

Regular Languages

Lecture #2

Computational Linguistics CMPSCI 591N, Spring 2006

University of Massachusetts Amherst



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A Little About Yourself

- Have you programmed before?
 - Almost none at all.
 - Not much.
 - I work for a software company.
 - Fortran, C, C++, C#, Lisp, Perl, Python, Java,...
 - Only Basic on my Tandy 286!

A Little About Yourself

- Hobbies?
 - Fencing!
 - Hiking, Singing, Cooking, Poker, ...
 - Working on machines, like cars, motorcycles, airplanes.
 - Drinking, Smoking.
 - Fencing.
 - Watching movies, especially awesomely bad ones.

A Little About Yourself

- Favorite authors:
 - Kurt Vonnegut, George Orwell, Noam Chomsky
 - Asimov, Tolkein, Pinger,
 - I avoid reading, sorry.
 - Tolkein (x6), CS Lewis, etc.
 - Stroustrup
 - Arthur C. Clark
 - Hemmingway, x2
 - Salman Rushdie
 - Obscure foreign names like Savyon Librecht.
 - Karel Capek, Milan Kundera, Bulgahov.

A Little About Yourself

- Why are you in the class?
 - Practical skills to help in my linguistic research: accessing data, building grammars...
 - Interested in how probabilistic methods can be integrated with algebraic grammars.
 - Possibilities of a computer that can make sense of language are very exciting!
 - I want to expand my knowledge of AI.
 - I want to focus my career in CL, especially translation.
 - Want to simulate “the mind’s big bang”.
 - I think this will help me get a job!

Today's Main Points

- Examples of computation helping in Linguistic goals
- What are regular languages, finite state automata and regular expressions?
- Writing regular expressions (in Python)
- Examples on several large natural language corpora
- Finite-state transducers, and morphology
- Homework assignment #1

Some brief history: 1950s

- Early CL on machines less powerful than pocket calculators.
- Foundational work on automata, formal languages, probabilities and information theory.
- First speech systems (Davis et al, Bell Labs).
- MT heavily funded by military, but basically just word substitution programs.
- Little understanding of natural language syntax, semantics, pragmatics.

Some brief history: 1960s

- Alvey report (1966) ends funding for MT in America - the lack of real results realized
- ELIZA (MIT): Fraudulent NLP in a simple pattern matcher psychologist
 - It's true, I am unhappy.
 - *Do you think coming here will make you not to be unhappy?*
 - I need some help; that much is certain.
 - *What would it mean to you if you got some help?*
 - Perhaps I could learn to get along with my mother.
 - *Tell me more about your family.*
- Early corpora: Brown Corpus (Kudera and Francis)

Some brief history: 1970s

- Winograd's SHRDLU (1971): existence proof of NLP (in tangled LISP code).
- Could interpret questions, statements commands.
 - Which cube is sitting on the table?
 - *The large green one which supports the red pyramid.*
 - Is there a large block behind the pyramid?
 - *Yes, three of them. A large red one, a large green cube, and the blue one.*
 - Put a small one onto the green cube with supports a pyramid.
 - *OK.*

Some brief history: 1980s

- Procedural --> Declarative (including logic programming)
- Separation of processing (parser) from description of linguistic knowledge.
- Representations of meaning: procedural semantics (SHRDLU), semantic nets (Schank), logic (perceived as answer; finally applicable to real languages (Montague))
- Perceived need for KR (Lenat and Cyc)
- Working MT in limited domains (METEO)

Some brief history: 1990s

- Resurgence of finite-state methods for NLP: in practice they are incredibly effective.
- Speech recognition becomes widely usable.
- Large amounts of digital text become widely available and reorient the field. The Web.
- Resurgence of probabilistic / statistical methods, led by a few centers, especially IBM (speech, parsing, Candide MT system), often replacing logic for reasoning.
- Recognition of *ambiguity* as key problem.
- Emphasis on machine learning methods.

Some brief history: 2000s

- A bit early to tell! But maybe:
 - Continued surge in probability, Bayesian methods of evidence combination, and joint inference.
 - Emphasis on meaning and knowledge representation.
 - Emphasis on discourse and dialog.
 - Strong integration of techniques, and levels: brining together statistical NLP and sophisticated linguistic representations.
 - Increased emphasis on unsupervised learning.

Examples of Computation Helping Linguistics



Figure 1: The Phaistos Disk (c. 1700BC). The disk is six inches wide, double-sided, and is the earliest known document printed with a form of movable type.

- Kevin Knight
“*A Computational Approach to Deciphering Unknown Scripts*”
Mayan Writing
Pronunciation model, by
Expectation Maximization
(which we will study in
about 5 weeks)

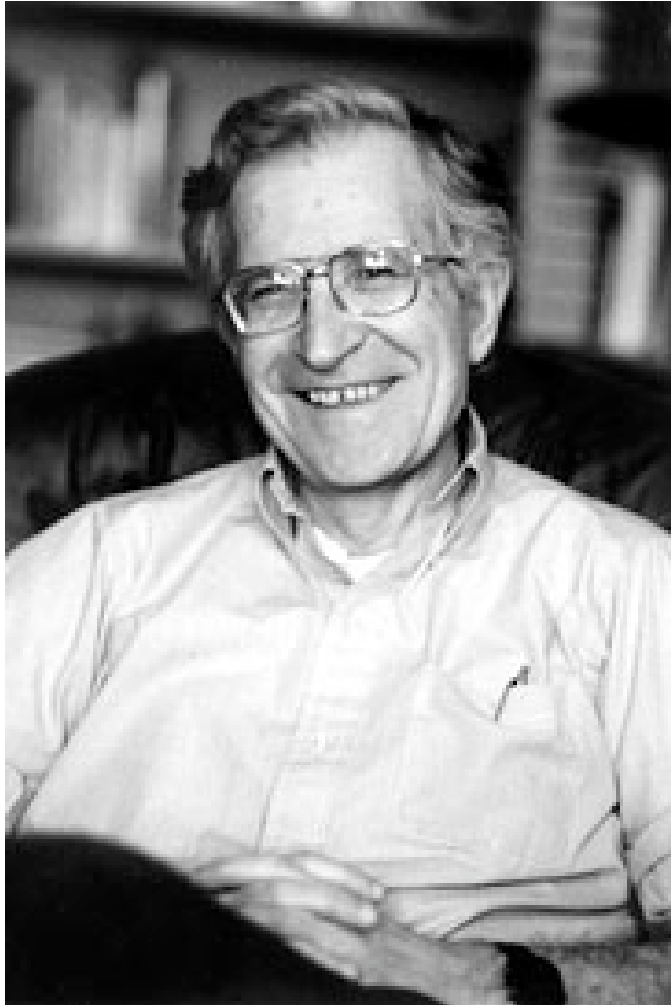
Examples of Computation Helping Linguistics

Other examples coming later:

- Learning Lexical Semantics
 - Augmenting WordNet by mining the Web.
- Automatically discovering English versus Japanese word order by grammar induction.
- Neural Network learners go through the same periods mistakes on irregular verbs as children do.
- ...and others.

Noun phrase parsing...?

Ed Hovy's thing?



Noam Chomsky
1928 -

Chomsky Hierarchy
Generative Grammar
Liberatarian-Socialist

The most cited person alive.

A Language

Some sentences in the language

- The man took the book. From [Chomsky, 1956], his first context-free parse tree.
- The purple giraffe hopped through the clouds.
- This sentence is false.

Some sentences not in the language

- *The girl, the sidewalk, the chalk, drew.
- *Backwards is sentence this.
- *loDvaD tlhIngan Hol ghojmoH be.

Compact description of a language

- Start with some “non-terminal” symbol, **S**.
- Expand that symbol, using some substitution **rules**.
- ...keep applying rules until all non-terminals are expanded to terminals.
- The string of terminals is in the sentence.

Chomsky Hierarchy

- Type 0 languages (Turing-equivalent)
Rewrite rules $a \rightarrow b$
where a, b are any string of terminals and non-terminals
- Context-sensitive languages
Rewrite rules $aXb \rightarrow acb$
where X is non-terminal and a, b as above
- Context-free languages
Rewrite rules $X \rightarrow a$
where X, a, b as above
- Regular languages
Rewrite rules $X \rightarrow aY$
where X, Y are non-terminals and a is a string of terminals

Linguistic
example:

ATNs

TAGs

PSGs

FSAs

*More detail on all this
again later.*

Regular language example

- Non-terminals:
 - **S, X, Y, Z**
- Terminals:
 - **m, o**
- Rules:
 - S** → **mX**
 - X** → **oY**
 - Y** → **o**
 - Y** →
- Start symbol:
 - S**

An expansion:

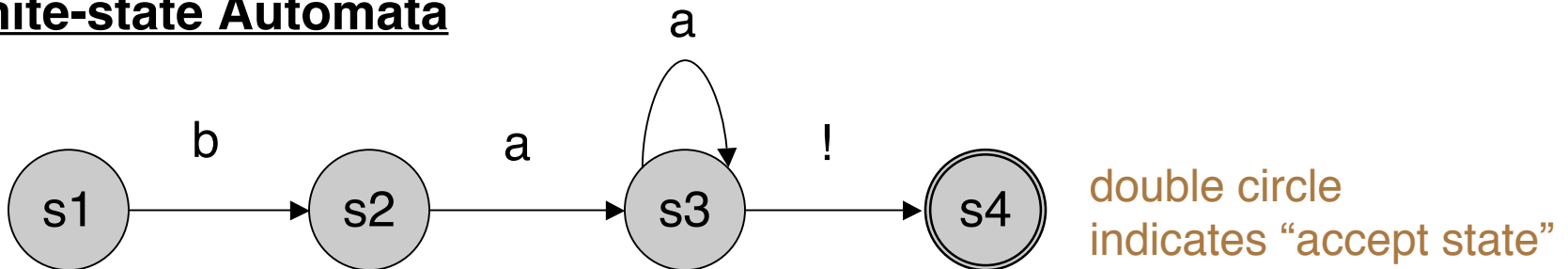
S
mX
moY
mooY
mooo

Example: Sheep Language

Strings in and out of the example Regular Language:

- In the language:
“ba!”, “baa!”, “baaaaa!”
- Not in the language:
“ba”, “b!”, “ab!”, “bbaaa!”, “alibaba!”

Finite-state Automata



Regular Expression

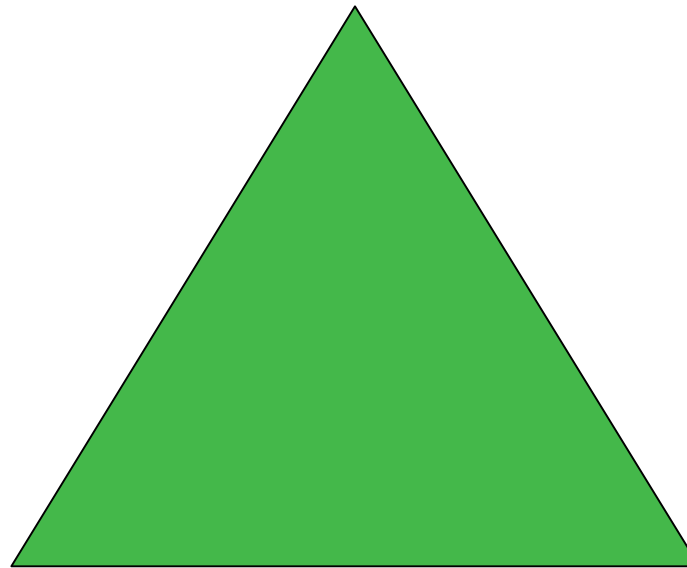
baa*

Recognizer

- A recognizer for a language is a program that takes as input a string W and answers “yes” if W is a sentence in the language, and answers “no” otherwise.
- We can think of this as a machine that emits only two possible responses it input.

Regular Languages: related concepts

Regular Languages
the accepted strings

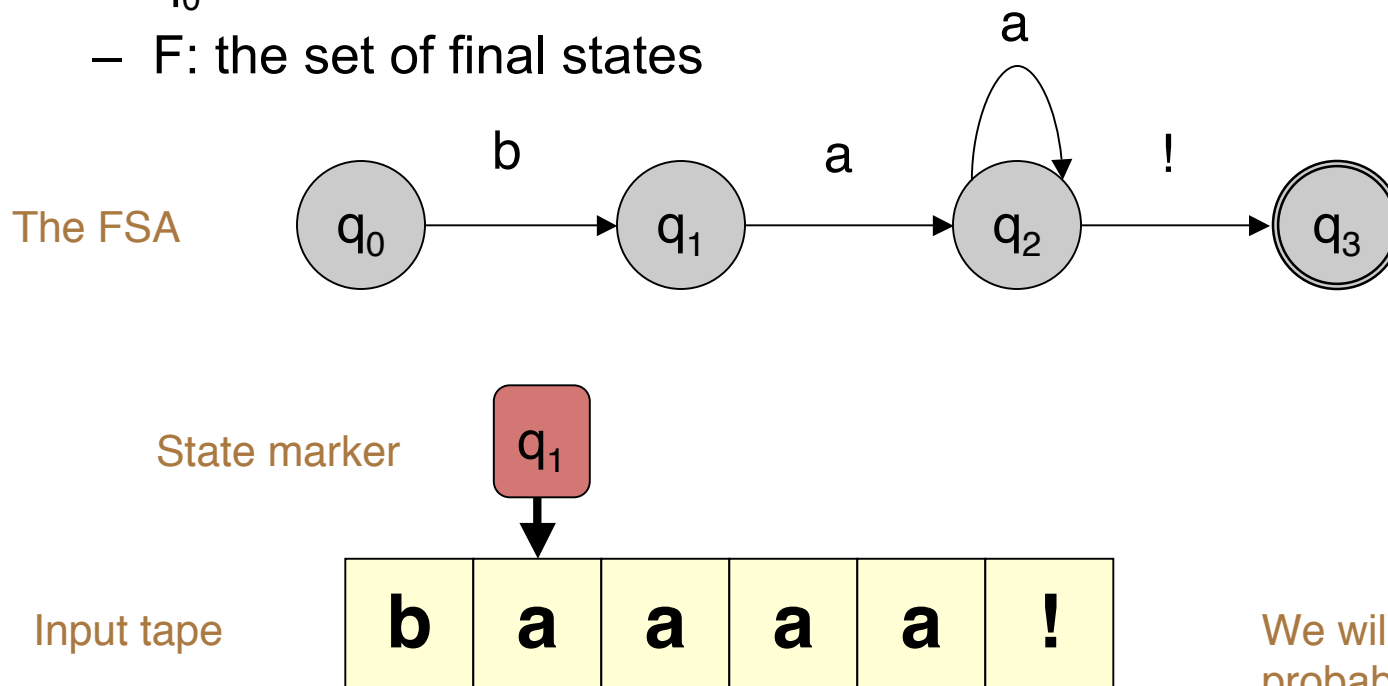


Finite-state Automata
machinery for accepting

Regular Expressions
a way to type the automata

Finite State Automata, more formally

- A finite state automata is a 5-tuple: $(Q, \Sigma, q_0, F, \delta(q,i))$
 - Q : finite set of N states, $q_0, q_1, q_2, \dots, q_N$ (non-terminals)
 - Σ : finite set of (terminals)
 - $\delta(q,i)$: transition function, given state and input, returns next state (production rules)
 - q_0 : the start state
 - F : the set of final states



We will later return to a probabilistic version of this with Hidden Markov Models!

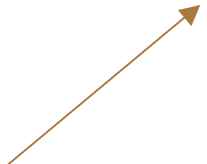
Transition Table, δ

State	Input		
	<i>b</i>	<i>a</i>	!
0	1	\emptyset	\emptyset
1	\emptyset	2	\emptyset
2	\emptyset	2	3
3	\emptyset	\emptyset	\emptyset

Regular Expressions

The “foundational” operations

	Pattern	Matches
Concatenation	abc	<i>abc</i>
Disjunction	a b	<i>a b</i>
	(a bb) d	<i>ad bbd</i>
Kleene star	a*	<i>ε a aa aaa ...</i>
	c (a bb) *	<i>ca cbba</i>


 The empty string

Regular expressions / Finite-state automata are “closed under these operations”

Stephen Kleene, 1909 - 1994



Attended Amherst College!

Best known for founding the branch of mathematical logic known as recursion theory, together with Alonzo Church, Kurt Godel, Alan Turing and others; and for inventing regular expressions.

“Kleeneliness is next to Godeliness.”

Practical Applications of RegEx's

- Web search
- Word processing, find, substitute
- Validate fields in a database (dates, email addr, URLs)
- Searching corpus for linguistic patterns
 - and gathering stats...
- Finite state machines extensively used for
 - acoustic modeling in speech recognition
 - information extraction (e.g. people & company names)
 - morphology
 - ...

Two types of characters in REs

- **Literal**
 - Every normal text character is an RE, and denotes itself.
- **Meta-characters**
 - Special characters that allow you to combine REs in various ways
 - Example:
 - a** denotes *a*
 - a*** denotes ϵ or *a* or *aa* or *aaa* or ...

Basic Regular Expressions

	Pattern	Matches
Character Concat	went	went
Alternatives	(go went)	go went
disjunc. negation	[aeiou]	a o u
wildcard char	[^aeiou]	b c d f g
	.	a z &
Loops & skips	a*	ϵ a aa aaa ...
one or more	a+	a aa aaa
zero or one	colou?r	color colour

More Fancy Regular Expressions

- Special characters

- `\t` tab `\v` vertical tab
- `\n` newline `\r` carriage return

- Aliases (shorthand)

- `\d` digits `[0-9]`
- `\D` non-digits `[^0-9]`
- `\w` alphabetic `[a-zA-Z]`
- `\W` non-alphabetic `[^a-zA-Z]`
- `\s` whitespace `[\t\n\r\f\v]`
- `\w` alphabetic `[a-zA-Z]`

- Examples

- `\d+ dollars` 3 dollars, 50 dollars, 982 dollars
- `\w*oo\w*` food, boo, oodles

- Escape character

- `\` is the general escape character; e.g. `\.` is not a wildcard, but matches a period `.`
- if you want to use `\` in a string it has to be escaped `\\`

Yet More Fancy Regular Expressions

- **Anchors**. AKA, “zero width characters”.
- They match positions in the text.
 - `^` beginning of line
 - `$` end of line
 - `\b` word boundary, i.e. location with `\w` on one side but not on the other.
 - `\B` ???
- Examples:
 - `\bthe\b` the together
- **Counters** `{1}`, `{1,2}`, `{3,}`

Even More Fancy Regular Expressions

- **Grouping**
 - **a (good|bad) movie**
 - **He said it (again and)*again.**
- Parens also indicate **Registers** (saved contents)
 - **b(\w+)b\1**
matches *boohoo* and *baha*, but not *boohaa*
The digit after the \ indicates which of multiple paren groups, as ordered by when then were opened.
- Grouping without the cost of register saving
 - **He went (?:this|that) way.**

Extra Fancy Regular Expressions

- Non-consuming tests
 - (? = . . .) - Positive lookAHEAD
 - (? ! . . .) - Negative lookAHEAD
 - (? < = . . .) - Positive lookBEHIND
 - (? < ! . . .) - Negative lookBEHIND
- Example
 - (? < = [H h] e) \ w + e d (? = \ w + l y)

Yowza! Regular Expressions

- Multi-line, with comments, etc.

```
FROM = re.compile( """
    ^From:           # Anchor to start of line
    \s*             # maybe some spaces
    (?P<user>\w+)   # 'user': group of word characters
    @
    (?P<domain>    # the 'domain':
    \S+)          # some non-space characters
    \s             # finally, a space character
    """, re.VERBOSE)
```

Oral Quiz:

Describe the strings these will accept

- `ab+a`
- `(ab)*`
- `([^aeiou][aeiou])\1`
- `\bdis\w+\b...`

Hands-on!

With little strings

- re_show function from NLTK

```
import re
def re_show(regex, string):
    print re.compile(regex, re.M).sub("{\g<0>}", string.rstrip())
```

Hands-on!

Investigating corpora

- `regexs.py`
- `regexcount.py`
- Corpora
 - 200k+ words of Wall Street Journal articles
 - 3 books by Jane Austin
 - Several plays by William Shakespeare
 - Several books by Mark Twain
 - and various other text from Walt Whitman, & others...

I will make these available on the course Web site

ELIZA (Weizenbaum 1966)

User: Men are all alike.

ELIZA: IN WHAT WAY?

User: They're always bugging us about something or other.

ELIZA: CAN YOU THINK OF A SPECIFIC EXAMPLE?

User: Well, my boyfriend made me come here.

ELIZA: YOUR BOYFRIEND MADE YOU COME HERE.

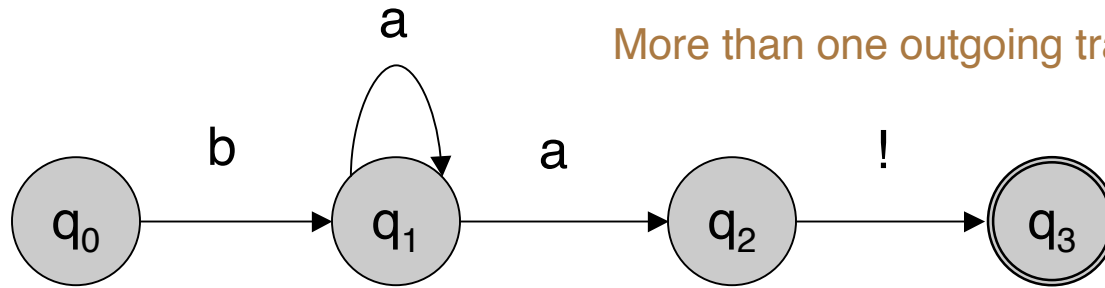
User: He say's I'm depressed much of the time.

ELIZA: I'M SORRY TO HEAR THAT YOU ARE DEPRESSED.

Implemented with regular expression substitution!

```
s/. * YOU ARE (depressed|sad) .*/I AM SORRY TO HEAR THAT YOU ARE \1/  
s/. * always .*/CAN YOU THINK OF A SPECIFIC EXAMPLE?/
```


Non-deterministic FSAs



More than one outgoing transition on 'a'

State	Input		
	<i>b</i>	<i>a</i>	<i>!</i>
0	1	\emptyset	\emptyset
1	\emptyset	1,2	\emptyset
2	\emptyset	\emptyset	3
3	\emptyset	\emptyset	\emptyset

Transition *relation*, rather than transition function.

Non-deterministic finite-state automata as “Recognizers”

- The problem:

When processing a string, we might follow the wrong transition, and reject the string when we should have accepted it!

One solution: turn the NFA into a DFA... (See CMPSCI 250)

- Ubiquitous problem in this course:
How to efficiently search through various possible “paths” (parses) to find one that works / the most likely one, etc.

How do humans do this?!

Solutions

- **Look-ahead**
 - Peek ahead to help decide which path to take.
- **Parallelism**
 - At each choice, take every path in parallel.
- **Backup**
 - At each choice point, mark the input / state
 - If we fail, go back and try another path
Need a *stack* (or *queue*) of markers
 - Marker = “Machine state”
 - Collection of current state & markers = “Search state”
 - Depth-first search (or Breadth-first search).

“Smart” heuristic search, “A*”. See CMPSCI 383
(Artificial Intelligence)

RE / FSA equivalence proof

- How would you do it?

Morphology

The study of the sub-word units of meaning.

disconnect

“not”

“to attach”

Making a word plural:

If word is regular,

add s

If word ends in y,

change y to i, and add s

If word ends in x,

add -es

...

Examples:

dog dogs

baby babies

fox foxes

Recognizing that **foxes** breaks down into morphemes **fox** and **-es** called *Morphological Parsing*

Parsing = taking an input and producing some sort of structure for it.

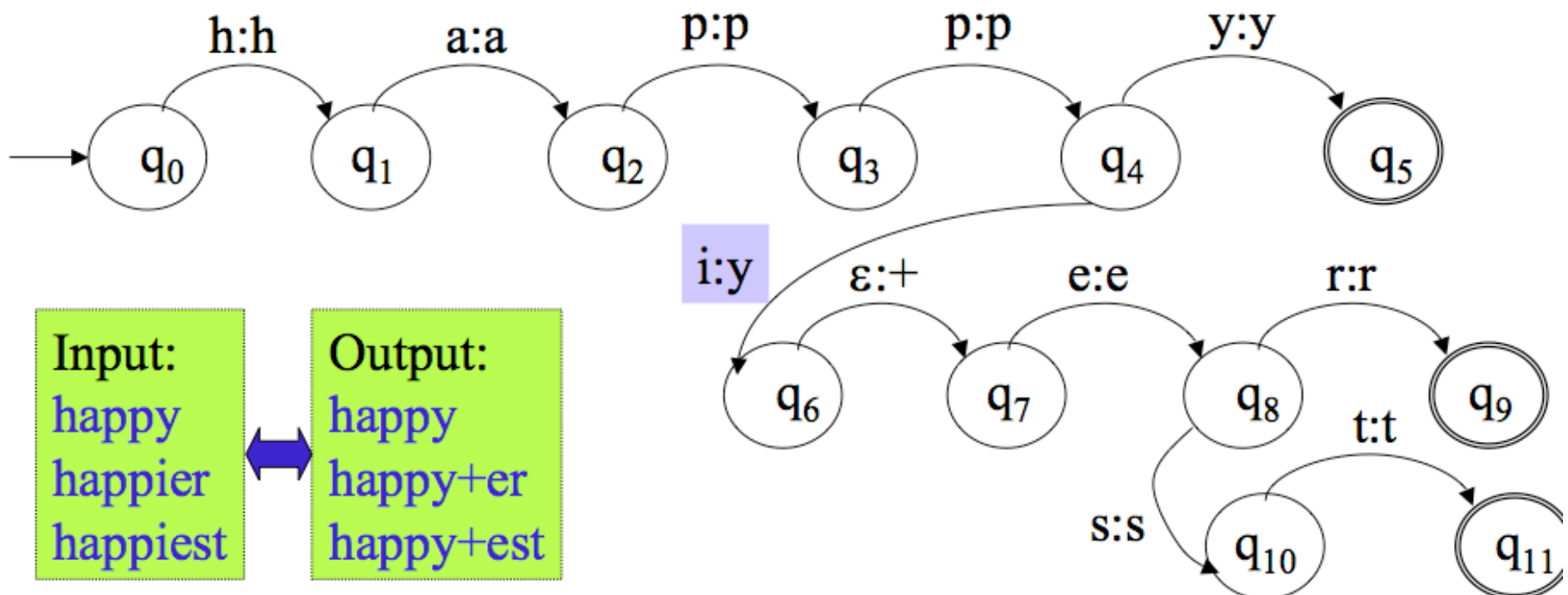
Morphology, briefly

- **morpheme**: minimal meaning-bearing unit
 - **stem**: “main” morpheme of a word, e.g. **fox**
 - **affixes**: add “additional” meanings, e.g. **+es**
includes **prefixes**, **suffixes**, **infixes**, **circumfixes**,
e.g. **un-**, **-ly**,
 - concatenative morphology, non-concatenative
- **inflection**: stem+morpheme in the same class as stem.
 - e.g. nouns plural **+s**, possessive **+’s**
- **derivation**: stem+morpheme in different class...
 - e.g. **+ly** makes and adverb from an adjective

Morphological Parsing with Finite State Transducers

- We want a system that given **foxes** will output a parse: **fox+es** or **fox +PL**
- FSAs will take input, but not produce output (other than “accept”/“reject”)
- Solution: **Finite State Transducers** (FST):
 - A FST is a two-tape automaton that recognizes or generates **pairs** of strings.

Example Finite-state Transducer



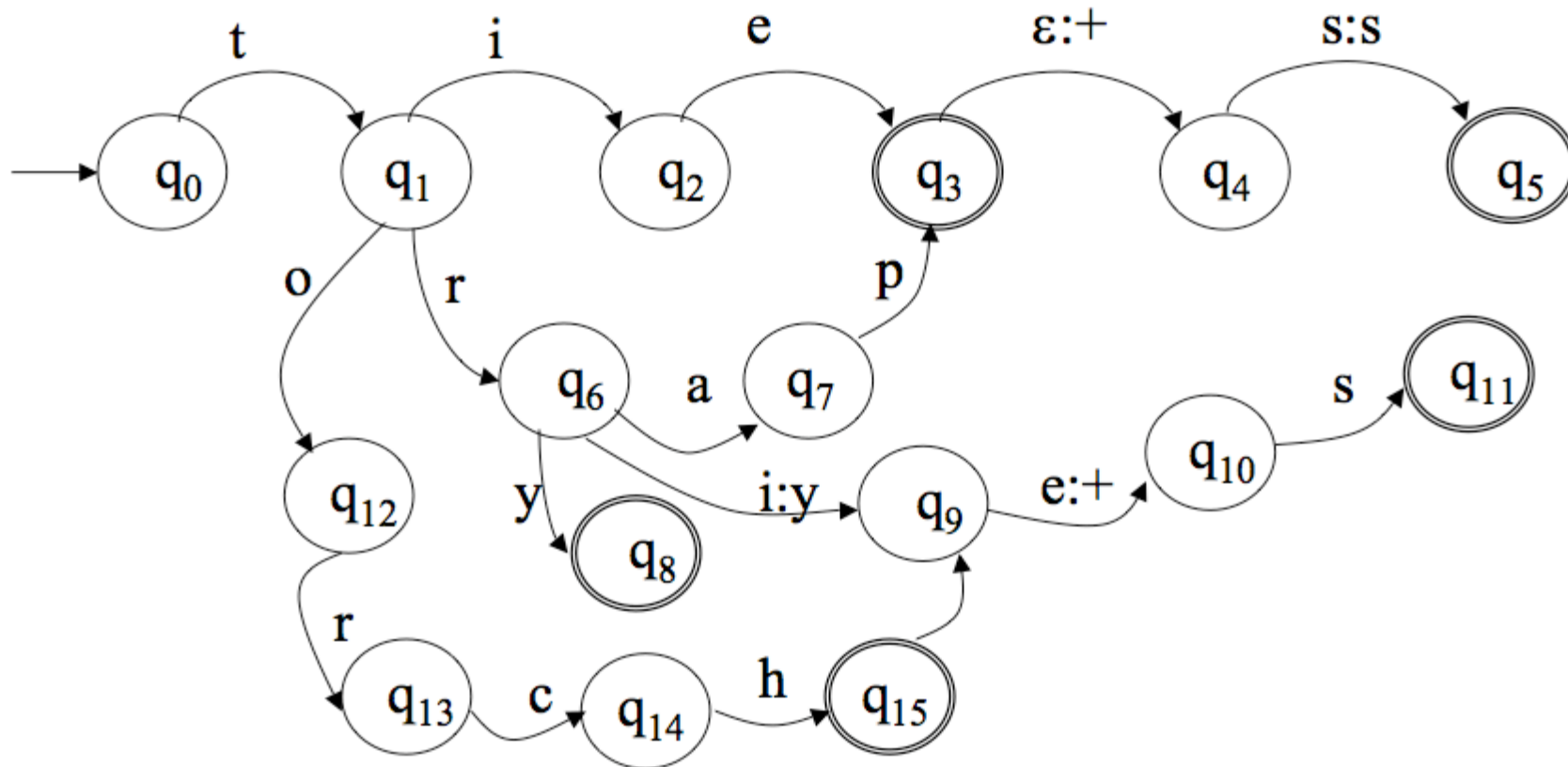
FSTs can be used to transform a word surface form into morphemes (or vice-versa!)

An entire lexicon can be encoded as a FST.

FST transition table

	Input									
State	h:h	a:a	p:p	y:y	i:y	ε:+	e:e	r:r	s:s	t:t
0	1	∅	∅	∅	∅	∅	∅	∅	∅	∅
1	∅	2	∅	∅	∅	∅	∅	∅	∅	∅
2	∅	∅	3	∅	∅	∅	∅	∅	∅	∅
3	∅	∅	4	∅	∅	∅	∅	∅	∅	∅
4	∅	∅	∅	5	6	∅	∅	∅	∅	∅
5:	∅	∅	∅	∅	∅	∅	∅	∅	∅	∅
6	∅	∅	∅	∅	∅	7	∅	∅	∅	∅
7	∅	∅	∅	∅	∅	∅	8	∅	∅	∅
8	∅	∅	∅	∅	∅	∅	10	9	∅	∅
9:	∅	∅	∅	∅	∅	∅	∅	∅	∅	∅

Fragment of a lexicon in a FST



Further Closure Properties of FSAs

Regular languages are also closed under the following operations

- **Reversal:** If L_1 is regular, so is the language consisting of the set of all reversals of strings in L_1 .
- **Intersection:** if L_1 and L_2 are regular languages, so is the language consisting of all strings that are in both L_1 and L_2 .
- **Difference:** If L_1 and L_2 are regular languages, so is the language consisting of all strings in L_1 that are not in L_2 .
- **Complementation:** If L_1 is a regular language, so is the set of all possible strings that are not in L_1 .

Announcement: Undergraduate CMPSSCI Meeting

- **“First Friday”**
 - Curriculum Information
 - Spring Events
 - Jobs/Co-ops/Research positions in and out of the Department
 - Library Carrels
 - And More!
- **Friday, February 3, 2005**
3:30 - 5:00 PM
CMPS 150/151 (Computer Science Building)
Refreshments will be served.

Next class (Tuesday Feb 7)

- Learning Python
 - Variables, operators, conditionals, iteration, etc.
 - functions, classes, modules
 - Gather statistics from Python-ized Penn Treebank.
 - Calculate statistics from 200k words of WSJ
 - Implement a phrase structure grammar, and generate sentences from it.
- **Install Python, and bring your laptop with you!**

First Homework, assigned today!

- Essentially:
 - Write some regular expressions
 - Run them on some corpora
 - Write ~1 page about your experience and findings
 - Extra credit for creativity and interesting application!
- Feel free to come do it in office hours!
- Due next Thursday, one week from today.
(Don't wait until Wednesday to install Python!)
- Recommended schedule:
 - Idea by Saturday
 - Coded/tested by Monday
 - Write-up by Wednesday

Office Hours, CS Building, Rm 264

- Friday, 2-4pm
- Monday, 10:30am-1pm
- Tuesday, 10:30am-1pm
- Wednesday, 10:30am-1pm
- Thursday, 10:30am-12:30pm

- If you can't make these times, let me know.

Aside: Grammar Induction

- Also called “Grammatical Inference”
- “Learning” finite-state automata from many examples of strings in (and out of) the language.
- <http://www.info.ucl.ac.be/~pdupont/pdupont/gram.html>
- Learning FSA and CFG structure from data!

Thank you!