

# COMPSCI 575/MATH 513

## Combinatorics and Graph Theory

Lecture #26: Inclusion-Exclusion

(Tucker Sections 8.1, 8.2)

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# Inclusion-Exclusion

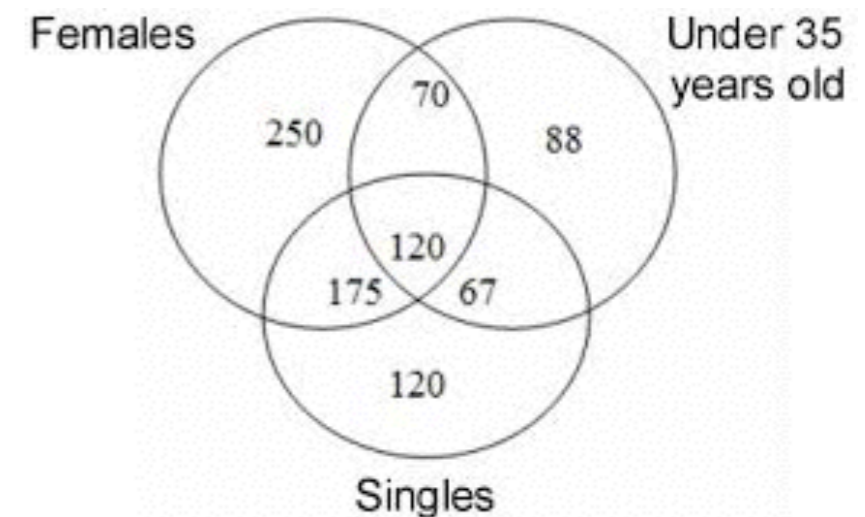
- Counting With Venn Diagrams
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# Counting With Venn Diagrams

- The single most fundamental rule for counting is the Sum Rule: if  $A$  and  $B$  are disjoint sets,  $|A \cup B| = |A| + |B|$ .
- Just after that is the Sum Rule With Overlap,  $|A \cup B| = |A| + |B| - |A \cap B|$ . We count both sets, and remove the double-counted items.
- We can expand this to more than two sets.

# Counting With Venn Diagrams

- If I tell you that my total set is made up of 615 females, 345 young people, and 482 singles, that is not enough to compute its size.
- If in addition I say that there are 190 young females, 187 young singles, and 120 young single females, you can complete this diagram and compute the total of 890.



# Restricted Arrangements

- Let's apply this to a counting problem. How many permutations of  $\{0, \dots, 9\}$  have first digit greater than 1 and last digit less than 8?
- It's all of them ( $10!$ ) minus those starting with 0 or 1 ( $2 \cdot 9!$ ), minus those ending with 8 or 9 ( $2 \cdot 9!$ ) *plus* those that both start with 0 or 1 and end with 8 or 9 ( $2 \cdot 2 \cdot 8!$ ). The elements of the last set were subtracted twice and so must be added back in once.

# Relatively Prime Numbers

- The set  $\mathbb{Z}_{70}$  of integers mod 70 has a subset  $\mathbb{Z}_{70}^*$  of elements relatively prime to 70.
- To count  $\mathbb{Z}_{70}^*$ , we start with 70, subtract the evens (35), multiples of 5 (14), and multiples of 7 (10), then add in the 7 multiples of 10, 5 multiples of 14, and 2 multiples of 35, then subtract 0 which has so far been subtracted twice and added in twice. The total of  $70 - (35+14+10) + (7+5+2) - 1 = (2-1)(5-1)(7-1) = 24$ , as we could also find from the Chinese Remainder Theorem.

# More Examples

- How many  $n$ -digit sequences over  $\{0,1,2\}$  have at least one 0, at least one 1, and at least one 2? This isn't hard, but let's develop some useful notation. Let  $A_0$ ,  $A_1$ , and  $A_2$  be the sequences without 0, 1, and 2 respectively.
- The set we want to count is  $A_0' \cap A_1' \cap A_2'$  where  $'$  means complement because I can't do overlines.
- We compute it as  $3^n - 3(2^n) + 3(1^n) - 0$ .

# More Examples

- If I have 100 students, 40 each are taking French, Latin, and German, and I am also given that 20 take only French, 20 only Latin, and 15 only German. 10 take both French and Latin.
- Is this enough to compute everything? We have eight possible combinations, represented in a three-set Venn diagram.
- We'll do this on the board, getting the conclusion that 15 students take none.



# The Inclusion-Exclusion Formula

- In general, when we have a set of  $N$  elements and subsets  $A_1, \dots, A_n$ , we will now write  $N(A_1' A_2' \dots A_n')$  to mean the number of elements not in any of the  $A_i$ 's, omitting the  $n$  symbols.
- We let  $S_1 = |A_1| + \dots + |A_n|$ ,  $S_2$  be the sum of sizes of all the intersections  $A_i \cap A_j$ ,  $S_3$  be the sum of sizes of all 3-set intersections, and so on through all  $S_k$  to  $S_n$ , the size of the intersection of all  $n$  sets.

# The Inclusion-Exclusion Formula

- With this notation, we have a theorem that the number  $N(A_1' \dots A_n')$  of elements in none of the sets is  $N - S_1 + S_2 - S_3 + \dots + (-1)^n S_n$ .
- To prove this, we look at an arbitrary element of the whole set and see how many times it is counted. If it is in  $m$  of the  $A_i$ 's, we count it once in  $N$ , subtract it  $m$  times in  $S_1$ , add it back in  $C(m, 2)$  times, subtract it  $C(m, 3)$ , etc.
- This sum is 1 if  $m = 0$  and 0 otherwise.

# Counting a Union

- Suppose I want the size of the union  $A_1 \cup \dots \cup A_n$ . This is just  $N$  minus the number we just computed, which is thus  $S_1 - S_2 + S_3 - \dots + (-1)^{n+1} S_n$ .
- This generalizes the formulas we have been using,  $|A_1 \cup A_2| = S_1 - S_2$  and  $|A_1 \cup A_2 \cup A_3| = S_1 - S_2 + S_3$ .

# Card Hands With No Voids

- Of the  $C(52, 6)$  possible six-card hands from a standard deck, how many have at least one card of each suit?
- We let  $A_1$  be the set of hands with no spades,  $A_2$  those with no hearts,  $A_3$  no diamonds, and  $A_4$  no clubs. We can easily see that  $|A_i| = C(39, 6)$ , that  $|A_i \cap A_j| = C(26, 6)$ , that  $|A_i \cap A_j \cap A_k| = C(13, 6)$ , and that  $|A_1 \cap A_2 \cap A_3 \cap A_4| = 0$ .
- The IE formula gives us  $C(52, 6) - 4C(39, 6) + 6C(26, 6) - 4C(13, 6) + 0$ .

# Upper Bounds on Solutions

- We earlier used generating functions to attack problems like the following: How many solutions to  $x_1 + \dots + x_6 = 20$  have  $0 \leq x_i \leq 8$ ?
- We let  $A_i$  be the set of non-negative solutions to this sum that have  $x_i > 8$ .
- We know  $N = C(20+6-1, 20)$ , that  $|A_i| = C(11+6-1, 11)$ , and that  $|A_i \cap A_j| = C(2+6-1, 2)$ . (The larger intersections are empty.) So our answer is  $C(25, 20) - 6C(15, 11) + 15C(7, 2)$ .

# The Derangement Problem

- If  $n$  people each check a hat and the hats are returned to them randomly, what is the probability that *no one* gets their own hat?
- We let  $N = n!$  be the set of all permutations, and let  $A_i$  be the ones where person  $i$  gets their own hat back.
- It is easy to see that  $|A_i| = (n-1)!$ , and that the intersection of  $k$   $A_i$ 's has size  $(n-k)!$ .

# The Derangement Problem

- The number  $S_k$  is thus  $C(n, k)(n-k)! = n!/k!$ .
- Applying the IE formula, we compute the number of permutations not in any  $A_i$  as  $n!$  times the sum for  $k$  from 0 to  $n$  of  $(-1)^k/k!$ .
- This last sum is the sum of the first  $n+1$  terms of the power series for  $e^{-1} = 1/e$ , and this number is the probability that no one gets their own hat.

# More on Derangements

- The number  $D_n$  of derangements is also given by the recurrence  $D_n = nD_{n-1} + (-1)^n$  for  $n \geq 2$ , with  $D_0 = 1$  and  $D_1 = 0$ .
- This recurrence can be derived by arithmetic from the more natural  $D_n = (n-1)(D_{n-1} + D_{n-2})$ , which comes from seeing where the first item goes and whether it is in an orbit of size 2.
- This recurrence can also be used to get the EGF for  $D_n$ , which is  $e^{-x}/(1-x)$ .



# Chromatic Polynomials Again

- Let  $G$  be a graph with vertices  $x_1, x_2, x_3,$  and  $x_4,$  and five edges (all but  $x_2x_4$ ). How many legal  $n$ -colorings does this graph have?
- We let  $N$  be the total number of colorings of the four vertices,  $n^4,$  and let  $A_1, \dots, A_5$  be the sets of colorings that fail the test for each of the five edges.
- We have  $|A_i| = n^3$  for each edge, and  $|A_i \cap A_j| = n^2$  for each pair of edges.

# Chromatic Polynomials Again

- What about the three-way intersections? Of the ten sets of three edges, two form triangles and can each have a common color in  $n^2$  ways, while the other eight involve all four vertices and allow only  $n$  colorings. The four-way and five-way intersections each allow  $n$  colorings.
- The chromatic polynomial, our answer, is thus  $n^4 - 5n^3 + 10n^2 - [2n^2 + 8n] + 5n - n$  which is  $n^4 - 5n^3 + 8n^2 - 4n$ .

# Number in $m$ or $\geq m$ Sets

- Again let us have a set of  $N$  elements, with  $n$  subsets  $A_1, \dots, A_n$ . How many of the elements are in exactly  $m$  of the subsets? Call this  $N_m$ .
- We can start with  $S_m$ , the sum of the sizes of all  $m$ -way intersections. Every element in exactly  $m$  sets will be in exactly one of these  $m$ -way intersections.
- But the elements in more than  $m$  subsets will also be counted multiple times in  $S_m$ .

# Number in $m$ or $\geq m$ Sets

- The correct formula for  $N_m$  is similar to the IE formula:  $N_m = S_m - C(m+1, m)S_{m+1} + C(m+2, m)S_{m+2} + \dots + (-1)^{n-m}C(n, m)S_n$ .
- The similar formula for  $N_m^*$ , the number of elements in at least  $m$  sets, is  $N_m^* = S_m - C(m, m-1)S_{m+1} + C(m+1, m-1)S_{m+2} + \dots + (-1)^{n-m}C(n-1, m-1)S_n$ .
- We verify each of these formulas by showing that each element we don't want is counted a net of zero times.

# One More Example

- Consider strings of length 4 over  $\{0, 1, 2\}$  with exactly two 1's. If  $A_i$  is the set of strings with a 1 in position  $i$ , we want the number  $N_2$ . By the formula this is  $S_2 - C(3, 2)S_3 + C(4, 2)S_4$ . (Tucker has some typos here.)
- Since  $S_2 = C(4, 2)3^2 = 54$ ,  $S_3 = C(4, 3)3^1 = 12$ , and  $S_4 = 1$ , we have  $54 - 3 \cdot 12 + 6 \cdot 1 = 24$ .
- Similarly  $N_2^* = S_2 - C(2, 1)S_3 + C(3, 1)S_4 = 54 - 2 \cdot 12 + 3 \cdot 1 = 33$ .