

## CS 312: Algorithms

### Flow Applications

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## First Application of Network Flows: Bipartite Matching

- ▶ Given a bipartite graph  $G = (L \cup R, E)$ , a subset of edges  $M \subseteq E \subseteq L \times R$  is a matching if each node appears in at most one edge in  $M$ .
- ▶ The **maximum matching problem** is to find the matching with the most edges.
- ▶ We'll design an efficient algorithm for maximum matching in a bipartite graph.

## Formulating it as a network flow problem

- ▶ **Goal:** given matching instance  $G = (L \cup R, E)$ , create a flow network  $G'$ , find a maximum flow  $f$  in  $G'$ , and use  $f$  to construct a maximum matching  $M$  in  $G$ . **Exercise.**
  - ▶ Add a source  $s$  and sink  $t$
  - ▶ For each edge  $(u, v) \in E$ , add a directed edge from  $u$  to  $v$  with capacity 1
  - ▶ Add an edge with capacity 1 from  $s$  to each node  $u \in L$
  - ▶ Add an edge with capacity 1 from each node  $v \in R$  to  $t$ .
  - ▶ Run F-F to get an **integral** max-flow  $f$
  - ▶ Set  $M$  to the set of edges from  $L$  to  $R$  with flow  $f(e) = 1$
- ▶ **Claim:** The set  $M$  is a maximum matching.

## Correctness

There is a bijection between integral flows  $f$  of value  $k$  and matchings  $M$  of size  $k$

1. Integral flow  $f$  of value  $k \Rightarrow$  matching  $M$  of size  $k$ 
  - ▶ Suppose  $f$  is a flow of value  $k$
  - ▶ Let  $M =$  edges from  $L$  to  $R$  carrying one unit of flow
  - ▶ There are  $k$  such edges, because the net flow across cut between  $L$  and  $R$  is  $k$ , and there are no edges from  $R$  to  $L$
  - ▶ There is at most 1 unit of flow entering  $u \in L$ , and therefore at most 1 unit of flow leaving  $u$
  - ▶ Since all flow values are 0 or 1, this means  $M$  has at most one edge incident to  $u$ .
  - ▶ A similar argument for  $v \in R$  means that  $M$  has at most one edge incident to  $v$
  - ▶ Therefore,  $M$  is a matching with size  $k$

## Correctness

2. Matching  $M$  of size  $k \Rightarrow$  integral flow  $f$  of value  $k$

- ▶ Suppose  $M$  is a matching of size  $k$
- ▶ Send one unit of flow from  $s$  to  $u \in L$  if  $u$  is matched
- ▶ Send one unit of flow from  $v \in R$  to  $t$  if  $t$  is matched
- ▶ Send one unit of flow on  $e$  if  $e$  is in  $M$
- ▶ All other edge flow values are zero
- ▶ Verify that capacity and flow conservation constraints are satisfied, and that  $v(f) = k$ .

For every integer flow of value  $k$  we can construct a matching  $M$  of size  $k$  and vice versa. Therefore, a maximum integer-valued flow yields a maximum matching.

## Second Application of Network Flows: Image Segmentation

- ▶ Using an expensive camera and appropriate lenses, you can get a "bokeh" effect on portrait photos in which the background is blurred and the foreground is in focus.



- ▶ But using cheap cameras in phones and appropriate software you can fake this effect...

## Formulating the problem

**Problem:** given set  $V$  of pixels, classify each as foreground or background. Assume you have:

- ▶ Numeric “cost” for assigning each pixel foreground/background
- ▶ Numeric penalty for assigning neighboring pixels to different classes

**Sketch of approach:** other slides, board work, demo.