# COMPSCI 514: Algorithms for Data Science

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- We released Problem Set 3 last night. It is due 11/17 at 11:59pm.
- Doing the first two Core Competency questions early might be helpful if you need linear algebra review.

#### Summary

#### Last Class:

- No-distortion embeddings for data lying in a k-dimensional subspace via an orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$  for that subspace.
- View as low-rank matrix factorization. Introduce concept of low-rank approximation.
- Idea of approximating a data matrix **X** with **XVV**<sup>T</sup> when the data points lie close to the subspace spanned by **V**'s columns.

#### This Class:

- 'Dual view' of low-rank approximation: data points that can be approximately reconstructed from a few basis vectors vs. linearly dependent features.
- How to find an optimal orthogonal basis  $V \in \mathbb{R}^{d \times k}$  to minimize  $\|X XVV^T\|_F^2$ .

### Low-Rank Factorization

**Claim:** If  $\vec{x}_1, \ldots, \vec{x}_n$  lie in a *k*-dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be written as

 $\mathbf{X} = \mathbf{X}\mathbf{V}\mathbf{V}^{\mathsf{T}}$  (implies rank( $\mathbf{X}$ )  $\leq k$ )

•  $VV^T$  is a projection matrix, which projects the rows of X (the data points  $\vec{x}_1, \ldots, \vec{x}_n$  onto the subspace  $\mathcal{V}$ .



#### Low-Rank Approximation

**Claim:** If  $\vec{x}_1, \ldots, \vec{x}_n$  lie close to a *k*-dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be approximated as:

 $\mathbf{X} \approx \mathbf{X} \mathbf{V} \mathbf{V}^{\mathsf{T}}$ 



**Note: XVV**<sup>*T*</sup> has rank *k*. It is a low-rank approximation of **X**.

$$XVV^{\mathsf{T}} = \underset{\mathsf{B with rows in }\mathcal{V}}{\arg\min} \|\mathsf{X} - \mathsf{B}\|_{F}^{2} = \sum_{i,j} (\mathsf{X}_{i,j} - \mathsf{B}_{i,j})^{2}.$$

**So Far:** If  $\vec{x_1}, \ldots, \vec{x_n}$  lie close to a *k*-dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be approximated as:

#### $\mathbf{X} \approx \mathbf{X} \mathbf{V} \mathbf{V}^T.$

This is the closest approximation to X with rows in  ${\cal V}$  (i.e., in the column span of V).

- Letting  $\mathbf{V}\mathbf{V}^T \vec{x}_i$ ,  $\mathbf{V}\mathbf{V}^T \vec{x}_j$  be the *i*<sup>th</sup> and *j*<sup>th</sup> projected data points,  $\|\mathbf{V}\mathbf{V}^T \vec{x}_i - \mathbf{V}\mathbf{V}^T \vec{x}_j\|_2 = \|\mathbf{V}^T \mathbf{V}\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \mathbf{V}\mathbf{V}^T \vec{x}_j\|_2 = \|\mathbf{V}^T \vec{x}_i - \mathbf{V}^T \vec{x}_j\|_2$ .
- I.e., we can use the rows of  $XV \in \mathbb{R}^{n \times k}$  as a compressed approximate data set.

Key question is how to find the subspace  ${\mathcal V}$  and correspondingly  ${\textbf V}.$ 

**Quick Exercise 1:** Show that  $VV^T$  is idempotent. I.e.,  $(VV^T)(VV^T)\vec{y} = (VV^T)\vec{y}$  for any  $\vec{y} \in \mathbb{R}^d$ .

Quick Exercise 2: Show that  $VV^{T}(I - VV^{T}) = 0$  (the projection is orthogonal to its complement).

#### Pythagorean Theorem

# **Pythagorean Theorem:** For any orthonormal $\mathbf{V} \in \mathbb{R}^{d \times k}$ and any $\vec{y} \in \mathbb{R}^d$ ,

$$\|\vec{y}\|_{2}^{2} = \|(\mathbf{V}\mathbf{V}^{T})\vec{y}\|_{2}^{2} + \|\vec{y} - (\mathbf{V}\mathbf{V}^{T})\vec{y}\|_{2}^{2}.$$

# A Step Back: Why Low-Rank Approximation?

**Question:** Why might we expect  $\vec{x_1}, \ldots, \vec{x_n} \in \mathbb{R}^d$  to lie close to a *k*-dimensional subspace?

• The rows of X can be approximately reconstructed from a basis of *k* vectors.



orthonormal basis  $v_1, \dots, v_{15}$ 



# Dual View of Low-Rank Approximation

**Question:** Why might we expect  $\vec{x_1}, \ldots, \vec{x_n} \in \mathbb{R}^d$  to lie close to a *k*-dimensional subspace?

• Equivalently, the columns of **X** are approx. spanned by *k* vectors.

Linearly Dependent Variables:

	bedrooms	bathrooms	sq.ft.	floors	list price	sale price		bedrooms
home 1	2	2	1800	2	200,000	195,000	home 1	2
home 2	4	2.5	2700	1	300,000	310,000	home 2	4
		•	•	•	•	•		
•	•	•	•	•	•	•		•
•	•	•	•	•	•	•	•	•
home n	5	3.5	3600	3	450,000	450,000	home n	5 <sup>10</sup>

### Best Fit Subspace

If  $\vec{x}_1, \ldots, \vec{x}_n$  are close to a *k*-dimensional subspace  $\mathcal{V}$  with orthonormal basis  $\mathbf{V} \in \mathbb{R}^{d \times k}$ , the data matrix can be approximated as  $\mathbf{XVV}^T$ . **XV** gives optimal embedding of **X** in  $\mathcal{V}$ .

How do we find  $\mathcal{V}$  (equivilantly V)?



# Solution via Eigendecomposition

**V** minimizing  $\|\mathbf{X} - \mathbf{X}\mathbf{V}\mathbf{V}^T\|_F^2$  is given by:

$$\underset{\text{orthonormal } \mathbf{V} \in \mathbb{R}^{d \times k}}{\arg \max} \|\mathbf{X}\mathbf{V}\|_{F}^{2} = \sum_{i=1}^{n} \|\mathbf{V}^{\mathsf{T}} \vec{x}_{i}\|_{2}^{2} = \sum_{j=1}^{k} \|\mathbf{X} \vec{v}_{j}\|_{2}^{2}$$

Surprisingly, can find the columns of V,  $\vec{v}_1, \ldots, \vec{v}_k$  greedily.

$$\vec{v}_1 = \underset{\vec{v} \text{ with } \|v\|_2 = 1}{\arg \max} \|\mathbf{X} \vec{v}\|_2^2 \vec{v}^T \mathbf{X}^T \mathbf{X} \vec{v}.$$

$$\vec{v}_2 = \underset{\vec{v} \text{ with } \|v\|_2=1, \ \langle \vec{v}, \vec{v}_1 \rangle = 0}{\arg \max} \vec{v}^T \mathbf{X}^T \mathbf{X} \vec{v}.$$

$$\vec{V}_k = \arg \max_{\vec{v} \text{ with } \|v\|_2 = 1, \ \langle \vec{v}, \vec{v}_j \rangle = 0 \ \forall j < k} \vec{V}^T \mathbf{X}^T \mathbf{X}^{\vec{v}}$$

 $\vec{v}_1, \ldots, \vec{v}_k$  are the top k eigenvectors of  $\mathbf{X}^T \mathbf{X}$  by the Courant-Fischer Principle.