.05]$ |
| $N P \rightarrow$ Pronoun | $[.35]$ |
| $N P \rightarrow$ Proper-Noun | $[.30]$ |
| $N P \rightarrow$ Det Nominal | $[.20]$ |
| $N P \rightarrow$ Nominal | $[.15]$ |
| Nominal $\rightarrow$ Noun | $[.75]$ |
| Nominal $\rightarrow$ Nominal Noun | $[.20]$ |
| Nominal $\rightarrow$ Nominal PP | $[.05]$ |
| $V P \rightarrow$ Verb | $[.35]$ |
| $V P \rightarrow$ Verb NP | $[.20]$ |
| $V P \rightarrow$ Verb NP PP | $[.10]$ |
| $V P \rightarrow$ Verb PP | $[.15]$ |
| $V P \rightarrow$ Verb NP NP | $[.05]$ |
| $V P \rightarrow$ VP PP | $[.15]$ |
| $P P \rightarrow$ Preposition NP | $[1.0]$ |

```
Det }->\mathrm{ that [.10]| | [.30] | the [.60]
Noun }->\mathrm{ book [.10] | flight [.30]
            | meal [.15] | money [.05]
    | flights [.40] | dinner [.10]
```

Verb $\rightarrow$ book [.30] | include [.30]
| prefer; [.40]
Pronoun $\rightarrow I$ [.40] $\mid$ she [.05]
| me [.15] | you [.40]
Proper-Noun $\rightarrow$ Houston [.60]
| TWA [.40]
Aux $\rightarrow$ does $[.60] \mid$ can [40]
Preposition $\rightarrow$ from [.30] | to [.30]
| on [.20] | near [.15]
| through [.05]

- Defines a probabilistic generative process for words in a sentence
- (How to learn? Fully supervised with a treebank...)
( (S
(NP-SBJ (NNP General) (NNP Electric) (NNP Co.) )
(VP (VBD said)
(SBAR (-NONE- 0)
(S
(NP-SBJ (PRP it) )
(VP (VBD signed)
(NP
(NP (DT a) (NN contract) )
(PP (-NONE- *ICH*-3) ))
(PP (IN with)
(NP
(NP (DT the) (NNS developers) )
(PP (IN of)
(NP (DT the) (NNP Ocean) (NNP State) (NNP Power) (NN project) ))))

Penn Treebank
(NP (DT the) (JJ second) (NN phase) ) (PP (IN of)
(NP
(NP (DT an) (JJ independent)
(ADJP
(QP (\$ \$) (CD 400) (CD million) )
(-NONE- *U*) )
(NN power) (NN plant) )
(, , , )
(WHNP-2 (WDT which) )
(S
(NP-SBJ-1 (-NONE- *T*-2) )
(VP (VBZ is)
(VP (VBG being)
(VP (VBN built)
(NP (-NONE-*-1) ) (PP-LOC (IN in)
(NP
(NP (NNP Burrillville) )
(, , )


PCFG as LM

## Is a PCFG a good LM? Yes...

## Is a PCFG a good LM? No...

## (P)CFG model, (P)CKY algorithm

- CKY: given CFG and sentence w
- Does there exist at least one parse?
- Enumerate parses (backpointers)
- Probabilistic CKY: given PCFG and sentence w
- Likelihood of sentence $P(w)$
- Most probable parse ("Viterbi parse") $\operatorname{argmaxy} \mathrm{P}(\mathrm{y} \mid \mathrm{w})=\operatorname{argmaxy} \mathrm{P}(\mathrm{y}, \mathrm{w})$
- Parsing model accuracy: lots of ambiguity!!
- PCFGs lack lexical information to resolve ambiguities (sneak in world knowledge?)
- Modern constituent parsers: enrich PCFG with lexical information and fine-grained nonterminals
- Modern dependency parsers: effectively the same trick
- Parsers' computational efficiency
- Grammar constant; pruning \& heuristic search
- $\mathrm{O}\left(\mathrm{N}^{3}\right)$ for CKY (ok? depends...)
- $\mathrm{O}(\mathrm{N})$ left-to-right incremental algorithms
- What was the syntactic training data?

