Classification & Information Theory

Lecture #8

Computational Linguistics CMPSCI 591N, Spring 2006

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Today's Main Points

- Automatically categorizing text
 - Parameter estimation and smoothing
 - a general recipe for a statistical CompLing model
 - Building a Spam Filter
- Information Theory
 - What is information? How can you measure it?
 - Entropy, Cross Entropy, Information gain

Maximum Likelihood Parameter Estimation Example: Binomial

- Toss a coin 100 times, observe *r* heads
- Assume a binomial distribution
 - Order doesn't matter, successive flips are independent
 - One parameter is q (probability of flipping a head)
 - Binomial gives p(r|n,q). We know r and n.
 - Find arg $max_q p(r|n, q)$

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(Notes for board) likelihood = $p(R = r \mid n, q) = \binom{n}{r} q^r (1-q)^{n-r}$ log - likelihood = $L = \log(p(r \mid n, q)) \propto \log(q^r (1-q)^{n-r}) = r \log(q) + (n-r) \log(1-q)$ $\frac{\partial L}{\partial q} = \frac{r}{q} - \frac{n-r}{1-q} \Rightarrow r(1-q) = (n-r)q \Rightarrow q = \frac{r}{n}$ Our familiar ratio-of-counts

is the maximum likelihood estimate!

Binomial Parameter Estimation Examples

- Make 1000 coin flips, observe 300 Heads
 - P(Heads) = 300/1000
- Make 3 coin flips, observe 2 Heads
 - P(Heads) = 2/3 ??
- Make 1 coin flips, observe 1 Tail
 - P(Heads) = 0 ???
- Make 0 coin flips
 - P(Heads) = ???
- We have some "prior" belief about P(Heads) before we see any data.
- After seeing some data, we have a "*posterior*" belief.

Maximum A Posteriori Parameter Estimation

• We've been finding the parameters that maximize

p(data|parameters),

not the parameters that maximize

- p(parameters|data) (parameters are random variables!)

• p(q|n,r) = p(r|n,q) p(q|n) = p(r|n,q) p(q)p(r|n) constant

• And let
$$p(q) = 2 q(1-q)$$

Maximum A Posteriori Parameter Estimation Example: Binomial

posterior =
$$p(r \mid n, q) p(q) = \binom{n}{r} q^r (1 - q)^{n-r} (2q(1 - q))$$

log - posterior = $L \propto \log(q^{r+1}(1 - q)^{n-r+1}) = (r+1)\log(q) + (n - r + 1)\log(1 - q)$
 $\frac{\partial L}{\partial q} = \frac{(r+1)}{q} - \frac{(n - r + 1)}{1 - q} \Rightarrow (r+1)(1 - q) = (n - r + 1)q \Rightarrow q = \frac{r+1}{n+2}$

Bayesian Decision Theory

• We can use such techniques for choosing among models:

– Which among several models best explains the data?

Likelihood Ratio

 P(model1 | data) = P(data|model1) P(model1)
 P(model2 | data) = P(data|model2) P(model2)

...back to our example: French vs English

- p(French | glacier, melange) versus p(English | glacier, melange) ?
- We have real data for
 - Shakespeare's Hamlet
 - Charles Dickens' Oliver Twist
- p(Hamlet | "hand", "death")
 p(Oliver | "hand", "death")

Statistical Spam Filtering



Document Classification by Machine Learning



Work out Naïve Bayes formulation interactively on the board

Recipe for Solving a NLP Task Statistically

- 1) Data: Notation, representation
- 2) **Problem**: Write down the problem in notation
- **3) Model**: Make some assumptions, define a parametric model
- **4) Inference**: How to search through possible answers to find the best one
- 5) Learning: How to estimate parameters
- 6) Implementation: Engineering considerations for an efficient implementation

(Engineering) Components of a Naïve Bayes Document Classifier

- Split documents into training and testing
- Cycle through all documents in each class
- Tokenize the character stream into words
- Count occurrences of each word in each class
- Estimate P(w|c) by a ratio of counts (+1 prior)
- For each test document, calculate P(c|d) for each class
- Record predicted (and true) class, and keep accuracy statistics

A Probabilistic Approach to Classification: "Naïve Bayes"

Pick the most probable class, given the evidence:

$$c^* = \operatorname{argmax}_{c_j} \Pr(c_j \mid d)$$

d - a document (like "language intelligence proof...")

Bayes Rule: $Pr(c_{j} \mid d) = \frac{Pr(c_{j})Pr(d \mid c_{j})}{Pr(d)} \approx \frac{Pr(c_{j})\prod_{i=1}^{ld}Pr(w_{d_{i}} \mid c_{j})}{\sum_{c_{k}}Pr(c_{k})\prod_{i=1}^{ld}Pr(w_{d_{i}} \mid c_{k})}$ $w_{d_{i}} - \text{the } i \text{ th word in } d \text{ (like "proof")}$

Parameter Estimation in Naïve Bayes

Estimate of P(c)

$$P(c_j) = \frac{1 + \operatorname{Count}(d \in c_j)}{|C| + \sum_k \operatorname{Count}(d \in c_k)}$$

Estimate of P(w|c)

$$P(w_i | c_j) = \frac{1 + \sum_{d_k \in c_j} \text{Count}(w_i, d_k)}{|V| + \sum_{t=1}^{|V|} \sum_{d_k \in c_j} \text{Count}(w_t, d_k)}$$

Information Theory

What is Information?

- "The sun will come up tomorrow."
- "Condi Rice was shot and killed this morning."

Efficient Encoding

- I have a 8-sided die. How many bits do I need to tell you what face I just rolled?
- My 8-sided die is unfair
 - P(1)=1/2, P(2)=1/8, P(3)=...=P(8)=1/16



Entropy (of a Random Variable)

- Average length of message needed to transmit the outcome of the random variable.
- First used in:
 - Data compression
 - Transmission rates over noisy channel

"Coding" Interpretation of Entropy

- Given some distribution over events P(X)...
- What is the average number of bits needed to encode a message (a event, string, sequence)
- = Entropy of P(X):

$$H(p(X)) = -\sum_{x \in X} p(x) \log_2(p(x))$$

• Notation: $H(X) = H_p(X)=H(p)=H_X(p)=H(p_X)$

What is the entropy of a fair coin? A fair 32-sided die? What is the entropy of an unfair coin that always comes up heads? What is the entropy of an unfair 6-sided die that always {1,2} Upper and lower bound? (Prove lower bound?)

Entropy and Expectation

- Recall $E[X] = \sum_{x \in X(\Omega)} x \cdot p(x)$
- Then

$$\begin{split} &\mathsf{E}[\mathsf{-log}_2(\mathsf{p}(\mathsf{x}))] = \Sigma_{\mathsf{x} \,\in\, \mathsf{X}(\Omega)} \,\mathsf{-log}_2(\mathsf{p}(\mathsf{x})) \,\cdot\, \mathsf{p}(\mathsf{x}) \\ &= \mathsf{H}(\mathsf{X}) \end{split}$$

Entropy of a coin



Entropy, intuitively

- High entropy ~ "chaos", fuzziness, opposite of order
- Comes from physics:
 - Entropy does not go down unless energy is used
- Measure of uncertainty
 - High entropy: a lot of uncertainty about the outcome, uniform distribution over outcomes
 - Low entropy: high certainty about the outcome

Claude Shannon



1950

- Claude Shannon
 1916 2001
 Creator of Information Theory
- Lays the foundation for implementing logic in digital circuits as part of his <u>Masters</u> Thesis! (1939)
- "A Mathematical Theory of Communication" (1948)

Joint Entropy and Conditional Entropy

- Two random variables: X (space Ω), Y (Ψ)
- Joint entropy
 - no big deal: (X,Y) considered a single event: $H(X,Y) = - \sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 p(x,y)$
- Conditional entropy: $H(X|Y) = - \sum_{x \in \Omega} \sum_{y \in \Psi} p(x,y) \log_2 p(x|y)$
 - recall that H(X) = E[-log₂(p(x))]
 (weighted average, and <u>weights are not conditional</u>)
 - How much extra information you need to supply to transmit X given that the other person knows Y.

Conditional Entropy (another way)

$$H(Y | X) = \sum_{x} p(x)H(Y | X = x)$$

= $\sum_{x} p(x)(-\sum_{y} p(y | x)\log_2(p(y | x)))$
= $-\sum_{x} \sum_{y} p(x)p(y | x)\log_2(p(y | x)))$
= $-\sum_{x} \sum_{y} p(x, y)\log_2(p(y | x))$

Chain Rule for Entropy

• Since, like random variables, entropy is based on an expectation..

H(X, Y) = H(X|Y) + H(X)

H(X, Y) = H(Y|X) + H(Y)

Cross Entropy

- What happens when you use a code that is sub-optimal for your event distribution?
 - I created my code to be efficient for a fair 8-sided die.
 - But the coin is unfair and always gives 1 or 2 uniformly.
 - How many bits on average for the optimal code?
 How many bits on average for the sub-optimal code?

$$H(p,q) = -\sum_{x \in X} p(x) \log_2(q(x))$$

KL Divergence

• What are the average number of bits that are wasted by encoding events from distribution *p* using distribution *q*?

$$D(p || q) = H(p,q) - H(p)$$

= $-\sum_{x \in X} p(x) \log_2(q(x)) + \sum_{x \in X} p(x) \log_2(p(x))$
= $\sum_{x \in X} p(x) \log_2(\frac{p(x)}{q(x)})$

A sort of "distance" between distributions *p* and *q*, but It is not symmetric! It does not satisfy the triangle inequality!

Mutual Information

- Recall: H(X) = average # bits for me to tell you which event occurred from distribution P(X).
- Now, first I tell you event y ∈ Y, H(X|Y) = average # bits necessary to tell you which event occurred from distribution P(X)?
- By how many bits does knowledge of Y lower the entropy of X?

$$I(X;Y) = H(X) - H(X | Y)$$

= $H(X) + H(Y) - H(X,Y)$
= $-\sum_{x} p(x) \log_2 \frac{1}{p(x)} - \sum_{y} p(y) \log_2 \frac{1}{p(y)} + \sum_{x,y} p(x,y) \log_2 p(x,y)$
= $\sum_{x,y} p(x,y) \log_2 \frac{p(x,y)}{p(x)p(y)}$

Mutual Information

- Symmetric, non-negative.
- Measure of independence.
 - I(X;Y) = 0 when X and Y are independent
 - I(X;Y) grows both with degree of dependence and entropy of the variables.
- Sometimes also called "information gain"

- Used often in NLP
 - clustering words
 - word sense disambiguation
 - feature selection...

Pointwise Mutual Information

- Previously measuring mutual information between two random variables.
- Could also measure mutual information between two events

$$I(x, y) = \log \frac{p(x, y)}{p(x)p(y)}$$

Continuing Homework Assignment?

- EC for whatever you hand in by tonight.
- For next week:
 - More time to create your own grammar
 - Modify parser to keep trace, and print parse trees
 - Try an additional grammar of your own creation, and investigate ambiguities
 - Work in small teams!