# COMPSCI 514: Algorithms for Data Science

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University of Massachusetts Amherst. Fall 2023. Lecture 24 (Final Lecture!)

- Problem Set 5 can be submitted up to 12/11 (Monday) at 11:59pm.
- Exam is next Thursday 12/14, from 10:30am-12:30pm in class.
- I am holding office hours Tuesday 12/12 1-3:30pm and Wednesday 12/13 2pm-3pm Both will be held in C<u>S140</u>.
- It would be really helpful if you could fill out SRTIs for this class.
- There is no quiz due this week.

Summary

#### Last Class:

• Analysis of gradient descent for convex and Lipschitz functions.

This Class:

• Extend gradient descent to constrained optimization via projected gradient descent.

• Course wrap up and review.

### GD Analysis Proof

A Theorem – GD on Convex Lipschitz Functions: For convex <u>G</u>-Lipschitz function f, GD run with  $t ≥ \frac{R^2G^2}{\epsilon^2}$  iterations,  $\eta = \frac{R}{G\sqrt{t}}$ , and starting point within radius R of  $\vec{\theta}_*$ , outputs  $\hat{\theta}$  satisfying:

$$\underline{f(\hat{\theta})} \leq f(\vec{\theta}_*) + \epsilon.$$

Step 1: For all 
$$i, f(\vec{\theta}_i) - f(\vec{\theta}_*) \leq \frac{\|\vec{\theta}_i - \vec{\theta}_*\|_2^2 - \|\vec{\theta}_{i+1} - \vec{\theta}_*\|_2^2}{\sqrt{2\eta}} + \frac{\eta G^2}{2} + \frac{\eta G^2}{2}$$

#### **GD** Analysis Proof

**Theorem – GD on Convex Lipschitz Functions:** For convex *G*-Lipschitz function *f*, GD run with  $t \ge \frac{R^2 G^2}{\epsilon^2}$  iterations,  $\eta = \frac{R}{G\sqrt{t}}$ , and starting point within radius *R* of  $\vec{\theta}_*$ , outputs  $\hat{\theta}$  satisfying:

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### **Constrained Convex Optimization**

Often want to perform convex optimization with convex constraints.

$$\underbrace{\vec{\theta}^*}_{\vec{\theta} \in \mathcal{S}} = \arg\min_{\vec{\theta} \in \mathcal{S}} f(\underline{\vec{\theta}}),$$

where  $\mathcal{S}$  is a convex set.

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where  $\mathcal{S}$  is a convex set.

**Definition – Convex Set:** A set  $S \subseteq \mathbb{R}^d$  is convex if and only if, for any  $\vec{\theta_1}, \vec{\theta_2} \in S$  and  $\lambda \in [0, 1]$ :

 $(1-\lambda)\vec{\theta_1} + \lambda\cdot\vec{\theta_2}\in\mathcal{S}$ 



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E.g. 
$$S = \{ \vec{\theta} \in \mathbb{R}^{d} : \|\vec{\theta}\|_{2} \leq 1 \}$$
.  $\rightarrow \{ \vec{\theta} \in \mathbb{R}^{d} : \|\vec{\theta}\|_{2} \leq 1 \}$ .  
 $O_{1, j} O_{2} \in S$   
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- For  $S = \{\vec{\theta} \in \mathbb{R}^d : \|\vec{\theta}\|_2 \le 1\}$  what is  $P_S(\vec{y})$ ?
- For S being a k dimensional subspace of  $\mathbb{R}^d$ , what is  $P_{\mathcal{S}}(\vec{y})$ ?

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#### **Projected Gradient Descent**

- Choose some initialization  $\vec{\theta}_1$  and set  $\eta = \frac{R}{G\sqrt{t}}$ .
- For i = 1, ..., t 1•  $\vec{\theta}_{i+1}^{(out)} = \vec{\theta}_i - \eta \cdot \vec{\nabla} f(\vec{\theta}_i)$  $\underbrace{ \cdot \vec{\theta}_{i+1}}_{i+1} = P_{\mathcal{S}}(\vec{\theta}_{i+1}^{(out)}).$
- Return  $\hat{\theta} = \arg \min_{\vec{\theta}_i} f(\vec{\theta}_i)$ .

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**Theorem – Projection to a convex set:** For any convex set  $S \subseteq \mathbb{R}^d$ ,  $\vec{y} \in \mathbb{R}^d$ , and  $\underline{\vec{\theta} \in S}$ ,

$$\|\mathsf{P}_{\mathcal{S}}(\vec{y}) - \vec{\theta}\|_2 \leq \|\vec{y} - \vec{\theta}\|_2.$$



$$S = s + b \le p = 2$$
  

$$|| b - W^{T} \times || \le || b - \chi ||$$

$$|| b - W^{T} \times || + || (t - W^{T}) \times ||)$$

**Theorem – Projected GD:** For convex *G*-Lipschitz function *f*, and convex set S, Projected GD run with  $t \ge \frac{R^2 G^2}{\epsilon^2}$  iterations,  $\eta = \frac{R}{G\sqrt{t}}$ , and starting point within radius *R* of  $\vec{\theta}_*$ , outputs  $\hat{\theta}$  satisfying:

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**Recall:** 
$$\vec{\theta}_{i+1}^{(out)} = \vec{\theta}_i - \eta \cdot \vec{\nabla} f(\vec{\theta}_i)$$
 and  $\vec{\theta}_{i+1} = P_{\mathcal{S}}(\vec{\theta}_{i+1}^{(out)})$ .

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Step 1: For all  $i, f(\vec{\theta}_i) - f(\vec{\theta}_*) \leq \frac{\|\vec{\theta}_i - \theta_*\|_2^2 - \|\vec{\theta}_{i+1}^{(out)} - \vec{\theta}_*\|_2^2}{2\eta} + \frac{\eta G^2}{2}.$ 

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Recall:  $\vec{\theta}_{i+1}^{(out)} = \vec{\theta}_i - \eta \cdot \vec{\nabla} f(\vec{\theta}_i)$  and  $\vec{\theta}_{i+1} = P_S(\vec{\theta}_{i+1}^{(out)})$ . Step 1: For all  $i, f(\vec{\theta}_i) - f(\vec{\theta}_*) \le \frac{\|\vec{\theta}_i - \theta_*\|_2^2 - \|\vec{\theta}_{i+1}^{(out)} - \vec{\theta}_*\|_2^2}{2\eta} + \frac{\eta G^2}{2}$ . Step 1.a: For all  $i, f(\vec{\theta}_i) - f(\vec{\theta}_*) \le \frac{\|\vec{\theta}_i - \vec{\theta}_*\|_2^2 - \|\vec{\theta}_{i+1} - \vec{\theta}_*\|_2^2}{2\eta} + \frac{\eta G^2}{2}$ . Step 2:  $\frac{1}{t} \sum_{i=1}^t f(\vec{\theta}_i) - f(\vec{\theta}_*) \le \frac{R^2}{2\eta \cdot t} + \frac{\eta G^2}{2} \implies$  Theorem.  $\le \varepsilon$ 

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#### Randomization as a computational resource for massive datasets.

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In the process covered probability/statistics tools that are very useful beyond algorithm design: concentration inequalities, higher moment bounds, law of large numbers, central limit theorem, linearity of expectation and variance, union bound, median as a robust estimator.

#### Methods for working with (compressing) high-dimensional data

• Started with randomized dimensionality reduction and the JL lemma: compression from *any* d-dimensions to  $O(\log n/\epsilon^2)$  dimensions while preserving pairwise distances.

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- Low-rank approximation of similarity matrices and entity embeddings (e.g., LSA, word2vec, DeepWalk).
- Spectral graph theory nonlinear dimension reduction and <u>spe</u>ctral clustering for community detection.
- In the process covered linear algebraic tools that are very broadly useful in ML and data science: eigendecomposition, singular value decomposition, projection, norm transformations.

#### Foundations of continuous optimization and gradient descent.

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- How to analyze gradient descent in a simple setting (convex Lipschitz functions).
- Simple extension to projected gradient descent for optimization over a convex constraint set.
- Lots that we didn't cover: online and stochastic gradient descent, accelerated methods, adaptive methods, second order methods (quasi-Newton methods), practical considerations.
   Gave mathematical tools to understand these methods.

Thanks for a great semester!

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( ) X $f(x) = x^{T}c \quad \forall f(x) = c^{T}$  $f(x) = \sum_{i=1}^{d} x(i) c(i)$  $\frac{\partial f(x)}{\partial x_{i}} = c(i)$  $\begin{bmatrix} \partial f x \\ \partial X_1 \\ \vdots \\ \partial X_2 \\ \vdots \\ \partial f x \\ \vdots \\ \partial f x \\ \partial f x$ 

 $f(x) = X^{T} A X$  $\langle x, A X \rangle$  $\nabla f(x) = 2Ax$ X·A·X x<sup>2</sup> A  $\frac{d}{\sum_{i=1}^{n} \chi(i)} \stackrel{\text{d}}{\underset{j=1}{\overset{j=1}{\overset{j=$ 7 x A  $\frac{\sum_{j \neq i} x_{(i)} \times (j) A_{ij} + x_{(i)} M_{ji}}{+ x_{(i)} A_{ij} + x_{(i)} A_{ij}}$  $E A J \left[ x \right] = \left[ 0 \right] \frac{\partial F(x)}{\partial x(i)}$  $z = 2x_j A_{ij} + 2x_{i}$  $j \neq j$  $z = 2[A \times ](i)$  16

