COMPSCI 514: ALGORITHMS FOR DATA SCIENCE

Cameron Musco University of Massachusetts Amherst. Fall 2019. Lecture 22

LOGISTICS

- Problem Set 4 released last night. Due Sunday 12/15 at 8pm.
- Final Exam Thursday 12/19 at 10:30am in Thompson 104.
- Exam prep materials (list of topics, practice problems) coming in next couple of days.

SUMMARY

Before Break:

- · Finished discussion of SGD.
- Gradient descent and SGD as applied to least squares regression.

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This Class:

- A quick tour of the counterintuitive properties of high-dimensional space.
- · Many connections to concentration inequalities.
- Implications for working with high-dimensional data (curse of dimensionality).

HIGH-DIMENSIONAL DATA

Modern data analysis often involves very high-dimensional data points.

- Websites record (tens of) thousands of measurements per user: who they follow, when they visit the site, timestamps for specific iteractions, etc.
- A 3 minute, 500 \times 500 pixel video clip at 15 FPS has \geq 2 billion pixel values.
- The human genome has 3 billion+ base pairs.

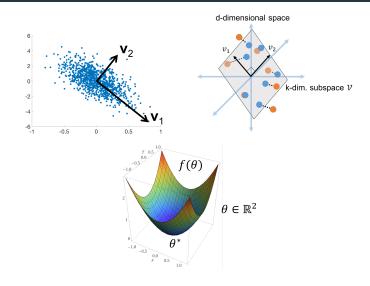
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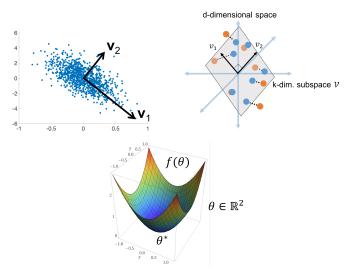
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Typically when discussing algorithm design we imagine data in much lower (usually 3) dimensional space.

LOW-DIMENSIONAL INTUITION



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This can be a bit dangerous as in reality high-dimensional space is very different from low-dimensional space.

ORTHOGONAL VECTORS

What is the largest set of mutually orthogonal unit vectors in *d*-dimensional space?



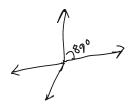
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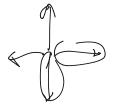
What is the largest set of mutually orthogonal unit vectors in *d*-dimensional space? Answer: *d*.





What is the largest set of unit vectors in d-dimensional space that have all pairwise dot products $|\langle \vec{x}, \vec{y} \rangle| \le \epsilon$? (think $\epsilon = .01$)





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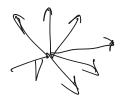
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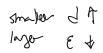
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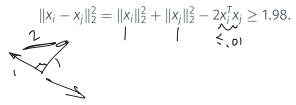
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- If we chose $t=\frac{1}{2}e^{\epsilon^2d/6}$ using a union bound over all $\leq t^2=\frac{1}{4}e^{\epsilon^2d/3}$ possible pairs, with probability > 1/2 all with be nearly orthogonal.

+= = 0,01.3/6

$$||x_i - x_j||_2^2$$

$$||x_i - x_j||_2^2 = ||x_i||_2^2 + ||x_j||_2^2 - \frac{2x_i^T x_j}{\xi}$$



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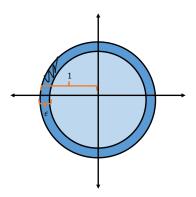
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· Only hope is if we have strong low-dimensional structure.

Let \mathcal{B}_d be the unit ball in d dimensions. $\mathcal{B}_d = \{x \in \mathbb{R}^d : ||x||_2 \le 1\}$.

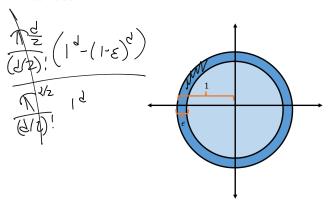
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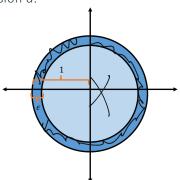
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What percentage of the volume of \mathcal{B}_d falls within ϵ distance of its surface? Answer: all but a $(1 - \epsilon)^d \le e^{-\epsilon d}$ fraction. Exponentially small in the dimension d!

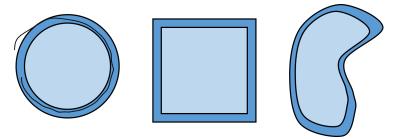


Volume of a radius R ball is $\frac{\pi^{\frac{d}{2}}}{(d/2)!} \cdot R^d$.

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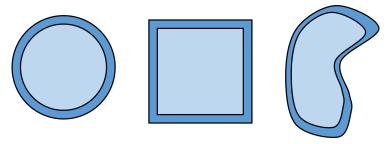
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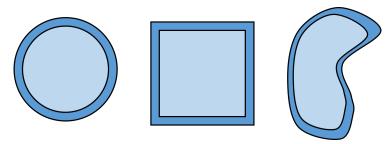
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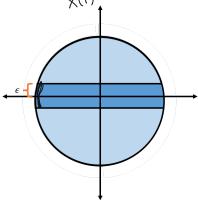
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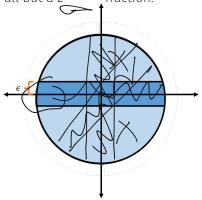
- If we randomly sample points from any high-dimensional shape, nearly all will fall near its surface.
- · 'All points are outliers.'

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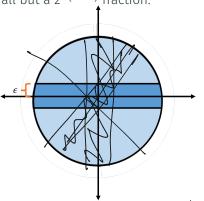
Formally: volume of set $S = \{x \in \mathcal{B}_d : |x(1)| \le \epsilon\}.$

What percentage of the volume of \mathcal{B}_d falls within ϵ distance of its equator? Answer: all but a $2^{\Theta(-\epsilon^2 d)}$ fraction.



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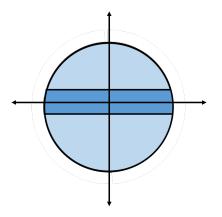
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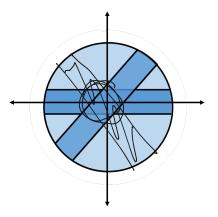
Formally: volume of set $S = \{x \in \mathcal{B}_d : |x(1)| \le \epsilon\}$. Let $\{x \in \mathcal{B}_d : |x(1)| \le \epsilon\}$. By symmetry, all but a $2^{\Theta(-\epsilon^2 d)}$ fraction of the volume falls within ϵ of any equator! $S = \{x \in \mathcal{B}_d : |\langle x, t \rangle| \le \epsilon\}$

Claim 1: All but a $2^{\Theta(-\epsilon^2 d)}$ fraction of the volume of a ball falls within ϵ of any equator.

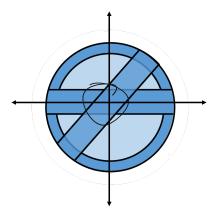
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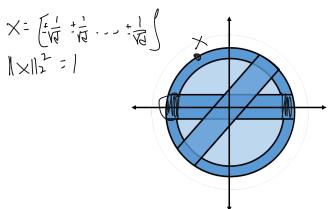


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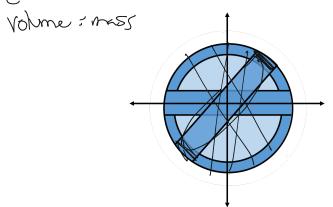
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How is this possible?

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How is this possible? High-dimensional space looks nothing like this picture!

Claim: All but a $2^{\Theta(-\epsilon^2 d)}$ fraction of the volume of a ball falls within ϵ of its equator. I.e., in $S = \{x \in \mathcal{B}_d : |x(1)| \le \epsilon\}$.

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Proof Sketch:

• Let x have entries set to independent Gaussians $\mathcal{N}(0,1)$ and let $\bar{x} = \frac{x}{\|x\|_2}$. \bar{x} is selected uniformly at random from the surface of the ball.

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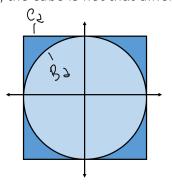
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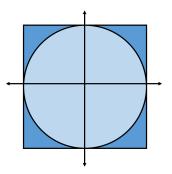
$$\begin{aligned} \Pr[|\bar{X}(1)| > \epsilon] &= \Pr[|X(1)| > \epsilon \cdot ||X||_2] \\ &\leq \Pr[|X(1)| > \epsilon \cdot \sqrt{d/2}] = 2^{\Theta(-(\epsilon \sqrt{d/2})^2)} = 2^{\Theta(-\epsilon^2 d)}. \end{aligned}$$

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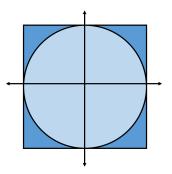


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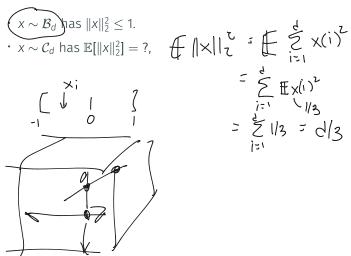


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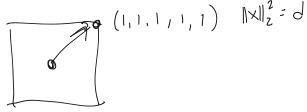
But volume of C_d is 2^d while volume of B^d is $\frac{\pi^{\frac{d}{2}}}{(d/2)!} = \frac{1}{d^{\Theta(d)}}$. A huge gap! So something is very different about these shapes...



- $x \sim \mathcal{B}_d$ has $||x||_2^2 \le 1$.
- $x \sim C_d$ has $\mathbb{E}[\|x\|_2^2] = d/3$,

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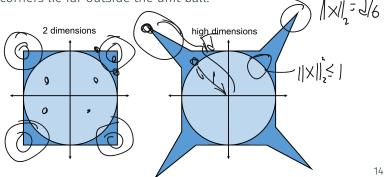
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But these different dimensional spaces have very different geometries, so how is this possible?

 x_1, \ldots, x_n are sampled from the surface of \mathcal{B}_d and $\mathbf{\Pi} x_1, \ldots, \mathbf{\Pi} x_n$ are (approximately) sampled from the surface of \mathcal{B}_m .

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- m is chosen just large enough so that the odd geometry of d-dimensional space will still hold on the n points in question.

Questions?