CS-690C: Homework 1

Problem 1. (100 points.) We define target-key recovery in what is called the *ideal cipher model*. In this model, first proposed by Shannon, a blockcipher is modeled by a different, independent random permutation for *every* key. That is, the adversary in addition to its usual procedures gets access to oracles that implement a family of independent random permutations (called the *ideal cipher*) and their inverses. For simplicity, we do not give oracle access to their inverses below since it won't matter for this homework.

For key-length $k \in \mathbb{N}$, block-length $\ell \in \mathbb{N}$, and an adversary A, define *ideal-cipher target-key* recovery game IC-TKR^A_{k,\ell} as follows:

Define the *IC-TKR-advantage* of A for key-length k and block-length ℓ as

$$\mathbf{Adv}_{k,\ell}^{\mathrm{ic-tkr}}(A) = \Pr\left[\mathrm{IC}\text{-}\mathrm{TKR}_{k,\ell}^{A} \text{ outputs } 1\right].$$

(20 points.) For $q \in \mathbb{N}$, define an appropriate notion of a q-query exhaustive key-search adversary A^q_{ic-eks} in the ideal cipher model.

(60 points.) Prove that

$$\mathbf{Adv}_{k,\ell}^{\mathrm{ic\text{-}tkr}}(A_{\mathrm{ic\text{-}eks}}^1) \geq 1 - \frac{2^k - 1}{2^{\ell+1}} \; .$$

Hint: First prove

$$\mathbf{Adv}_{k,\ell}^{\mathrm{ic-tkr}}(A_{\mathrm{ic-eks}}^{1}) \geq 2^{\ell-k} \left(1 - \left(1 - \frac{1}{2^{\ell}}\right)^{2^{k}}\right)$$

and use the inequality

$$(1-x)^n \leq \sum_{i=0}^m \binom{n}{i} (-x)^i$$

for any $x \in \mathbb{R}$ and $m, n \in \mathbb{Z}$ such that $0 \le x \le 1$ and $0 \le m \le n$ and m is even.

(20 points.) What does a result in the ideal cipher model "mean" in practice? Is DES an ideal cipher? Is AES an ideal cipher? Discuss.