

COMPSCI 690RA: Randomized Algorithms and Probabilistic Data Analysis

Prof. Cameron Musco

University of Massachusetts Amherst. Spring 2022.

Lecture 10

- Problem Set 3 is due this Friday 4/15 at 8pm.
- We have no class next Wednesday – it's a Monday at UMass.
- I will post a quiz due Tuesday 4/26 at 8pm.
- Remember that office hours are now Thursday at 4pm.

Summary

Last Week: Markov Chains.

- Finish spectral graph sparsification and physical interpretation
- Start on Markov chains and their analysis
- Markov chain based algorithms for satisfiability: $\approx n^2$ time for 2-SAT, and $\approx (4/3)^n$ for 3-SAT. 2^n

Summary

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- Finish spectral graph sparsification and physical interpretation
- Start on Markov chains and their analysis
- Markov chain based algorithms for satisfiability: $\approx n^2$ time for 2-SAT, and $\approx (4/3)^n$ for 3-SAT.

Today: Markov Chains Continued

- The gambler's ruin problem.
- Aperiodicity and stationary distribution of a Markov chain.
- Mixing time and its analysis via coupling.
- Markov Chain Monte Carlo (MCMC) methods.

Question 3

Not complete

Points out of 1.00

Flag question

 [Edit question](#)

Consider a matrix $A \in \mathbb{R}^{5 \times 3}$ such that $x = [0, 2, 2, 1, 1]$ is in the column span of A .

$$5 \begin{bmatrix} 3 \\ A \end{bmatrix} \begin{bmatrix} e \\ \end{bmatrix} = \begin{bmatrix} 0 \\ 2 \\ \vdots \end{bmatrix}$$

What can we say about the leverage score of the second row of A , i.e., τ_2 ?

- a. $\tau_2 \leq 2$
 b. $\tau_2 \leq 0.4$
 c. $\tau_2 \geq 2$
 d. Nothing without knowing A .
 e. $\tau_2 \geq 0.4$

Check

$$\tau_2 \in [0, 1]$$

$$\tau_2 \in [0.4, 1]$$

$$\tau_2 = \max_{y \in \mathbb{R}^3} \frac{\|Ay\|_2^2}{\|y\|_2^2}$$

$$\frac{x(2)^2}{\|x\|_2^2} = \frac{4}{10} = .4$$

$$5 \begin{bmatrix} \begin{bmatrix} 1 \\ 2 \\ 2 \\ 1 \\ 1 \end{bmatrix} \\ \vdots \end{bmatrix} \begin{bmatrix} \tau_1 \\ \tau_2 \\ \vdots \end{bmatrix} \quad \text{--- norm 1}$$

Question 4

Not complete

Points out of 1.00

Flag question

 [Edit question](#)

Consider a matrix $A \in \mathbb{R}^{n \times d}$ with full column rank. Let $U \in \mathbb{R}^{n \times d}$ be its left singular vector matrix. Let $Q \in \mathbb{R}^{n \times d}$ be an orthonormal basis for its column span, computed e.g., using Gram-Schmidt.

True or False: for all $i \in [n]$, $\|Q_{i,:}\|_2 = \|U_{i,:}\|_2$.

Select one:

 True

 False

$$Q_{i,:} = U_{i,:} R$$

$$Q = UR$$

$R \in \mathbb{R}^{d \times d}$ is orthonormal



$$\begin{bmatrix} u_1 & u_2 \end{bmatrix} \begin{bmatrix} r_{11} & r_{12} \\ r_{21} & r_{22} \end{bmatrix} = \begin{bmatrix} | & | \\ u_1 r_{11} + u_2 r_{21} & u_1 r_{12} + u_2 r_{22} \\ | & | \end{bmatrix} \begin{matrix} q_1 \\ Q \\ q_2 \end{matrix}$$

$$A = QR$$

$$A = U \Sigma V^T$$

Question 5

Not complete

Points out of 1.00

Flag question

 [Edit question](#)

Let E_0, E_1, \dots be independent, identically distributed random variables. Which of the following are Markov chains? Select all that apply.

a. X_0, X_1, \dots where $X_0 = E_0$ and $X_{i+1} = X_i + E_i$

b. E_0, E_1, \dots themselves.

c. X_0, X_1, \dots where $X_0 = E_0$ and $X_{i+1} = X_i \cdot E_i$

d. X_0, X_1, \dots where $X_0 = E_0$, and $X_{i+1} = X_i + E_{i-1} + E_i$

e. X_0, X_1, \dots where $X_0 = E_0, X_1 = E_1$, and $X_{i+1} = X_i + X_{i-1} + E_i$

$$\Pr(X_{i+1} = v | X_i) \neq \Pr(X_{i+1} = v | X_i, X_{i-1})$$

$$X_0 = E_0$$

$$X_1 = E_0 + E_1 + E_0$$

$$X_2 = X_1 + E_1 + E_0$$

$$X_2 = X_0 + X_1 + E_0$$

Markov Chain Review

- A discrete time stochastic process is a **Markov chain** if it is **memoryless**:

$$\Pr(\underline{X_t = a_t} | \underline{X_{t-1} = a_{t-1}, \dots, X_0 = a_0}) = \Pr(\underline{X_t = a_t} | \underline{X_{t-1} = a_{t-1}})$$

- If each X_t can take m possible values, the Markov chain is specified by the **transition matrix** $P \in [0, 1]^{m \times m}$ with

$$P_{i,j} = \Pr(\underline{X_{t+1} = j} | \underline{X_t = i}).$$

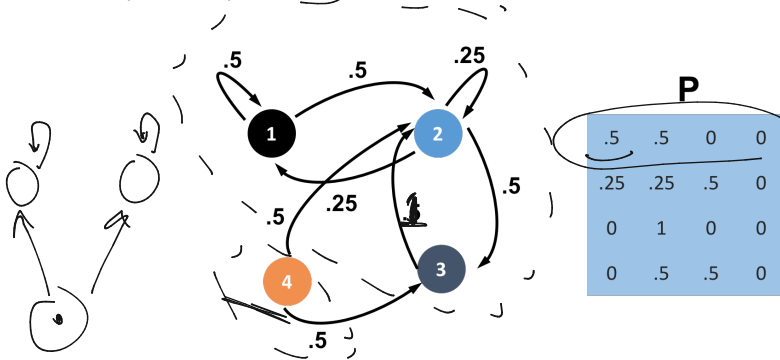
- Let $\underline{q_t} \in [0, 1]^{1 \times m}$ be the distribution of $\underline{X_t}$. Then $q_{t+1} = q_t P$.

q_1	P	q_2
<u>.5</u> <u>.5</u> 0 0	<u>.5</u> <u>.5</u> 0 0	<u>.375</u> <u>.375</u> .25 0
	.25 .25 .5 0	
	0 1 0 0	
	0 .5 .5 0	

(Note: A hand-drawn arrow points from the underlined elements of q_1 to the underlined elements of q_2 .)

Markov Chain Review

Often viewed as an underlying state transition graph. Nodes correspond to possible values that each X_t can take.



The Markov chain is irreducible if the underlying graph consists of single strongly connected component.

Gambler's Ruin

Gambler's Ruin

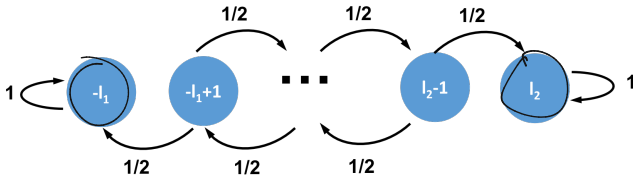


- You and 'a friend' repeatedly toss a fair coin. If it hits heads, you give your friend \$1. If it hits tails, they give you \$1.
- You start with $\$l_1$ and your friend starts with $\$l_2$. When either of you runs out of money the game terminates.
- What is the probability that you win $\$l_2$?

Gambler's Ruin Markov Chain

Let X_0, X_1, \dots be the Markov chain where X_t is your profit at step t . $X_0 = 0$ and:

$$\begin{aligned} P_{-l_1, -l_1} &= P_{l_2, l_2} = 1 \\ P_{i, i+1} &= P_{i, i-1} = 1/2 \text{ for } -l_1 < i < l_2 \end{aligned}$$

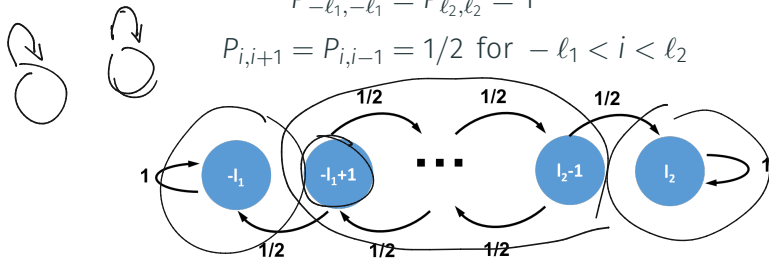


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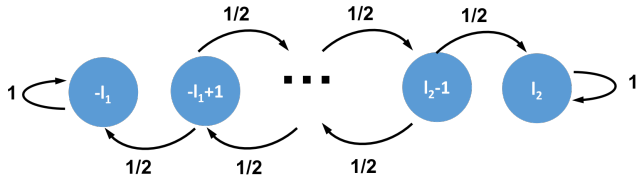
- $-l_1$ and l_2 are **absorbing states**.
- All i with $-l_1 < i < l_2$ are **transient states**. I.e., $\Pr[X_{t'} = i \text{ for some } t' > t \mid X_t = i] < 1$.

Gambler's Ruin Markov Chain

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 $\Pr[X_{t'} = i \text{ for some } t' > t \mid X_t = i] < 1$.

Observe that this Markov chain is also a **Martingale** since

$$\mathbb{E}[X_{t+1} | X_t] = X_t.$$

Gambler's Ruin Analysis

Let X_0, X_1, \dots be the Markov chain where X_t is your profit at step t .
 $X_0 = 0$ and:

$$P_{-l_1, -l_1} = P_{l_2, l_2} = 1$$
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We want to compute $q = \lim_{t \rightarrow \infty} \Pr[X_t = l_2]$.

Gambler's Ruin Analysis

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$$P_{i, i+1} = P_{i, i-1} = 1/2 \text{ for } -l_1 < i < l_2 \quad \mathbb{E}[X_0] = 0$$

We want to compute $q = \lim_{t \rightarrow \infty} \Pr[X_t = l_2]$.

By linearity of expectation, for any i , $\mathbb{E}[X_i] = 0$. Further, for
 $q = \lim_{t \rightarrow \infty} \Pr[X_t = l_2]$, since $-l_1, l_2$ are the only non-transient states,

$$\lim_{t \rightarrow \infty} \Pr[X_t = i] = 0 \quad \text{if } i \neq -l_1, l_2$$
$$\lim_{t \rightarrow \infty} \mathbb{E}[X_t] = \sum_i \lim_{t \rightarrow \infty} \Pr(X_t = i) \cdot i$$
$$\lim_{t \rightarrow \infty} \mathbb{E}[X_t] = \underbrace{l_2 q}_{\text{from } l_2} + \underbrace{-l_1(1-q)}_{\text{from } -l_1} = 0$$

Gambler's Ruin Analysis

Let X_0, X_1, \dots be the Markov chain where X_t is your profit at step t .
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$$\lim_{t \rightarrow \infty} \mathbb{E}[X_t] = \underbrace{l_2 q + (-l_1)(1 - q)}_{= 0} = 0.$$

Solving for q , we have $q = \frac{l_1}{l_1 + l_2}$.

$$q l_2 = q, (1 - q)$$
$$q(l_1 + l_2) = l_1$$

Gambler's Ruin Thought Exercise

What if you always walk away as soon as you win just \$1. Then what is your probability of winning, and what are your expected winnings?

$$\frac{q_1}{q_1 + q_2} \quad q_2^{-1} \quad \left(\frac{q_1}{q_1 + 1} \right)^0 \cdot 1 + \left(-\frac{q_1}{q_1 + 1} \right)^0 \cdot -q_1$$
$$= 0$$

$$\mathbb{E}[X_i] = 0$$

optional stopping theorem.

Stationary Distributions

Stationary Distribution

A **stationary distribution** of a Markov chain with transition matrix $P \in [0, 1]^{m \times m}$ is a distribution $\pi \in [0, 1]^m$ such that $\pi = \pi P$.

i.e. if $X_t \sim \pi$, then $X_{t+1} \sim \pi P = \pi$.

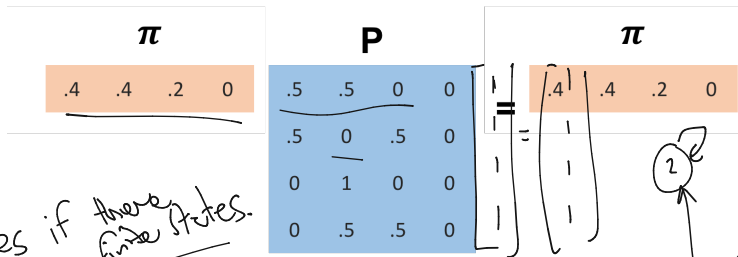
π		P	π
.4 .4 .2 0		.5 .5 0 0 .5 0 .5 0 0 1 0 0 0 .5 .5 0	.4 .4 .2 0

The diagram illustrates the equation $\pi = \pi P$. On the left, a row vector π is shown with values .4, .4, .2, and 0. In the middle, a 4x4 transition matrix P is shown with values: row 1: .5, .5, 0, 0; row 2: .5, 0, .5, 0; row 3: 0, 1, 0, 0; row 4: 0, .5, .5, 0. On the right, the resulting row vector π is shown with values .4, .4, .2, and 0. Hand-drawn arrows and brackets highlight the calculation of the first element of the resulting π vector, which is $.4 \cdot .5 + .4 \cdot .5 + .2 \cdot 0 + 0 \cdot 0 = .2 + .2 = .4$.

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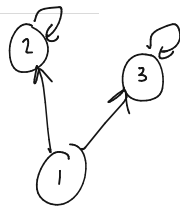
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Yes if there are finite states.

Think-pair-share: Do all Markov chains have a stationary distribution?

$X_{i+1} = X_i \pm 1$ — infinite # states



$\begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$

Claim (Existence of Stationary Distribution)

Any Markov chain with a finite state space, and transition matrix $P \in [0, 1]^{m \times m}$ has a stationary distribution $\pi \in [0, 1]^m$ with $\pi = \pi P$.

$$f(\pi) = \pi P = \pi$$

Perron-Frobenius theorem

Claim (Existence of Stationary Distribution)

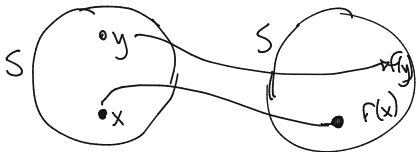
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$$\underline{q}P = \underline{q}'$$

$$P: S \rightarrow S$$

S is the m -dim simplex

Follows from the **Brouwer fixed point theorem**: for any continuous function $f: \underline{S} \rightarrow \underline{S}$, where S is a compact convex set, there is some x such that $f(x) = x$.



$$f(x) = x+1$$

$$\begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$\rightarrow \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

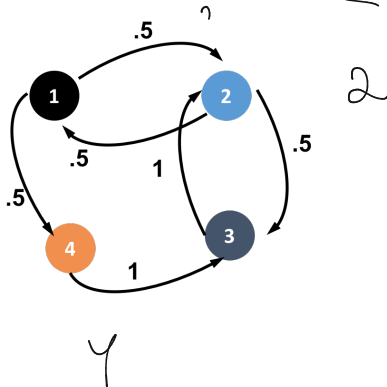
$$\begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

Periodicity

The **periodicity** of a state i is defined as:

$$T = \gcd\{t > 0 : \Pr(X_t = i \mid X_0 = i) > 0\}.$$

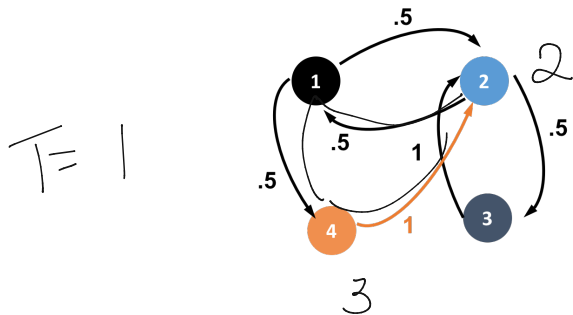
$$T = 2$$



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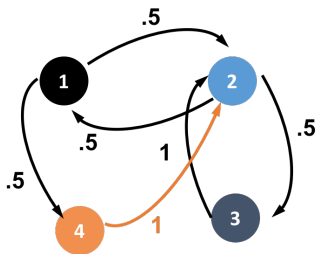
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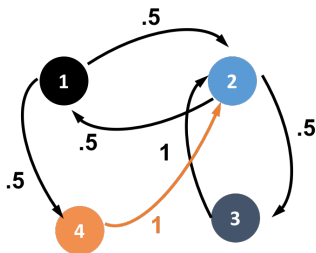


The state is **aperiodic** if it has periodicity $T = 1$.

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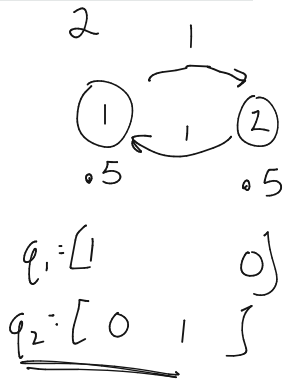
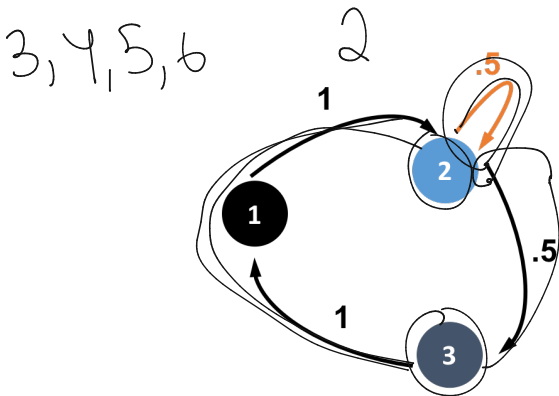
The state is **aperiodic** if it has periodicity $T = 1$.

A Markov chain is aperiodic if all states are aperiodic.

Periodicity

Claim

If a Markov chain is irreducible, and has at least one self-loop, then it is aperiodic.



Fundamental Theorem

Theorem (The Fundamental Theorem of Markov Chains)

Let X_0, X_1, \dots be a Markov chain with a finite state space and transition matrix $P \in [0, 1]^{m \times m}$. If the chain is both irreducible and aperiodic,

1. There exists a **unique** stationary distribution $\pi \in [0, 1]^m$ with $\pi = \pi P$.
2. For any states i, j , $\lim_{t \rightarrow \infty} \Pr[X_t = i | X_0 = j] = \underline{\pi(i)}$. I.e., for any initial distribution q_0 , $\lim_{t \rightarrow \infty} q_t = \lim_{t \rightarrow \infty} \underline{q_0 P^t} = \underline{\pi}$.
3. $\underline{\pi(i)} = \frac{1}{\mathbb{E}[\min\{t: X_t = i\} | X_0 = i]}$. I.e., $\pi(i)$ is the inverse of the average expected return time from state i back to i .

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$$, \sum \pi(i) = 1$$

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3. $\pi(i) = \frac{1}{\mathbb{E}[\min\{t: X_t = i\} | X_0 = i]}$. I.e., $\pi(i)$ is the inverse of the average expected return time from state i back to i .

In the limit, the probability of being at any state i is **independent of the starting state**.

Stationary Distribution Example 1

Shuffling Markov Chain: Given a pack of c cards. At each step draw a random card, place it on top, and repeat.

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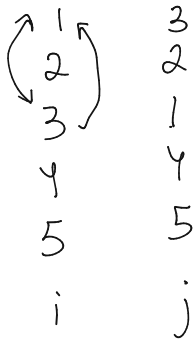
- What is the state space of this chain?

$c!$ all possible permutations of c distinct cards

Stationary Distribution Example 1

Shuffling Markov Chain: Given a pack of c cards. At each step draw ^{two} random cards, ~~place it on top, and repeat.~~ ^{swap them} swap with the top card.

- What is the state space of this chain? - all permutations, $c!$
- What is the transition probability P_{ij} ? How does it compare to P_{ji} ?



$$P_{ij} = P_{ji}$$

This slide has many errors

$$P_{ii} = 1/c$$

$$P_{ij} = 1/c$$

$$P_{ji} = 1/c$$

$$P_{ij} = 0$$

$$P_{ji} = 0$$

Stationary Distribution Example 1

Shuffling Markov Chain: Given a pack of c cards. At each step draw a random card, place it on top, and repeat.

- What is the state space of this chain?
- What is the transition probability $P_{i,j}$? How does it compare to $P_{j,i}$?

$$P_{ij} = P_{ji}$$

- This Markov chain is **symmetric** and thus its stationary distribution is uniform, $\pi(i) = \frac{1}{c}$.

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Letting $m = c!$ denote the size of the state space,

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$$\begin{bmatrix} \pi \end{bmatrix} P = \begin{bmatrix} \pi \\ \pi \\ \vdots \\ \pi \end{bmatrix}$$

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$$\left[\begin{array}{c} \pi \\ \pi \\ \pi \end{array} \right]$$

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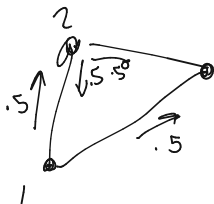
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Once we have exhibited a stationary distribution, we know that it is unique and that the chain converges to it in the limit!

Stationary Distribution Example 2

Random Walk on an Undirected Graph: Consider a random walk on an undirected graph. If it is at node i at step t , then it moves to any of i 's neighbors at step $t + 1$ with probability $\frac{1}{d_i}$.

- What is the state space of this chain? ✓
- What is the transition probability $P_{i,j}$?

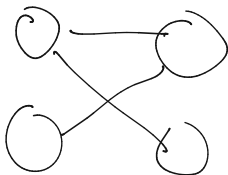


$$\frac{1}{d_i}$$

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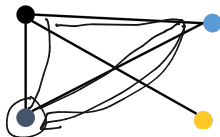
- What is the state space of this chain?
- What is the transition probability $P_{i,j}$?
- Is this chain aperiodic?



Stationary Distribution Example 2

Random Walk on an Undirected Graph: Consider a random walk on an undirected graph. If it is at node i at step t , then it moves to any of i 's neighbors at step $t + 1$ with probability $\frac{1}{d_i}$.

- What is the state space of this chain?
- What is the transition probability $P_{i,j}$?
- Is this chain aperiodic?
- If the graph is not bipartite, then there is at least one odd cycle, making the chain aperiodic.



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$$\pi P_{:,i} = \sum_j \pi(j) P_{j,i} = \frac{1}{d_j} \pi(j) \frac{d_j}{2|E|}$$

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$\frac{1}{d_j}$ if $j \in N(i)$
0 if $j \notin N(i)$

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I.e., the probability of being at a given node i is dependent only on the node's degree, not on the structure of the graph in any other way.

Mixing Times

Total Variation Distance

[wasserstein]

Definition (Total Variation (TV) Distance)

For two distributions $p, q \in [0, 1]^m$ over state space $[m]$, the total variation distance is given by:

$$\|p - q\|_{TV} = \frac{1}{2} \sum_{i \in [m]} |p(i) - q(i)| = \max_{A \subseteq [m]} |p(A) - q(A)|.$$

p

$$\begin{bmatrix} \frac{1}{6} \\ \frac{1}{6} \\ \frac{1}{6} \\ \vdots \\ \frac{1}{6} \end{bmatrix}$$

$$\begin{bmatrix} \frac{1}{2} \\ \frac{1}{2} \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$\left(\frac{1}{2} + \frac{1}{2}\right) - \frac{1}{6} - \frac{1}{6} = \frac{2}{3}$$

$$\frac{1}{2} \left(4 \cdot \left|\frac{1}{6}\right| + 2 \cdot \left|\frac{1}{2} - \frac{1}{6}\right| \right) = \frac{2}{3}$$

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Kantorovich-Rubinstein duality: Let P, Q be possibly correlated random variables with marginal distributions p, q . Then

$$p = \begin{bmatrix} \frac{1}{2} \\ \frac{1}{2} \end{bmatrix}$$

$$q = \begin{bmatrix} \frac{2}{3} \\ \frac{1}{3} \end{bmatrix}$$

$$\|p - q\|_{TV} \leq \Pr[P \neq Q].$$

$$P = Q$$

$$\|p - q\|_{TV} = \min_{P, Q} \Pr[P \neq Q]$$

$$P = Q \text{ so } \|p - q\|_{TV} = 0$$

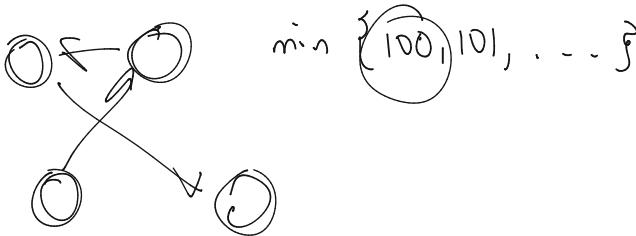
coupling

Mixing Time

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Consider a Markov chain X_0, X_1, \dots with unique stationary distribution π . Let $q_{i,t}$ be the distribution over states at time t assuming $X_0 = i$. The mixing time is defined as:

$$\tau(\epsilon) = \min \left\{ t : \max_{i \in [m]} \|q_{i,t} - \pi\|_{TV} \leq \epsilon \right\}.$$



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Claim: If X_0, X_1, \dots is finite, irreducible, and aperiodic, then $\tau(\epsilon) \leq \tau(1/2) \cdot c \log(1/\epsilon)$ for large enough constant c .

$$q_0 P^t = q_t \approx \pi$$

$$\begin{aligned} q_t &= \pi + e \\ q_{2t} &= \pi P^t + e P^t \end{aligned}$$

Coupling Motivation

Claim: $\max_{i \in [m]} \|q_{i,t} - \pi\|_{TV} \leq \max_{i,j \in [m]} \|q_{i,t} - q_{j,t}\|_{TV}$.

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i) Starting
States of
Markov chain.

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$$\begin{aligned}\|q_{i,t} - \pi\|_{TV} &= \|q_{i,t} - \pi P^t\|_{TV} \quad \underbrace{q_{j,t}} \\ &= \|q_{i,t} - \sum_j \pi(j) e_j P^t\|_{TV}\end{aligned}$$

$$\pi = \pi(j) e_j \quad e_j = [0 \ 0 \ 1 \ 0 \ 0]$$

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Coupling: A common technique for bounding the mixing time by showing that $\max_{i,j \in [m]} \|q_{i,t} - q_{j,t}\|_{TV}$ is small.