CS 312: Algorithms **Spring 2018**

Discussion 4

Your Name: _____ Collaborators: _____

You will be randomly assigned groups to work on these problems in discussion section.

Problem 1. Graphs, proofs. The degree of node v in a graph is the number of edges incident to v (equivalently, the number of neighbors of v). Consider the claim and proof below. Identify at least two statements or phrases in the proof that are imprecise or not fully justified.

Claim: Let G be an undirected graph where every node has degree two or more. Then G has a cycle.

Proof: Since every node has two or more neighbors, we can construct a path starting from any node and keep going on that path without ever turning around. The path must eventually loop back on itself, so there is a cycle.

Now discuss with your group how you could make this more precise.

Problem 2. BFS. Recall the definition of BFS. Here is a concrete description from the book pp. 90-91.

- Explore out from node s by distance
- Layer i = nodes at distance i from s
- Returns a tree T
- Runs in $O(n+m)$ time.

Run BFS on the examples drawn on the board. Clearly identify the tree T, the layers $L_0, L_2, \ldots$, and the non-tree edges.

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BFS(s):
  Set Discovered[s] = true
  Set Discovered[v] = false for all other v
  Set L[0] = {s}
  Set i = 0                                  ▷ layer counter
  Set T = ∅                                  ▷ BFS tree
  while L[i] is not empty do
    Initialize an empty list L[i+1]          ▷ next layer
    for each node u ∈ L[i] do
      for each edge (u, v) incident to u do
        if Discovered[v]= false then
          Set Discovered[v] = true
          Add edge (u, v) to T
          Add v to L[i+1]
        end if
      end for
    end for
    Set i = i + 1
  end while
```

Problem 3. (Graph Traversal) K&T Chapter 3, Exercise 4.

This is the "butterfly problem" from the homework. Hint: try to modify BFS.